



- □ Brief Clustering Review
- $\ \square$  Similarity Graph
- □ Graph Partitioning
- □ Graph Laplacian Properties
- □ Spectral Clustering Algorithm
- □ Graph Cut Point of View
- □ Random Walk Point of View
- Practical Details



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### Clustering Review



#### Clustering

Groups together "similar" instances in the data sample

#### Basic clustering problem:

- distribute data into k different groups such that data points similar to each other are in the same group
- Similarity between data points is defined in terms of some distance metric (can be chosen)

#### Clustering is useful for:

- Similarity/Dissimilarity analysis
   Analyze what data points in the sample are close to each other
- Dimensionality reduction
   High dimensional data replaced with a group (cluster) label

CS 2750 Machine Learning

### K-means Clustering



Given a set of observations  $(\mathbf{x}_1, \mathbf{x}_2, ..., \mathbf{x}_n)$ , where each observation is a d-dimensional real vector, k-means clustering aims to partition the n observations into k sets  $(k \le n)$   $\mathbf{S} = \{S_1, S_2, ..., S_k\}$  so as to minimize the within-cluster sum of squares (WCSS):

$$\underset{S}{\operatorname{arg\,min}} \sum_{i=1}^{k} \sum_{x_j \in S_i} \|x_j - u_i\|^2$$

where  $\mu_i$  is the mean of points in  $S_i$ .

### K-means Clustering



#### Standard Algorithm



1) k initial "means" (in this case k=3) are randomly selected from the data set.



2) k clusters are created by associating every observation with the nearest mean.



3) The centroid of each of the k clusters becomes the new means.

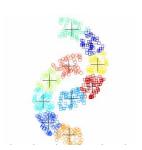


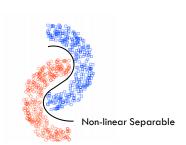
4) Steps 2 and 3 are repeated until convergence has been reached.

### K-means Shortcomings



- □ K-means
  - □ Criteria: Compactness of clusters
  - Major drawback:
    - It can't separate clusters that are non-linearly separable





 $\arg\min_{S} \sum_{i=1}^{k} \sum_{x_j \in S_i} \|x_j - u_i\|^2$ 

### K-means Shortcomings



- Spectral Clustering
  - □ Criteria: Connectivity









- Instead of linear similarities use non-linear distance/similarity of data or geodesic distance of points
  - Similarity Graph



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### General Idea

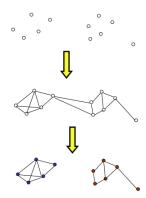


First – graph representation of data (largely, application dependent)

Second - graph partitioning:

#### **Partition Properties:**

- Weakly connections in between components
- Strongly connections within components



# Graph Notation



#### G=(V,E):

- $\quad \quad \square \quad \text{Vertex set} \qquad V = \{v_1, \dots, v_n\}$
- □ Weighted adjacency matrix:

$$v_m$$

$$W = \left(w_{ij}\right)i, j = 1, ..., n \quad w_{ij} \geq 0 \text{ and } w_{ij} = w_{ji}$$



figspace **Degree matrix D** Diagonal matrix with the degrees  $d_1,\ldots,d_n$  on the diagonal.

### **Graph Notation**



#### G=(V,E):

- □ A is a **subset** of V
  - $\bar{A} = V \setminus A$



- $\label{eq:location} \ensuremath{\blacksquare} \ensuremath{\text{Indicator Vector}} \quad \mathbb{1}_{A} = (f_1, \dots, f_n)' \in \ensuremath{\mathbb{R}}^n \quad f_i \in \{0, 1\}$
- $\square$  "Size" of a subset  $A \subseteq V$

$$|A| :=$$
the number of vertices in  $A$   $vol(A) := \sum_{i \in A} d_i$ 

Similarity/Connectivity of two subsets A and B

$$W(A,B) \coloneqq \sum_{i \in A, j \in B} w_{ij}$$

### **Graph Notation**



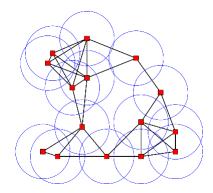


- □ **Connected Subset** A subset A of a graph is connected if any two vertices in A can be joined by a path such that all intermediate points also lie in A.
- fill Connected Component it is connected and if there are no connections between vertices in A and  $\bar{A}$ .
- **Partition:** A set of nonempty sets  $A_1,\ldots,A_k$  form a partition of the graph if  $A_i\cap A_j=\emptyset \text{ and } A_1\cup\cdots\cup A_k=V.$

# Similarity Graph Construction



- □ Goal: Model local neighborhood relationships between data points
- $\square$   $\varepsilon$ -neighborhood graph
  - lacksquare Connect all points whose pairwise distances are smaller than arepsilon
    - Considered unweighted graph

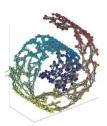




## Similarity Graph Construction

- □ k-nearest neighbor graph
  - lacksquare Connect vertex  $v_i$  with vertex  $v_i$  if  $v_i$  is among the k-nearest neighbors of  $v_i$ .
  - Is directed graph
    - Knn graph: obtained by loosing directions
    - Mutual knn graph: connect only if both vertices have the other in their knn





### Similarity Graph Construction



- Fully connected graph
  - Connect all points with positive similarity with each other
  - Not modeling local neighborhoods
    - We have to use similarity function has to do it (Gaussian Similarity Function)

$$W_{ij} = e^{\frac{\left|x_i - x_j\right|^2}{2\sigma^2}}$$

σ controls size of neighborhood

All of introduced graphs are regularly used in spectral clustering!



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### **Graph Partitioning**

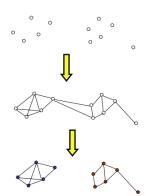


First – graph representation of data (largely, application dependent)

**Second**— graph partitioning:

#### **Partition Properties:**

- Weakly connections in between components
- Strongly connections within components



# Graph Cut



G=(V,E):

 $\hfill\Box$  For two not necessarily disjoint set  $A,B\subset V$  , we define

$$W(A,B) \coloneqq \sum_{i \in A, j \in B} w_{ij}$$

 $\hfill\Box$  Minicut: choosing a partition  $A_1,A_2,\ldots,A_K$  which minimizes

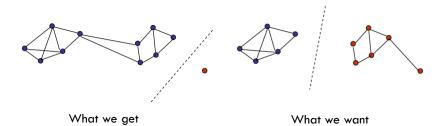
$$cut(A_1,\dots,A_k):=\frac{1}{2}\sum_{i=1}^k W(A_i,\overline{A_i})$$



### Graph Cut Shortcoming



Sensitive to outliers



### Solutions



"Size" of a subset  $A \subset V$  could be defined:

|A| := the number of vertices in A

 $vol(A) := \sum_{i \in A} d_i$ 

Solutions:

RatioCut(Hagen and Kahng, 1992)

$$RatioCut(A_1, ..., A_k) := \frac{1}{2} \sum_{i=1}^k \frac{W(A_i, \overline{A_i})}{|A_i|} = \sum_{i=1}^k \frac{cut(A_i, \overline{A_i})}{|A_i|}$$

Ncut(Shi and Malik, 2000)

$$Ncut(A_1, ..., A_k) := \frac{1}{2} \sum_{i=1}^k \frac{W(A_i, \overline{A_i})}{vol(A_i)} = \sum_{i=1}^k \frac{cut(A_i, \overline{A_i})}{vol(A_i)}$$

### Problem



#### NP hard!!!: can't be solved in polynomial time

$$\min RatioCut(A_1,\dots,A_k) := \frac{1}{2} \sum_{i=1}^k \frac{W(A_i,\overline{A_i})}{|A_i|} = \sum_{i=1}^k \frac{cut(A_i,\overline{A_i})}{|A_i|}$$

$$\min Ncut(A_1, \dots, A_k) := \frac{1}{2} \sum_{i=1}^k \frac{W(A_i, \overline{A_i})}{vol(A_i)} = \sum_{i=1}^k \frac{cut(A_i, \overline{A_i})}{vol(A_i)}$$

**Solution:** Approximation

### Solving Simple Case



- □ Approximation RatioCut for **k=2** 
  - Our goal is to solve the optimization problem:

$$\min_{A \subseteq V} RatioCut(A, \bar{A})$$

We can also write:

$$\begin{aligned} \min_{A \subset V} RatioCut(A, \bar{A}) &= \min_{A \subset V} \left( \frac{1}{|V|} |V| RatioCut(A, \bar{A}) \right) \\ \min_{A \subset V} \left( |V| RatioCut(A, \bar{A}) \right) \end{aligned}$$

### Solving Simple Case



We know:

$$RatioCut(A_1, \dots, A_k) := \frac{1}{2} \sum_{i=1}^k \frac{W(A_i, \overline{A_i})}{|A_i|} = \sum_{i=1}^k \frac{cut(A_i, \overline{A_i})}{|A_i|}$$

Hence, we can rewrite the objective function as:

$$|V|RatioCut(A, \bar{A}) = |V|(\frac{cut(A, \bar{A})}{|A|} + \frac{cut(A, \bar{A})}{|\bar{A}|})$$

$$= cut(A, \bar{A}) \left(\frac{|A| + |\bar{A}|}{|A|} + \frac{|A| + |\bar{A}|}{|\bar{A}|}\right)$$

$$= cut(A, \bar{A}) \left(\frac{|\bar{A}|}{|A|} + \frac{|A|}{|\bar{A}|} + 2\right)$$

Since we know:

$$cut(A_i,\overline{A_i}) = W(A,B) := \sum_{i \in A, j \in B} w_{ij}$$

$$=\frac{1}{2}\sum_{i\in A, j\in \bar{A}}w_{ij}(\sqrt{\frac{|\bar{A}|}{|A|}}+\sqrt{\frac{|A|}{|\bar{A}|}})^2+\sum_{i\in A, j\in A}w_{ij}(-\sqrt{\frac{|\bar{A}|}{|A|}}-\sqrt{\frac{|A|}{|\bar{A}|}})^2$$

### Solving Simple Case



$$=\frac{1}{2}\sum_{i\in A, j\in \bar{A}}w_{ij}(\sqrt{\frac{|\bar{A}|}{|A|}}+\sqrt{\frac{|A|}{|\bar{A}|}})^2+\sum_{i\in A, j\in A}w_{ij}(-\sqrt{\frac{|\bar{A}|}{|A|}}-\sqrt{\frac{|A|}{|\bar{A}|}})^2$$

If we define  $f_i$  as:

$$f_i = \begin{cases} \sqrt{|\bar{A}|/|A|}, & \text{if } v_i \in A \\ -\sqrt{|A|/|\bar{A}|}, & \text{if } v_i \in \bar{A} \end{cases}$$

$$\sum_{j=1}^n \sum_{i=1}^n w_{ij} f_j^2 = \sum_{j=1}^n d_j f_j^2$$
ion as:

We can rewrite the objective function as:

$$= \frac{1}{2} \sum_{i,j=1}^{n} w_{ij} (f_i - f_j)^2$$

$$= \frac{1}{2} (\sum_{i,j=1}^{n} w_{ij} f_i^2 - 2 \sum_{i,j=1}^{n} w_{ij} f_i f_j + \sum_{i,j=1}^{n} w_{ij} f_j^2)$$

$$= \frac{1}{2} (\sum_{i=1}^{n} d_i f_i^2 - 2 \sum_{i,j=1}^{n} f_i f_j w_{ij} + \sum_{j=1}^{n} d_j f_j^2)$$

$$= \sum_{i=1}^{n} d_i f_i^2 - \sum_{i,j=1}^{n} f_i f_j w_{ij}$$

$$= \int_{i=1}^{n} d_i f_i^2 - \sum_{i,j=1}^{n} f_i f_j w_{ij}$$

$$= \int_{i=1}^{n} f_i f_j w_{ij}$$

### Solving Simple Case



$$\min_{A \subset V} f'Lf \quad \text{if} \quad f_i = \begin{cases} \sqrt{|\bar{A}|/|A|}, & \text{if } v_i \in A \\ -\sqrt{|A|/|\bar{A}|}, & \text{if } v_i \in \bar{A} \end{cases}$$

We can say that f is orthogonal to the constant vector  $\mathbb{I}$ 

$$f\mathbb{I} = \ \textstyle \sum_{i=1}^n f_i = \ \textstyle \sum_{i \in A} \sqrt{\frac{|A|}{|A|}} - \sum_{i \in \bar{A}} \sqrt{\frac{|A|}{|\bar{A}|}} = |A| \sqrt{\frac{|\bar{A}|}{|A|}} - |\bar{A}| \sqrt{\frac{|A|}{|\bar{A}|}} = 0$$

f satisfies:

$$||f||^2 = \sum_{i=1}^n f_i^2 = |A| \sqrt{\frac{|\bar{A}|}{|A|}} + |\bar{A}| \sqrt{\frac{|A|}{|\bar{A}|}} = n$$

### New Objective



 $\min_{A \subset V} f' L f$  subject to  $f \perp \mathbb{I}$  ,  $f_i$  as defined in (1),  $\left| |f| \right| = \sqrt{n}$ 

$$f_i = \begin{cases} \sqrt{|A|/|A|}, & \text{if } v_i \in A \\ -\sqrt{|A|/|A|}, & \text{if } v_i \in \bar{A} \end{cases}$$

- $\ \square$  A discrete optimization problem in which f is only allowed to have two particular values.
- There are no efficient algorithms for solving optimization problems with discrete values

NP hard!!! => RELXATION

### Relaxed Objective



- Relaxation:
  - We remove constraints to create a relax version of the objective function which can be solved in polynomial time

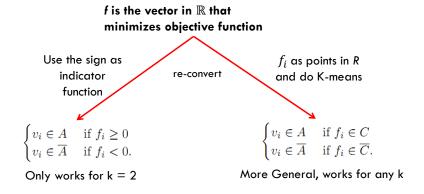
$$\min_{A \subseteq V} f' L f$$
 , subject to  $f \perp \mathbb{I} ||f|| = \sqrt{n}$ 

 $\hfill\Box$  The Relaxed version can be solved efficiency to get f that minimizes the objective function

### Creating the Partition



□ Approximation RatioCut for k=2



### Eigenvectors and Eigenvalues



An **eigenvector** of a square matrix A is a non-zero vector v that, when the matrix multiplies v, yields the same as when some scalar multiplies v, the scalar multiplier often being denoted by  $\lambda$ . That is:

$$Av = \lambda v => v^{-1}Av$$

 $\lambda$  is also called the eigenvalue of A

We also know that eigenvectors are orthogonal to each other

$$\min_{A \subseteq V} f' L f$$
 , subject to  $f \perp \mathbb{I} \mid ||f|| = \sqrt{n}$ 

f is the second eigenvector of L with eigenvalue equal to

$$\lambda = \text{RatioCut}(A, \bar{A})$$

 $\mathbb{I}$  is the first one with eigenvalue 0.

### Idea of Spectral Clustering



- □ We found K=2 eigenvectors of matrix L
  - L is Laplacian Matrix of Laplacian Graph
- □ We took data to a new dimensional system where data has linear structure
- □ We used K-Means to find best partition (clustering)

The fundamental idea of Spectral Clustering

### Agenda



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# STATE OF THE PARTY OF THE PARTY

### **Graph Laplacians**

Unnormalized Graph Laplacian

D= degree matrix where 
$$d_i = \sum_{j=1}^n w_{ij}$$
 
$$L = D - W \longrightarrow W = (w_{ij}) \ i, j = 1, ..., n \quad w_{ij} \geq 0$$

#### Properties of L

Proposition 1 The matrix L satisfies the following properties:

1. For every  $f \in \mathbb{R}^n$  we have

$$f'Lf = \frac{1}{2} \sum_{i,j=1}^{n} w_{ij} (f_i - f_j)^2$$



### **Graph Laplacians**

Unnormalized Graph Laplacian

$$L = D - W$$

$$f'Lf = \frac{1}{2} \sum_{i,j=1}^{n} w_{ij} (f_i - f_j)^2$$

- 1. L is symmetric and positive semi-definite.
  - L is symmetric since W and D are
  - $f'Lf \ge =$  therfore positive semidefinite
- 2. The smallest eigenvalue of L is 0, the corresponding eigenvector is the constant one vector  $\mathbb{1}$
- 3. L has n non-negative, real-valued eigenvalues  $0=\lambda_i\leq \lambda_2\leq \cdots \leq \lambda_n$ .

# STATE OF THE PARTY OF THE PARTY

### **Graph Laplacians**

Unnormalized Graph Laplacian

$$L = D - W$$

**Proposition 2 (Number of connected components and the spectrum of L)** Let G be an undirected graph with non-negative weights. The multiplicity k of the eigenvalue 0 of L equals the number of connected components  $A_1, \ldots, A_k$  in the graph. The eigenspace of eigenvalue 0 is spanned by the indicator vectors  $\mathbb{1}_{A_1}, \ldots, \mathbb{1}_{A_k}$  of those components.

Proof:

When k = 1

$$Lf = 0 f \quad (\lambda = 0) \quad \ => \quad \quad 0 = f' L f = \frac{1}{2} \sum_{i,j=1}^n w_{ij} \big( f_i - f_j \big)^2$$

Only happens if  $f_i = f_j$ 



### Graph Laplacians



Unnormalized Graph Laplacian

$$L = D - W$$

**Proposition 2** (Number of connected components and the spectrum of L) Let G be an undirected graph with non-negative weights. The multiplicity k of the eigenvalue 0 of L equals the number of connected components  $A_1, \ldots, A_k$  in the graph. The eigenspace of eigenvalue 0 is spanned by the indicator vectors  $1_{Proof}$ ;  $1_{A_k}$  of those components.

When k > 1, L can be written in a block form. the spectrum of L is given by the union of the spectra of  $L_i$ , and the corresponding eigenvectors of L are the eigenvectors of  $L_i$ , filled with 0 at the positions of the other blocks.

$$L = \begin{pmatrix} L_1 & & & \\ & L_2 & & \\ & & \ddots & \\ & & & L_k \end{pmatrix}$$

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### **Graph Laplacians**

Normalized Graph Laplacian

$$\begin{split} L_{sym} &:= D^{-\frac{1}{2}} L D^{-\frac{1}{2}} = I - D^{-\frac{1}{2}} W D^{-\frac{1}{2}} \\ L_{rw} &:= D^{-1} L = I - D^{-1} W \end{split}$$

We denote the first matrix by  $L_{sym}$  as it is a symmetric matrix, and the second one by  $L_{rw}$  as it is closely related to a random walk.



### **Graph Laplacians**

**Proposition 3 (Properties of**  $L_{sym}$  **and**  $L_{rw}$ **)** The normalized Laplacians statisfy the following properties:

1. For every  $f \in \mathbb{R}^n$  we have

$$f'L_{sym}f = \frac{1}{2} \sum_{i,j=1}^{n} w_{ij}^{2} \left( \frac{f_{i}}{\sqrt{d_{i}}} - \frac{f_{j}}{\sqrt{d_{j}}} \right)^{2} \qquad L_{sym} \coloneqq D^{-\frac{1}{2}}LD^{-\frac{1}{2}} = I - D^{-\frac{1}{2}}WD^{-\frac{1}{2}} \\ L_{rw} \coloneqq D^{-1}L = I - D^{-1}W$$

- 2.  $\lambda$  is an eigenvalue of  $L_{rw}$  with eigenvector u if and only if  $\lambda$  is an eigenvalue of  $L_{sym}$  with eigenvector  $w = D^{1/2}u$ .
- 3.  $\lambda$  is an eigenvalue of  $L_{rw}$  with eigenvector u if and only if  $\lambda$  and u solve the generalized eigen problem  $Lu=\lambda Du$ . ( $\Theta D^{-\frac{1}{2}}Lu=\lambda Du$ )

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### **Graph Laplacians**

$$f'L_{sym}f = \frac{1}{2}\sum_{i,j=1}^{n}w_{ij}^{2}\left(\frac{f_{i}}{\sqrt{d_{i}}} - \frac{f_{j}}{\sqrt{d_{j}}}\right)^{2} \qquad L_{sym} \coloneqq D^{-\frac{1}{2}}LD^{-\frac{1}{2}} = I - D^{-\frac{1}{2}}WD^{-\frac{1}{2}} = I - D^{-\frac{1}{2}}WD^{-\frac{1}{2}}$$

- 4. 0 is an eigenvalue of  $L_{rw}$  with the constant one vector  $\mathbb 1$  as eigenvector. 0 is an eigenvalue of  $L_{svm}$  with eigenvector  $D^{1/2}\mathbb 1$ . (former is obvious, latter derives from (2))
- 5.  $L_{sym}$  and  $L_{rw}$  are positive semi-definite and have n non-negative real-valued eigenvalues  $0=\lambda_i\leq \lambda_2\leq \cdots \leq \lambda_n.$

# Graph Laplacians



Normalized Graph Laplacian

$$L_{sym} := D^{-\frac{1}{2}} L D^{-\frac{1}{2}} = I - D^{-\frac{1}{2}} W D^{-\frac{1}{2}}$$
  
$$L_{rw} := D^{-1} L = I - D^{-1} W$$

Proposition 4 (Number of connected components and spectra of  $L_{sym}$  and  $L_{rw}$ ) Let G be an undirected graph with non-negative weights. Then the multiplicity k of the eigenvalue 0 of both  $L_{sym}$  and  $L_{rw}$  equals the number of connected components  $A_1, \ldots, A_k$  in the graph. For  $L_{rw}$  the eigenspace of 0 is spanned by the indicator vectors  $\mathbb{1}_{A_i}$  of those components. For  $L_{sym}$ , the eigenspace of 0 is spanned by the vectors  $D^{1/2}\mathbb{1}_{A_i}$ .

Proof. The proof is analogous to the one of Proposition 2, using Proposition 3.



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# Algorithm



- $\hfill\Box$  Main trick is to change the representation of the abstract data points  $x_i$  to points  $y_i \in \Re^k$ 
  - Unnormalized Spectral Clustering
  - Normalized Spectral Clustering 1
  - Normalized Spectral Clustering 2

### Algorithm



#### Unnormalized Graph Laplacian

$$L = D - W$$

#### Unnormalized spectral clustering

Input: Similarity matrix  $S \in \mathbb{R}^{n \times n}$ , number k of clusters to construct.

- ullet Construct a similarity graph by one of the ways described in Section 2. Let W be its weighted adjacency matrix.
- ullet Compute the unnormalized Laplacian L.
- Compute the first k eigenvectors  $u_1, \ldots, u_k$  of L.
- ullet Let  $U\in\mathbb{R}^{n imes k}$  be the matrix containing the vectors  $u_1,\dots,u_k$  as columns.
- For  $i=1,\ldots,n$ , let  $y_i\in\mathbb{R}^k$  be the vector corresponding to the i-th row of U.
- ullet Cluster the points  $(y_i)_{i=1,\dots,n}$  in  $\mathbb{R}^k$  with the k-means algorithm into clusters  $C_1,\dots,C_k$ .

Output: Clusters  $A_1, \ldots, A_k$  with  $A_i = \{j | y_j \in C_i\}$ .

## Algorithm



#### Normalized Graph Laplacian

$$L_{rw} := D^{-1}L = I - D^{-1}W$$

Normalized spectral clustering according to Shi and Malik (2000)

Input: Similarity matrix  $S \in \mathbb{R}^{n \times n}$ , number k of clusters to construct.

- ullet Construct a similarity graph by one of the ways described in Section 2. Let W be its weighted adjacency matrix.
- ullet Compute the unnormalized Laplacian L.
- Compute the first k generalized eigenvectors  $u_1, \ldots, u_k$  of the generalized eigenproblem  $\underline{Lu = \lambda Du}$ .
- ullet Let  $U\in\mathbb{R}^{n imes k}$  be the matrix containing the vectors  $u_1,\dots,u_k$  as columns.
- For  $i=1,\ldots,n$ , let  $y_i\in\mathbb{R}^k$  be the vector corresponding to the i-th row of U.
- Cluster the points  $(y_i)_{i=1,\dots,n}$  in  $\mathbb{R}^k$  with the k-means algorithm into clusters  $C_1,\dots,C_k$ .

Output: Clusters  $A_1, \ldots, A_k$  with  $A_i = \{j | y_j \in C_i\}$ .

Proposition 3 (Properties of  $L_{\text{sym}}$  and  $L_{\text{rw}}$ ) The normalized Laplacians satisfy the following properties:

 λ is an eigenvalue of L<sub>rw</sub> with eigenvector u if and only if λ and u solve the generalized eigenproblem Lu = λDu.

### Algorithm



#### Normalized Graph Laplacian

$$L_{sym} := D^{-\frac{1}{2}}LD^{-\frac{1}{2}} = I - D^{-\frac{1}{2}}WD^{-\frac{1}{2}}$$

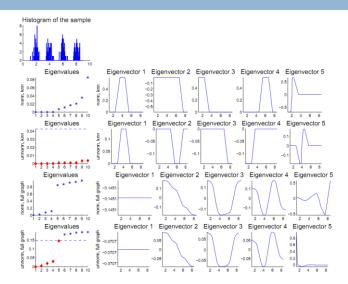
#### Normalized spectral clustering according to Ng, Jordan, and Weiss (2002)

Input: Similarity matrix  $S \in \mathbb{R}^{n \times n}$ , number k of clusters to construct.

- ullet Construct a similarity graph by one of the ways described in Section 2. Let Wbe its weighted adjacency matrix.
- ullet Compute the normalized Laplacian  $L_{
  m sym}$ .
- Compute the first k eigenvectors  $u_1, \ldots, u_k$  of  $L_{\text{sym}}$ .
- Let  $U \in \mathbb{R}^{n \times k}$  be the matrix containing the vectors  $u_1, \dots, u_k$  as columns.
- Form the matrix  $T \in \mathbb{R}^{n \times k}$  from U by normalizing the rows to norm 1,
- that is set  $t_{ij}=u_{ij}/(\sum_k u_{ik}^2)^{1/2}$ . For  $i=1,\ldots,n$ , let  $y_i\in\mathbb{R}^k$  be the vector corresponding to the i-th row of T.
- ullet Cluster the points  $(y_i)_{i=1,\dots,n}$  with the k-means algorithm into clusters  $C_1,\dots,C_k$  . Output: Clusters  $A_1, \ldots, A_k$  with  $A_i = \{j | y_j \in C_i\}$ .

# Algorithm







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- □ Spectral Clustering Algorithm
- □ Graph Cut Point of View
- □ Random Walk Point of View
- Practical Details

### Graph Cut



□ Approximation RatioCut for k=2

$$\min_{A\subset V} \mathrm{RatioCut}(A,\overline{A}).$$



$$f_i = \begin{cases} \sqrt{|\bar{A}|/|A|}, & \text{if } v_i \in A \\ -\sqrt{|A|/|\bar{A}|}, & \text{if } v_i \in \bar{A} \end{cases}$$

 $\min_{A \subset V} f' L f$  subject to  $f \perp \mathbb{I}$ ,  $f_i$  as defined in (1),  $||f|| = \sqrt{n}$ 



Relaxation !!!

 $\perp$ 

 $\min_{A\subset V} f'Lf$  , subject to  $f\perp_{\mathbb{T}} ||f||=\sqrt{n}$ 



Rayleigh-Ritz Theorem

f is the eigenvector corresponding to the second smallest eigenvalue of L (the smallest eigenvalue of L is 0 with eigenvector  $\mathbb{I}$ )

### Graph Cut Point of View



- □ It can be shown that:
  - Unnormalized Spectral Clustering algorithm is approximation of Ratio Cut Graph Partitioning
  - Normalized Spectral Clustering algorithms are approximations of Ncut Graph Partition

### **Abstract**



- Spectral Clustering gives an approximate solution to the solution of graph cut problem
- □ There is no guarantee on getting on the quality of the solution
- In general, efficient algorithms to approximate balanced graphs cut to a constant factor is also a NP-hard problem

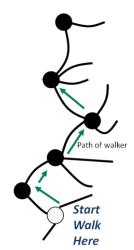


- □ Brief Clustering Review
- □ Similarity Graph
- □ Graph Laplacian
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- □ Random Walk Point of View
- □ Perturbation Theory Point of View
- Practical Details

### Random Walk



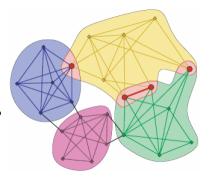
 A random walk on a graph is a stochastic process which randomly jumps from vertex to vertex.



### Random Walk in Similarity Graph



- Random walk stays long within the same cluster and seldom jumps between clusters.
- A balanced partition with a low cut will also have the property that the random walk does not have many opportunities to jump between clusters.



### Random walk



 $lue{}$  Transition probability  $p_{ij}$  of jumping from  $v_i$  to  $v_j$ 

$$p_{ij} = w_{ij}/d_i$$

lacksquare The **transition matrix**  $P=(p_{ij})$  i,j = 1,...,n of random walk is defined by

$$P = D^{-1}W$$

 $\square$  If the graph is connected and non-bipartite, the random walk always processes a unique stationary distribution  $\pi=(\pi_1,\ldots,\pi_n)'$ , where

$$\pi_i = d_i/vol(V)$$

$$d_i = \sum_{j=1}^n w_{ij} \qquad vol(V) := \sum_{i \in V} d_i$$

### Random walk



 $\ \square$  Relationship between  $L_{rw}$  and  $\emph{P}.$ 

$$L_{rw} = I - P$$
  $L_{rw} := D^{-1}L = I - D^{-1}W$ 

- $\ \square$  The largest eigenvectors of P and the smallest eigenvectors of  $L_{rw}$  can be used to describe cluster properties of the graph.

### Random walk



Random walks and Ncut

**Proposition 5** (Ncut **via transition probabilities**) Let G be connected and non bi-partite. Assume that we run the random walk  $(X_t)_{t\in N}$  starting with  $X_0$  in the stationary distribution  $\pi$ . For disjoint subsets  $A,B\subset V$ , denote by  $P(B|A):=P(X_1\in B|X_0\in A)$ . Then:

$$Ncut(A, \bar{A}) = P(\bar{A}|A) + P(A|\bar{A}).$$

It tells us that when minimizing Ncut, we actually look for a cut through the graph such that **A random walk seldom transitions from A to**  $\overline{A}$  and vice versa.

### Random walk



□ What is commute distance

The commute distance (resistance distance)  $c_{ij}$  between two vertices  $v_i$  and  $v_j$  is the expected time it takes the random walk to travel from vertex  $v_i$  to vertex  $v_i$  and back.

The commute distance between two vertices **decrease** if there are **many different short ways** to get from vertex  $v_i$  to vertex  $v_j$ .

instead of only shortest path

### Random walk



The commute distance between two vertices **decrease** if there are **many different short ways** to get from vertex  $v_i$  to vertex  $v_j$ .

· instead of only shortest path

Points which are **connected by a short path** in the graph and **lie in the same high-density region** of the graph are considered closer to each other than points which are connected by a short path but lie in different high-density regions of the graph.

**Well-suited for Clustering** 



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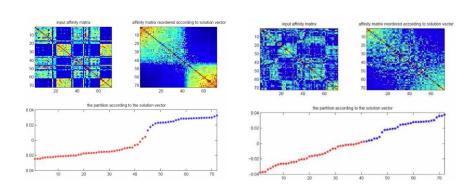
### Similarity function



- Constructing the similarity graph
  - 1. Similarity Function Itself
    - Make sure that points which are considered to be "very similar" by the similarity function are also closely related in the application the data comes from.
    - Global-range behavior of similarity function is not important
    - Domain specific

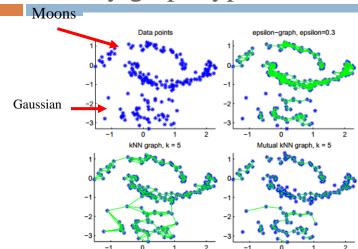
## Similarity function





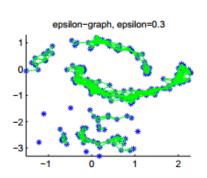
# Similarity graph type





## Similarity graph type

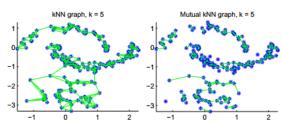




- Difficult to choose ε
  - Points on the moon are strongly connected
  - Points on Gaussian are barely connected
- Problem of "data on different scales"

### Similarity graph type





- K-nearest neighbor graph
  - Can handle "data on different scales"
  - Can break into disconnected components if there are densely connected regions
- Mutual k-nearest neighbor
  - Can act in presence of data at different densities but doesn't mix them together
  - Well suited

### Parameters of similarity graphs



- Significance of having connected similarity graph
  - In general, if the similarity graph contains more connected components than the number of clusters we ask the algorithm to detect, then spectral clustering will trivially return connected components as clusters
  - Unless one is perfectly sure that those connected components are the correct clusters, one should make sure that the similarity graph is connected, or only consists of "few" connected components and very few or no isolated vertices.
  - □ To prevent **outliers** to affect result
  - Set parameters small enough to address local structure of data while insuring connectivity

### Parameters of similarity graphs



- 1. Parameters of Similarity Graph(k or  $\varepsilon$ )
  - 1. KNN:
    - k in order of log(n);
  - . Mutual KNN:
    - 1. Has much fewer edges than knn
      - 1. k significantly larger than standard KNN;
    - Maintain the aforementioned advantage of not connected regions of different density
      - 1. Keep it small enough
    - 3. Unfortunately, no general heuristics

### Parameters of similarity graphs



#### 1. Parameters of Similarity Graph(k or $\varepsilon$ )

#### £-neighborhood graph:

- longest edge of MST;
- 2. Order of  $(\log(n)/n)^d$
- 3. A significantly large  $\varepsilon$  will be selected when (too large):
  - 1. Data contains outlies
  - 2. Several tight components very far from each other

#### fully connected graph:

- 1. Select  $\sigma$  in a way that similarity is "not too small not too large"
- 2. Select  $\sigma$  in order of the mean distance of a point to its k-th nearest neighbor. Or choose  $k=\varepsilon$ .

### Computing eigenvectors



#### Computing Eigenvectors

- 1. How to compute the first eigenvectors efficiently for large L
- 2. For sparse graphs from knn or  $\varepsilon$ 
  - Efficient methods exist to compute the first eigenvectors of sparse matrices. The speed of convergence of those algorithms depends on the size of the eigengap (also called spectral gap)
  - 2. Eigengap:  $\gamma_k = |\lambda_k \lambda_{k+1}|$
  - 3. The large the Eigengap the faster the algorithms

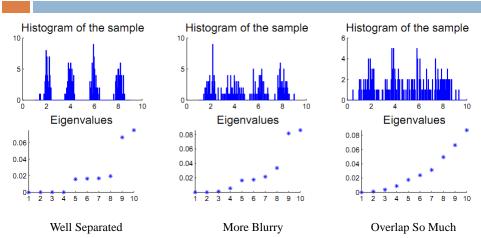
### Determining k



- Well-justified criteria in model-based clustering settings
  - Based on log-likelihood of the data
- Methods specific to Spectral Clustering
  - Eigengap Heuristic
    - Used for all 3 graph laplacian
    - Choose K that  $\lambda_1 ... \lambda_k$  are very small, but,  $\lambda_{k+1}$  is relatively large
    - based on **perturbation theory**, where we observe that in the ideal case of k completely disconnected clusters, the eigenvalue 0 has multiplicity k, and then there is a gap to the (k+1)th eigenvalue  $\lambda k+1 > 0$

#### **DETERMINING K**





Eigengap Heuristic usually works well if the data contains very well pronounced clusters, but in ambiguous cases it also returns ambiguous results.

#### WHICH GRAPH LAPLACIAN



#### The k-means step

It is not necessary. People also use other techniques

Which graph Laplacian should be used?

When does it become important?

If the distribution of degrees in graph is very regular and most vertices have approximately the same degree, then all the Laplacians are very similar to each other, and will work equally well for clustering. However, if the degrees in the graph are very broadly distributed, then the Laplacians differ considerably.

#### WHICH GRAPH LAPLACIAN



Which graph Laplacian should be used?

Why normalized is better than unnormalized spectral clustering?

#### Objective1:

1. We want to find a partition such that points in different clusters are dissimilar to each other, that is we want to minimize the between-cluster similarity. In the graph setting, this means to minimize  $cut(A, \bar{A})$ .

#### Both RatioCut and Ncut directly implement

#### Objective2:

2. We want to find a partition such that points in the same cluster are similar to each other, that is we want to maximize the within-cluster similarities W(A,A), and  $W(\bar{A},\bar{A})$ .

#### Only Ncut implements

Normalized spectral clustering implements both clustering objectives mentioned above, while unnormalized spectral clustering only implements the first obejctive.



#### WHICH GRAPH LAPLACIAN

#### □ Ncut:

- $W(A,A) = W(A,V) W(A,\bar{A}) = vol(A) cut(A,\bar{A})$
- lacktriangledown the within-cluster similarity is maximized if  $cut(A, \bar{A})$  is small and if vol(A) is large.
- As this is exactly what we achieve by minimizing Ncut
- Ncut criterion implements the second objective

#### □ RatioCut:

- Objective is to maximize |A| and  $|\bar{A}|$  instead of vol(A) and vol( $\bar{A}$ )
- $\square$  |A| and | $\bar{A}$ | are not necessary related to the within-cluster similarity since within-cluster similarity depends on the edges and not the number of nodes
- A graph with a lot of vertices with low weight edges to each other



#### WHICH GRAPH LAPLACIAN

- Second perspective: Consistency Issue
  - $\mathbf{x}_1 \dots \mathbf{x}_k$  have been sampled i.i.d according to some probability distribution P on some underlying data space X
  - The most fundamental question is then the question of consistency: if we draw more and more data points, do the clustering results of spectral clustering converge to a useful partition of the underlying space X ?

#### WHICH GRAPH LAPLACIAN

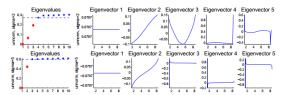


#### □ Normalized Spectral clustering:

All consistency statements about normalized spectral clustering hold, for both Lsym and Lrw, under very mild conditions which are usually satisfied in real world applications

#### Unnromalized:

All consistency statements about normalized spectral clustering hold, for both Lsym and Lrw, under very mild conditions which are usually satisfied in real world applications



#### WHICH GRAPH LAPLACIAN



Which graph Laplacian should be used?

Why the eigenvectors of  $L_{rw}$  are better than those of  $L_{sym}$ ?

- 1. Eigenvectors of  $L_{rw}$  are cluster indicator vectors  $\mathbb{I}_{A_i}$ , while the eigenvectors of  $L_{sym}$  are additionally multiplied with  $D^{1/2}$ , which might lead to undesired artifacts.
- 2. Using  $L_{\mathit{Sym}}$  also does not have any computational advantages.

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# **Thanks**