

CS 441 Discrete Mathematics for CS
Lecture 23

Relations IV

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Closures on relations

- **Relations can have different properties:**
 - **reflexive,**
 - **symmetric**
 - **transitive**
 - **Because of that we can have:**
 - **symmetric,**
 - **reflexive and**
 - **transitive**
- closures.**

Closures

Definition: Let R be a relation on a set A . A relation S on A with property P is called **the closure of R with respect to P** if S is a subset of every relation Q ($S \subseteq Q$) with property P that contains R ($R \subseteq Q$).

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Example (a symmetric closure):

- Assume $R = \{(1,2), (1,3), (2,2)\}$ on $A = \{1,2,3\}$.
- What is the symmetric closure S of R ?
- $S = \{(1,2), (1,3), (2,2)\} \cup \{(2,1), (3,1)\}$
 $= \{(1,2), (1,3), (2,2), (2,1), (3,1)\}$

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Example (transitive closure):

- Assume $R = \{(1,2), (2,2), (2,3)\}$ on $A = \{1,2,3\}$.
- **Is R transitive? No.**
- **How to make it transitive?**
- **$S = ?$**

Closures

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Example (transitive closure):

- Assume $R = \{(1,2), (2,2), (2,3)\}$ on $A = \{1,2,3\}$.
- **Is R transitive? No.**
- **How to make it transitive?**
- **$S = \{(1,2), (2,2), (2,3)\} \cup \{(1,3)\}$**
 $= \{(1,2), (2,2), (2,3), (1,3)\}$
- **S is the transitive closure of R**

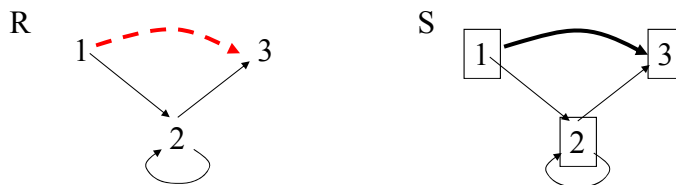
Transitive closure

We can represent the relation on the graph. Finding a transitive closure corresponds to finding all pairs of elements that are connected with a directed path (or digraph).

Example:

Assume $R = \{(1,2), (2,2), (2,3)\}$ on $A = \{1,2,3\}$.

Transitive closure $S = \{(1,2), (2,2), (2,3), (1,3)\}$.



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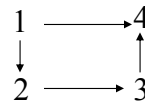
Connectivity relation

Definition: Let R be a relation on a set A . The **connectivity relation** R^* consists of all pairs (a,b) such that there is a path (of any length, ie. 1 or 2 or 3 or ...) between a and b in R .

$$R^* = \bigcup_{k=1}^{\infty} R^k$$

Example:

- $A = \{1,2,3,4\}$
- $R = \{(1,2), (1,4), (2,3), (3,4)\}$
- $R^2 = \{(1,3), (2,4)\}$
- $R^3 = \{(1,4)\}$
- $R^4 = \emptyset$
- ...
- $R^* = \{(1,2), (1,3), (1,4), (2,3), (2,4), (3,4)\}$



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Transitivity closure and connectivity relation

Theorem: The transitive closure of a relation R **equals** the connectivity relation R^* .

Lemma: Let A be a set with n elements, and R a relation on A. If there is a path from a to b, then there exists a path of length $< n$ in between (a,b). Consequently:

$$R^* = \bigcup_{k=1}^n R^k$$

Equivalence relation

Definition: A relation R on a set A is called an **equivalence relation** if it is **reflexive, symmetric and transitive**.

Example: Let $A = \{0,1,2,3,4,5,6\}$ and

- $R = \{(a,b) \mid a,b \in A, a \equiv b \pmod{3}\}$ (a is congruent to b modulo 3)

Congruencies:

- $0 \pmod{3} = 0 \quad 1 \pmod{3} = 1 \quad 2 \pmod{3} = 2 \quad 3 \pmod{3} = 0$
- $4 \pmod{3} = 1 \quad 5 \pmod{3} = 2 \quad 6 \pmod{3} = 0$

Relation R has the following pairs:

- $(0,0) \quad (0,3), (3,0), (0,6), (6,0)$
- $(3,3), (3,6), (6,3), (6,6) \quad (1,1), (1,4), (4,1), (4,4)$
- $(2,2), (2,5), (5,2), (5,5)$

Equivalence relation

Example:

- Suppose that R is the relation on the set of strings of English letters such that aRb if and only if $l(a) = l(b)$, where $l(x)$ is the length of the string x . Is R an equivalence relation?
- **Is the relation reflexive?**

Equivalence relation

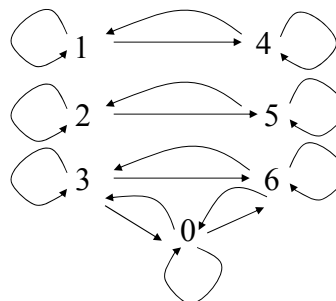
- **Relation R on $A=\{0,1,2,3,4,5,6\}$ has the following pairs:**

(0,0) (0,3), (3,0), (0,6), (6,0)
(3,3), (3,6), (6,3), (6,6) (1,1), (1,4), (4,1), (4,4)
(2,2), (2,5), (5,2), (5,5)

- Is R reflexive? **Yes.**
- Is R symmetric? **Yes.**
- Is R transitive. **Yes.**

Then

- **R is an equivalence relation.**



Equivalence relation

Example:

- Suppose that R is the relation on the set of strings of English letters such that aRb if and only if $l(a) = l(b)$, where $l(x)$ is the length of the string x . Is R an equivalence relation?
- **Is the relation reflexive? Yes.**
 - Because $l(a) = l(a)$, it follows that aRa for all strings a .
- **Is the relation symmetric? Yes.**
 - Suppose that aRb . Since $l(a) = l(b)$, $l(b) = l(a)$ also holds and bRa .
- **Is the relation transitive? Yes.**
 - Suppose that aRb and bRc . Since $l(a) = l(b)$, and $l(b) = l(c)$, $l(a) = l(c)$ also holds and aRc .
- **Conclusion: R an equivalence relation**

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Equivalence class

Definition: Let R be an equivalence relation on a set A . The set $\{x \in A \mid a R x\}$ is called **the equivalence class of a** , denoted by $[a]_R$ or simply $[a]$ when there is only one relation R . If $b \in [a]$ then b is called **a representative of this equivalence class**.

Example:

- Assume $R = \{(a,b) \mid a \equiv b \pmod{3}\}$ for $A = \{0,1,2,3,4,5,6\}$
- Pick an element $a = 0$.
- $[0]_R = \{0,3,6\}$
- Element 1: $[1]_R = \{1,4\}$
- Element 2: $[2]_R = \{2,5\}$
- Element 3: $[3]_R = \{0,3,6\} = [0]_R = [6]_R$
- Element 4: $[4]_R = \{1,4\} = [1]_R$
- Element 5: $[5]_R = \{2,5\} = [2]_R$

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Equivalence class

Example:

- Assume $R = \{(a,b) \mid a \equiv b \pmod{3}\}$ for $A = \{0,1,2,3,4,5,6\}$

Three different equivalence classes all together:

- $[0]_R = [3]_R = [6]_R = \{0,3,6\}$
- $[1]_R = [4]_R = \{1,4\}$
- $[2]_R = [5]_R = \{2,5\}$

Partition of a set S

Definition: Let S be a set. A collection of nonempty subsets of S A_1, A_2, \dots, A_k is called a **partition of S** if:

- $A_i \cap A_j = \emptyset, i \neq j$ and $S = \bigcup_{i=1}^k A_i$

Example: Let $S = \{1,2,3,4,5,6\}$ and

- $A_1 = \{0,3,6\}$ $A_2 = \{1,4\}$ $A_3 = \{2,5\}$
- Is A_1, A_2, A_3 a partition of S ?

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- $A_1 = \{0, 3, 6\}$ $A_2 = \{1, 4\}$ $A_3 = \{2, 5\}$
- Is A_1, A_2, A_3 a partition of S ? **Yes.**
- Give a partition of S ?

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- $A_1 = \{0, 3, 6\}$ $A_2 = \{1, 4\}$ $A_3 = \{2, 5\}$
- Is A_1, A_2, A_3 a partition of S ? **Yes.**
- Give a partition of S ?
- $\{0, 2, 4, 6\}$ $\{1, 3, 5\}$
- $\{0\}$ $\{1, 2\}$ $\{3, 4, 5\}$ $\{6\}$

Equivalence classes

Theorem: Let R be an equivalence relation on a set A . The union of all the equivalence classes of R is all of A , since an element a of A is in its own equivalence class $[a]_R$. In other words,

$$A = \bigcup_{a \in A} [a]_R$$

Equivalence class

Theorem: Let R be an **equivalence relation** on a set A . The following statements are equivalent:

- i) $a R b$
- ii) $[a] = [b]$
- iii) $[a] \cap [b] \neq \emptyset$.

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- iii) $[a] \cap [b] \neq \emptyset$.

Proof: (i) \rightarrow (ii)

- Suppose $a R b$, i.e., $(a,b) \in R$. Want to show $[a] = [b]$.
- Let $x \in [a] \rightarrow (a,x) \in R$.
- Since R is symmetric $(b,a) \in R$.
- Since R is transitive, $(b,a) \in R$ and $(a,x) \in R \rightarrow (b,x) \in R$. Thus, $x \in [b]$.
- Let $x \in [b] \rightarrow (b,x) \in R$.
- Since R is transitive, $(a,b) \in R$ and $(b,x) \in R \rightarrow (a,x) \in R$. Thus $x \in [a]$.
- Therefore $[a] = [b]$.

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Theorem: Let R be an **equivalence relation** on a set A . The following statements are equivalent:

- i) $a R b$
- ii) $[a] = [b]$
- iii) $[a] \cap [b] \neq \emptyset$.

Proof: (ii) \rightarrow (iii)

- Suppose $[a] = [b]$. Want to show $[a] \cap [b] \neq \emptyset$.
- Since R is reflexive, $a \in [a] \rightarrow [a] \neq \emptyset$ and the result follows.

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Theorem: Let R be an **equivalence relation** on a set A . The following statements are equivalent:

- i) $a R b$
- ii) $[a] = [b]$
- iii) $[a] \cap [b] \neq \emptyset$.

Proof: (iii) \rightarrow (i)

- Suppose $[a] \cap [b] \neq \emptyset$, want to show $a R b$.
- $[a] \cap [b] \neq \emptyset \rightarrow x \in [a] \cap [b] \rightarrow x \in [a]$ and $x \in [b] \rightarrow (a,x)$ and $(b,x) \in R$.
- Since R is symmetric $(x,b) \in R$. By the transitivity of R $(a,x) \in R$ and $(x,b) \in R$ implies $(a,b) \in R \rightarrow a R b$.

Partial orderings

Definition: A relation R on a set S is called a *partial ordering*, or *partial order*, if it is reflexive, antisymmetric, and transitive. A set together with a partial ordering R is called a *partially ordered set*, or *poset*, and is denoted by (S, R) . Members of S are called *elements* of the poset.

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Example: Assume R denotes the “greater than or equal” relation (\geq) on the set $S = \{1, 2, 3, 4, 5\}$.

- Is the relation reflexive?

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- Is the relation reflexive? Yes
- Is it antisymmetric?

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- Is it transitive?

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- Is it transitive? Yes
- **Conclusion:** R is a partial ordering.

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Example: Assume R denotes the “greater than or equal” relation (\geq) on the set $S = \{1, 2, 3, 4, 5\}$.

- Is the relation reflexive? Yes
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- Is it transitive? Yes
- **Conclusion:** R is a partial ordering.

Partial orderings

Example: Assume R is the divisibility relation ($|$) on the set of integers $S = \{1, 2, 3, 4, 5, 6\}$

- Is the relation reflexive? Yes
- Is it antisymmetric? Yes
- Is it transitive? Yes
- **Conclusion:** R is a partial ordering.