# Spectral Clustering 

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## Agenda

- Brief Clustering Review
- Similarity Graph
- Graph Laplacian
- Spectral Clustering Algorithm
- Graph Cut Point of View
- Random Walk Point of View
- Perturbation Theory Point of View
- Practical Details
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## CLUSTERING REVIEW

## Clustering

Groups together "similar" instances in the data sample

## Basic clustering problem:

- distribute data into $k$ different groups such that data points similar to each other are in the same group
- Similarity between data points is defined in terms of some distance metric (can be chosen)

Clustering is useful for:

- Similarity/Dissimilarity analysis

Analyze what data points in the sample are close to each other

- Dimensionality reduction

High dimensional data replaced with a group (cluster) label

## K-MEANS CLUSTERING

- Description

Given a set of observations $\left(\mathbf{x}_{1}, \mathbf{x}_{2}, \ldots, \mathbf{x}_{n}\right)$, where each observation is a $d$-dimensional real vector, $k$-means clustering aims to partition the $n$ observations into $k$ sets $(k \leq n) \mathbf{S}=\left\{S_{1}, S_{2}, \ldots, S_{k}\right\}$ so as to minimize the within-cluster sum of squares (WCSS):

$$
\underset{S}{\arg \min } \sum_{i=1}^{k} \sum_{x_{j} \in S_{i}}\left\|x_{j}-u_{i}\right\|^{2}
$$

where $\boldsymbol{\mu}_{i}$ is the mean of points in $S_{i}$.

- Standard Algorithm


1) k initial "means" (in this case $\mathrm{k}=3$ ) are randomly selected from the data set.

2) k clusters are created by associating every observation with the nearest mean.

3) The centroid of each of the $k$ clusters becomes the new means.

4) Steps 2 and 3 are repeated until convergence has been reached.

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## GENERAL

First - graph representation of data (largely, application dependent)


O


Then - graph partitioning


Disconnected graph components

Groups of points(Weakly connections in between components Strongly connections within components)

## GRAPH NOTATION

$G=(V, E)$ :

- Vertex set $V=\left\{v_{1}, \ldots, v_{n}\right\}$
- Weighted adjacency matrix $W=\left(w_{i j}\right) i, j=1, \ldots, n \quad w_{i j} \geq 0$

- Degree $d_{i}=\sum_{j=1}^{n} w_{i j}$

- Degree matrix Diagonal matrix with the degrees $d_{1}, \ldots, d_{n}$ on the diagonal.


## GRAPH NOTATION

$G=(V, E)$ :

- Indicator Vector $\mathbb{1}_{A}=\left(f_{1}, \ldots, f_{n}\right)^{\prime} \in \mathbb{R}^{n} \quad f_{i} \in\{0,1\}$
- "Size" of a subset $A \subset V$

$$
|A|:=\text { the number of vertices in } A
$$

$$
\operatorname{vol}(A):=\sum_{i \in A} d_{i}
$$



- Connected A subset $A$ of a graph is connected if any two vertices in $A$ can be joined by a path such that all intermediate points also lie in $A$.
- Connected Component it is connected and if there are no connections between vertices in $A$ and $\bar{A}$. The nonempty sets $A_{1}, \ldots, A_{k}$ form a partition of the graph if $A_{i} \cap A_{j}=\emptyset$ and $A_{1} \cup \cdots \cup A_{k}=V$.


## SIMILARITY GRAPH

- $\varepsilon$-neighborhood graph

Connect all points whose pairwise distances are smaller than $\varepsilon$

- $k$-nearest neighbor graph

Connect vertex $v_{i}$ with vertex $v_{j}$ if $v_{j}$ is among the $k$-nearest neighbors of $v_{i}$.

- fully connected graph

Connect all points with positive similarity with each other

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## GRAPH LAPLACIANS

- Unnormalized Graph Laplacian

$$
d_{i}=\sum_{j=1}^{n} w_{i j}
$$

$$
L=D-W
$$

Proposition 1 (Properties of L) The matrix L satisfies the following properties:

1. For every $f \in \mathbb{R}^{n}$ we have

$$
\begin{gathered}
f^{\prime} L f=\frac{1}{2} \sum_{i, j=1}^{n} w_{i j}^{2}\left(f_{i}-f_{j}\right)^{2} \\
f^{\prime} L f=f^{\prime} D f-f^{\prime} W f=\sum_{i=1}^{n} d_{i} f_{i}^{2}-\sum_{i, j=1}^{n} f_{i} f_{j} w_{i j} \\
=\frac{1}{2}\left(\sum_{i=1}^{n} d_{i} f_{i}^{2}-2 \sum_{i, j=1}^{n} f_{i} f_{j} w_{i j}+\sum_{j=1}^{n} d_{j} f_{j}^{2}\right)=\frac{1}{2} \sum_{i, j=1}^{n} w_{i j}\left(f_{i}-f_{j}\right)^{2}
\end{gathered}
$$

## GRAPH LAPLACIANS

- Unnormalized Graph Laplacian

$$
L=D-W
$$

Proposition 1 (Properties of L) The matrix L satisfies the following properties:

1. For every $f \in \mathbb{R}^{n}$ we have

$$
f^{\prime} L f=\frac{1}{2} \sum_{i, j=1}^{n} w_{i j}^{2}\left(f_{i}-f_{j}\right)^{2}
$$

2. $L$ is symmetric and positive semi-definite.
3. The smallest eigenvalue of $L$ is 0 , the corresponding eigenvector is the constant one vector $\mathbb{1}$
4. $\quad L$ has $n$ non-negative, real-valued eigenvalues $0=\lambda_{i} \leq \lambda_{2} \leq \cdots \leq \lambda_{n}$.

## GRAPH LAPLACIANS

- Unnormalized Graph Laplacian

$$
L=D-W
$$

Proposition 2 (Number of connected components and the spectrum of $L$ ) Let G be an undirected graph with non-negative weights. The multiplicity k of the eigenvalue 0 of $L$ equals the number of connected components $A_{1}, \ldots, A_{k}$ in the graph. The eigenspace of eigenvalue 0 is spanned by the indicator vectors $\mathbb{1}_{A_{1}}, \ldots, \mathbb{1}_{A_{k}}$ of those components.

## Proof:

When $k=1$, a graph consisting of only one connected component we thus only have the constant one vector $\mathbb{1}$ as eigenvector with eigenvalue 0 , which obviously is the indicator vector of the connected component.

When $k>1, L$ can be written in a block form. the spectrum of $L$ is given by the union of the spectra of $L_{i}$, and the corresponding eigenvectors of $L$ are the eigenvectors of $L_{i}$, filled with 0 at the positions of the other blocks.

$$
L=\left(\begin{array}{cccc}
L_{1} & & & \\
& L_{2} & & \\
& & \ddots & \\
& & & L_{k}
\end{array}\right)
$$

## GRAPH LAPLACIANS

- Normalized Graph Laplacian

$$
\begin{aligned}
& L_{\text {sym }}:=D^{-\frac{1}{2}} L D^{-\frac{1}{2}}=I-D^{-\frac{1}{2}} W D^{-\frac{1}{2}} \\
& L_{r w}:=D^{-1} L=I-D^{-1} W
\end{aligned}
$$

We denote the first matrix by $L_{\text {sym }}$ as it is a symmetric matrix, and the second one by $L_{r w}$ as it is closely related to a random walk.

## GRAPH LAPLACIANS

- Normalized Graph Laplacian

$$
\begin{aligned}
& L_{s y m}:=D^{-\frac{1}{2}} L D^{-\frac{1}{2}}=I-D^{-\frac{1}{2}} W D^{-\frac{1}{2}} \\
& L_{r w}:=D^{-1} L=I-D^{-1} W
\end{aligned}
$$

Proposition 3 (Properties of $\boldsymbol{L}_{\boldsymbol{s y m}}$ and $\boldsymbol{L}_{\boldsymbol{r} \boldsymbol{w}}$ ) The normalized Laplacians statisfy the following properties:

1. For every $f \in \mathbb{R}^{n}$ we have $f^{\prime} L_{s y m} f=\frac{1}{2} \sum_{i, j=1}^{n} w_{i j}^{2}\left(\frac{f_{i}}{\sqrt{d_{i}}}-\frac{f_{j}}{\sqrt{d_{j}}}\right)^{2}$
2. $\lambda$ is an eigenvalue of $L_{r w}$ with eigenvector $u$ if and only if $\lambda$ is an eigenvalue of $L_{\text {sym }}$ with eigenvector $w=D^{1 / 2} u$.
3. $\lambda$ is an eigenvalue of $L_{r w}$ with eigenvector u if and only if $\lambda$ and u solve the generalized eigen problem $L u=\lambda D u$.
4. 0 is an eigenvalue of $L_{r w}$ with the constant one vector $\mathbb{1}$ as eigenvector. 0 is an eigenvalue of $L_{\text {sym }}$ with eigenvector $D^{1 / 2} \mathbb{1}$.
5. $L_{s y m}$ and $L_{r w}$ are positive semi-definite and have $n$ non-negative real-valued eigenvalues $0=\lambda_{i} \leq \lambda_{2} \leq \cdots \leq \lambda_{n}$.

## GRAPH LAPLACIANS

- Normalized Graph Laplacian

$$
\begin{aligned}
& L_{s y m}:=D^{-\frac{1}{2}} L D^{-\frac{1}{2}}=I-D^{-\frac{1}{2}} W D^{-\frac{1}{2}} \\
& L_{r w}:=D^{-1} L=I-D^{-1} W
\end{aligned}
$$

Proposition 4 (Number of connected components and spectra of $\boldsymbol{L}_{\boldsymbol{s y m}}$ and $\boldsymbol{L}_{\boldsymbol{r w}}$ ) Let $G$ be an undirected graph with non-negative weights. Then the multiplicity $k$ of the eigenvalue 0 of both $L_{s y m}$ and $L_{r w}$ equals the number of connected components $A_{1}, \ldots, A_{k}$ in the graph. For $L_{r w}$ the eigenspace of 0 is spanned by the indicator vectors $\mathbb{1}_{A_{i}}$ of those components. For $L_{s y m}$, the eigenspace of 0 is spanned by the vectors $D^{1 / 2} \mathbb{1}_{A_{i}}$.

Proof. The proof is analogous to the one of Proposition 2, using Proposition 3.

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## ALGORIGHM

Main trick is to change the representation of the abstract data points $x_{i}$ to points $y_{i} \in \mathfrak{R}^{k}$

1. Unnormalized Spectral Clustering
2. Normalized Spectral Clustering 1
3. Normalized Spectral Clustering 2

## ALGORIGHM

- Unnormalized Graph Laplacian

$$
L=D-W
$$

Unnormalized spectral clustering
Input: Similarity matrix $S \in \mathbb{R}^{n \times n}$, number $k$ of clusters to construct.

- Construct a similarity graph by one of the ways described in Section 2. Let $W$ be its weighted adjacency matrix.
- Compute the unnormalized Laplacian $L$.
- Compute the first $k$ eigenvectors $u_{1}, \ldots, u_{k}$ of $L$.
- Let $U \in \mathbb{R}^{n \times k}$ be the matrix containing the vectors $u_{1}, \ldots, u_{k}$ as columns.
- For $i=1, \ldots, n$, let $y_{i} \in \mathbb{R}^{k}$ be the vector corresponding to the $i$-th row of $U$.
- Cluster the points $\left(y_{i}\right)_{i=1, \ldots, n}$ in $\mathbb{R}^{k}$ with the $k$-means algorithm into clusters $C_{1}, \ldots, C_{k}$.
Output: Clusters $A_{1}, \ldots, A_{k}$ with $A_{i}=\left\{j \mid y_{j} \in C_{i}\right\}$.


## ALGORIGHM

## - Normalized Graph Laplacian

$$
L_{r w}:=D^{-1} L=I-D^{-1} W
$$

Normalized spectral clustering according to Shi and Malik (2000)
Input: Similarity matrix $S \in \mathbb{R}^{n \times n}$, number $k$ of clusters to construct.

- Construct a similarity graph by one of the ways described in Section 2. Let $W$ be its weighted adjacency matrix.
- Compute the unnormalized Laplacian $L$.
- Compute the first $k$ generalized eigenvectors $u_{1}, \ldots, u_{k}$ of the generalized eigenproblem $L u=\lambda D u$.
- Let $U \in \mathbb{R}^{n \times k}$ be the matrix containing the vectors $u_{1}, \ldots, u_{k}$ as columns.
- For $i=, \ldots, n$, let $y_{i} \in \mathbb{R}^{k}$ be the vector corresponding to the $i$-th row of $U$.
- Cluster the points $\left(y_{i}\right)_{i=1, \ldots, n}$ in $\mathbb{R}^{k}$ with the $k$-means algorithm into clusters $C_{1}, \ldots, C_{k}$.
Output: Clusters $A_{1}, \ldots, A_{k}$ with $A_{i}=\left\{j \mid y_{j} \in C_{i}\right\}$.

Proposition 3 (Properties of $L_{\mathrm{sym}}$ and $L_{\mathrm{rw}}$ ) The normalized Laplacians satisfy the following properties:
3. $\lambda$ is an eigenvalue of $L_{r w}$ with eigenvector $u$ if and only if $\lambda$ and $u$ solve the generalized eigenproblem $L u=\lambda D u$.

## ALGORIGHM

- Normalized Graph Laplacian

$$
L_{\text {sym }}:=D^{-\frac{1}{2}} L D^{-\frac{1}{2}}=I-D^{-\frac{1}{2}} W D^{-\frac{1}{2}}
$$

Normalized spectral clustering according to Ng, Jordan, and Weiss (2002)
Input: Similarity matrix $S \in \mathbb{R}^{n \times n}$, number $k$ of clusters to construct.

- Construct a similarity graph by one of the ways described in Section 2. Let $W$ be its weighted adjacency matrix.
- Compute the normalized Laplacian $L_{\text {sym }}$.
- Compute the first $k$ eigenvectors $u_{1}, \ldots, u_{k}$ of $L_{\mathrm{sym}}$.
- Let $U \in \mathbb{R}^{n \times k}$ be the matrix containing the vectors $u_{1}, \ldots, u_{k}$ as columns.
- Form the matrix $T \in \mathbb{R}^{n \times k}$ from $U$ by normalizing the rows to norm 1 , that is set $t_{i j}=u_{i j} /\left(\sum_{k} u_{i k}^{2}\right)^{1 / 2}$.
- For $i=1, \ldots, n$, let $y_{i} \in \mathbb{R}^{k}$ be the vector corresponding to the $i$-th row of $T$.
- Cluster the points $\left(y_{i}\right)_{i=1, \ldots, n}$ with the $k$-means algorithm into clusters $C_{1}, \ldots, C_{k}$. Output: Clusters $A_{1}, \ldots, A_{k}$ with $A_{i}=\left\{j \mid y_{j} \in C_{i}\right\}$.


## ALGORIGHM

Histogram of the sample


Eigenvector 1 Eigenvector 2 Eigenvector 3 Eigenvector 4
Eigenvector 5







Eigenvector 1 Eigenvector 2 Eigenvector 3 Eigenvector 4 Eigenvector 5


Eigenvalues







Eigenvalues







## ALGORIGHM

## On Spectral Clustering: <br> Analysis and an algorithm

| Andrew Y. Ng | Michael I. Jordan | Yair Weiss |
| :---: | :---: | :---: |
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## GRAPH CUT

First - graph representation of data (largely, application dependent)


Disconnected graph components

Groups of points(Weakly connections in between components Strongly connections within components)

## GRAPH CUT

## $G=(V, E)$ :

- For two not necessarily disjoint set $A, B \subset V$, we define

$$
W(A, B):=\sum_{i \in A, j \in B} w_{i j}
$$

- Minicut: choosing a partition $A_{1}, A_{2}, \ldots, A_{K}$ which minimizes

$$
\operatorname{cut}\left(A_{1}, \ldots, A_{k}\right):=\frac{1}{2} \sum_{i=1}^{k} W\left(A_{i}, \overline{A_{i}}\right)
$$

Cut between 2 sets

$$
\operatorname{cut}\left(A_{1}, A_{2}\right)=\sum_{n \in A_{1}} \sum_{m \in A_{2}} w_{n m}
$$



## GRAPH CUT

## Problems!!!

- Sensitive to outliers


What we get
What we want

## GRAPH CUT

## Solutions

$|A|:=$ the number of vertices in $A$

$$
\operatorname{vol}(A):=\sum_{i \in A} d_{i}
$$

- RatioCut(Hagen and Kahng, 1992)

$$
\operatorname{RatioCut}\left(A_{1}, \ldots, A_{k}\right):=\frac{1}{2} \sum_{i=1}^{k} \frac{W\left(A_{i}, \overline{A_{i}}\right)}{\left|A_{i}\right|}=\sum_{i=1}^{k} \frac{\operatorname{cut}\left(A_{i}, \overline{A_{i}}\right)}{\left|A_{i}\right|}
$$

- Ncut(Shi and Malik, 2000)

$$
\operatorname{Ncut}\left(A_{1}, \ldots, A_{k}\right):=\frac{1}{2} \sum_{i=1}^{k} \frac{W\left(A_{i}, \overline{A_{i}}\right)}{\operatorname{vol}\left(A_{i}\right)}=\sum_{i=1}^{k} \frac{\operatorname{cut}\left(A_{i}, \overline{A_{i}}\right)}{\operatorname{vol}\left(A_{i}\right)}
$$

## GRAPH CUT

Problem!!!
$\operatorname{RatioCut}\left(A_{1}, \ldots, A_{k}\right):=\frac{1}{2} \sum_{i=1}^{k} \frac{W\left(A_{i}, \overline{A_{i}}\right)}{\left|A_{i}\right|}=\sum_{i=1}^{k} \frac{\operatorname{cut}\left(A_{i}, \overline{A_{i}}\right)}{\left|A_{i}\right|}$

- NP hard

$$
\operatorname{Ncut}\left(A_{1}, \ldots, A_{k}\right):=\frac{1}{2} \sum_{i=1}^{k} \frac{W\left(A_{i}, \overline{A_{i}}\right)}{\operatorname{vol}\left(A_{i}\right)}=\sum_{i=1}^{k} \frac{\operatorname{cut}\left(A_{i}, \overline{A_{i}}\right)}{\operatorname{vol}\left(A_{i}\right)}
$$

## Solution!!!

- Approximation



## GRAPH CUT

- Approximation RatioCut for $\mathrm{k}=2$

Our goal is to solve the optimization problem:

$$
\min _{A \subset V} \operatorname{RatioCut}(A, \bar{A})
$$

Rewrite the problem in a more convenient form:
Given a subset $A \subset V$, we define the vector $f=\left(f_{1}, \ldots, f_{n}\right)^{\prime} \in \mathbb{R}^{n}$ with entries

$$
f_{i}= \begin{cases}\sqrt{|\bar{A}| /|A|}, & \text { if } v_{i} \in A \\ -\sqrt{|\bar{A}| /|A|}, & \text { if } v_{i} \in \bar{A}\end{cases}
$$

## Magic happens!!!

## GRAPH CUT

- Approximation RatioCut for $\mathrm{k}=2$

$$
f_{i}= \begin{cases}\sqrt{|\bar{A}| /|A|,} & \text { if } v_{i} \in A \\ -\sqrt{|\bar{A}| /|A|}, & \text { if } v_{i} \in \bar{A}\end{cases}
$$

$$
\begin{aligned}
& \quad \operatorname{RatioCut}\left(A_{1}, \ldots, A_{k}\right):=\frac{1}{2} \sum_{i=1}^{k} \frac{W\left(A_{i}, \overline{A_{i}}\right)}{\left|A_{i}\right|}=\sum_{i=1}^{k} \frac{\operatorname{cut}\left(A_{i}, \overline{A_{i}}\right)}{\left|A_{i}\right|} \\
f^{\prime} L f & =\frac{1}{2} \sum_{i, j=1}^{n} w_{i j}\left(f_{i}-f_{j}\right)^{2} \quad \operatorname{cut}\left(A_{1}, \ldots, A_{k}\right):=\frac{1}{2} \sum_{i=1}^{k} W\left(A_{i}, \overline{A_{i}}\right) \quad W(A, B):=\sum_{i \in A, \in \in B} w_{i j} \\
& =\frac{1}{2} \sum_{i \in A, j \in \bar{A}} w_{i j}\left(\sqrt{\frac{|\bar{A}|}{|A|}}+\sqrt{\frac{|A|}{|\bar{A}|}}\right)^{2}+\frac{1}{2} \sum_{i \in \bar{A}, j \in A} w_{i j}\left(-\sqrt{\frac{|\bar{A}|}{|A|}}-\sqrt{\frac{|A|}{|\bar{A}|}}\right)^{2} \\
& =\operatorname{cut}(A, \bar{A})\left(\frac{|\bar{A}|}{|A|}+\frac{|A|}{|\bar{A}|}+2\right) \\
& =\operatorname{cut}(A, \bar{A})\left(\frac{|A|+|\bar{A}|}{|A|}+\frac{|A|+|\bar{A}|}{|\bar{A}|}\right) \\
& =|V| \cdot \operatorname{RatioCut}(A, \bar{A}) .
\end{aligned}
$$

## GRAPH CUT

- Approximation RatioCut for $\mathrm{k}=2$

Additionally, we have

$$
\sum_{i=1}^{n} f_{i}=\sum_{i \in A} \sqrt{\frac{|\bar{A}|}{|A|}}-\sum_{i \in \bar{A}} \sqrt{\frac{|A|}{|\bar{A}|}}=|A| \sqrt{\frac{|\bar{A}|}{|A|}}-|\bar{A}| \sqrt{\frac{|A|}{|\bar{A}|}}=0
$$

The vector $f$ as defined before is orthogonal to the constant one vector $\mathbb{l}$.
$f$ satisfies

$$
\|f\|^{2}=\sum_{i=1}^{n} f_{i}^{2}=|A| \frac{|\bar{A}|}{|A|}+|\bar{A}| \frac{|A|}{|\bar{A}|}=|\bar{A}|+|A|=n
$$

## GRAPH CUT

- Approximation RatioCut for $\mathrm{k}=2$

$$
f_{i}= \begin{cases}\sqrt{|\bar{A}| /|A|} & \text { if } v_{i} \in A \\ -\sqrt{|A| /|\bar{A}|} & \text { if } v_{i} \in \bar{A}\end{cases}
$$

## $\min _{A \subset V} \operatorname{RatioCut}(A, \bar{A})$.

$$
\begin{aligned}
& f^{\prime} L f=\frac{1}{2} \sum_{i, j=1}^{n} w_{i j}\left(f_{i}-f_{j}\right)^{2} \\
& =|V| \cdot \operatorname{RatioCut}(A, \bar{A}) .
\end{aligned}
$$

$\min _{A \subset V} f^{\prime} L f$ subject to $f \perp \mathbb{1}\|f\|=\sqrt{n}$.
Relaxation !!!

$$
\min _{f \in \mathbb{R}^{n}} f^{\prime} L f \text { subject to } f \perp \mathbb{1},\|f\|=\sqrt{n} .
$$

Rayleigh-Ritz Theorem
$f$ is the eigenvector corresponding to the second smallest eigenvalue of $L$ (the smallest eigenvalue of $L$ is 0 with eigenvector $\mathbb{1}$ )

## GRAPH CUT

- Approximation RatioCut for $\mathrm{k}=2$


Only works for $\mathrm{k}=2$
More General, works for any k

## GRAPH CUT

- Approximation RatioCut for arbitrary k

Given a partition of $V$ into $k$ sets $A_{1}, A_{2}, \ldots, A_{k}$, we define $k$ indicator vectors $h_{j}=\left(h_{1, j}, \ldots, h_{n, j}\right)^{\prime}$ by

$$
h_{i, j}=\left\{\begin{array}{ll}
\frac{1}{\sqrt{\left|A_{j}\right|}}, & \text { if } v_{i} \in A_{j} \\
0, & \text { otherwise }
\end{array} \quad(i=1, \ldots, n ; j=1, \ldots, k)\right.
$$

$$
h_{i}^{\prime} L h_{i}=\frac{\operatorname{cut}\left(A_{i}, \overline{A_{i}}\right)}{\left|A_{i}\right|}
$$

$H \in \mathbb{R}^{n \times k}$, containing those $k$ Indicator vectors as columns.
Columns in $H$ are orthonormal to each other, that is $H^{\prime} H=I$

## GRAPH CUT

- Approximation RatioCut for arbitrary k

$$
h_{i, j}= \begin{cases}\frac{1}{\sqrt{\left|A_{j}\right|}}, & \text { if } v_{i} \in A_{j} \\ 0, & \text { otherwise }\end{cases}
$$

Problem reformulation:


Optimal $H$ is the first $k$ eigenvectors of $L$ as columns.

## GRAPH CUT

- Approximation Ncut for $\mathrm{k}=2$

$$
\operatorname{Ncut}\left(A_{1}, \ldots, A_{k}\right):=\frac{1}{2} \sum_{i=1}^{k} \frac{W\left(A_{i}, \overline{A_{i}}\right)}{\operatorname{vol}\left(A_{i}\right)}=\sum_{i=1}^{k} \frac{\operatorname{cut}\left(A_{i}, \overline{A_{i}}\right)}{\operatorname{vol}\left(A_{i}\right)}
$$

Our goal is to solve the optimization problem:

$$
\min _{A \subset V} \operatorname{Ncut}(A, \bar{A})
$$

Rewrite the problem in a more convenient form:
Given a subset $A \subset V$, we define the vector $f=\left(f_{1}, \ldots, f_{n}\right)^{\prime} \in \mathbb{R}^{n}$ with entries

$$
f_{i}= \begin{cases}\sqrt{\frac{\operatorname{vol}(\bar{A})}{\operatorname{vol}(A)}} & \text { if } v_{i} \in A \\ -\sqrt{\frac{\operatorname{vol}(A)}{\operatorname{vol}(\bar{A})}} & \text { if } v_{i} \in \bar{A}\end{cases}
$$

Similar to above one can check that:

$$
(D f)^{\prime} \mathbb{1}=0, f^{\prime} D f=\operatorname{vol}(V), \text { and } f^{\prime} L f=\operatorname{vol}(V) N c u t(A, \bar{A})
$$

## GRAPH CUT

- Approximation Ncut for k=2

$$
f_{i}= \begin{cases}\sqrt{\frac{\operatorname{vol}(\bar{A})}{\operatorname{vol} A}} & \text { if } v_{i} \in A  \tag{6}\\ -\sqrt{\frac{\operatorname{vol}(\bar{A})}{\operatorname{vol}(\bar{A})}} & \text { if } v_{i} \in \bar{A}\end{cases}
$$

$\min _{A \subset V} \operatorname{Ncut}(A, \bar{A})$

$$
f^{\prime} L f=\operatorname{vol}(V) N c u t(A, \bar{A})
$$

$\min _{A} f^{\prime} L f$ subject to $f$ as in (6), $D f \perp \mathbb{1}, f^{\prime} D f=\operatorname{vol}(V)$

## Relaxation !!!

$\min _{f \in \mathbb{R}^{n}} f^{\prime} L f$ subject to $D f \perp \mathbb{1}, f^{\prime} D f=\operatorname{vol}(V)$
Substitute $g:=D^{1 / 2} f$
$\min _{g \in \mathbb{R}^{n}} g^{\prime} D^{-1 / 2} L D^{-1 / 2} g$ subject to $g \perp D^{\frac{1}{2}} \mathbb{1}, \quad\|g\|^{2}=\operatorname{vol}(V)$

## GRAPH CUT

- Approximation Ncut for arbitrary k

$$
h_{i, j}= \begin{cases}1 / \sqrt{\operatorname{vol}\left(A_{j}\right)} & \text { if } v_{i} \in A_{j} \\ 0 & \text { otherwise }\end{cases}
$$

Problem reformulation:


## Relaxation !!!

Re-substituting $H=D^{-1 / 2} T$
$\min _{T \in \mathbb{R}^{n \times k}} \operatorname{Tr}\left(T^{\prime} D^{-1 / 2} L D^{-1 / 2} T\right)$ subject to $T^{\prime} T=I$
Rayleigh-Ritz Theorem

T contains the first k eigenvectors of $L_{\text {sym }}$ as columns.
Re-substituting $H=D^{-1 / 2} T$, solution $H$ contains the first k eigenvectors of $L_{r w}$.

## Spectral Clustering

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## RANDOM WALK

- A random walk on a graph is a stochastic process which randomly jumps from vertex to vertex.
- Random walk stays long within the same cluster and seldom jumps between clusters.
- A balanced partition with a low cut will also have the property that the random walk does not have many opportunities to jump between clusters.


## RANDOM WALK

- Transition probability $p_{i j}$ of jumping from $v_{i}$ to $v_{j}$

$$
p_{i j}=w_{i j} / d_{i}
$$

- The transition matrix $P=\left(p_{i j}\right) \mathrm{i}, \mathrm{j}=1, \ldots, \mathrm{n}$ of random walk is defined by

$$
P=D^{-1} W
$$

- If the graph is connected and non-bipartite, the random walk always processes a unique stationary distribution $\pi=\left(\pi_{1}, \ldots, \pi_{n}\right)^{\prime}$, where $\pi_{i}=d_{i} / \operatorname{vol}(V) .\left(d_{i}=\sum_{j=1}^{n} w_{i j}, v o l(V)=\sum_{i \in v} d_{i}\right)$


## RANDOM WALK

- Relationship between $L_{r w}$ and $P$.

$$
L_{r w}=I-P
$$

- $\lambda$ is an eigenvalue of $L_{r w}$ with eigenvector $u$ if and only if $1-\lambda$ is an eigenvalue of $P$ with eigenvector $u$.
- The largest eigenvectors of $P$ and the smallest eigenvectors of $L_{r w}$ can be used to describe cluster properties of the graph.


## RANDOM WALK

- Random walks and Ncut

Proposition 5 (Ncut via transition probabilities) Let $G$ be connected and non bipartite. Assume that we run the random walk $\left(X_{t}\right)_{t \in N}$ starting with $X_{0}$ in the stationary distribution $\pi$. For disjoint subsets $A, B \subset V$, denote by $P(B \mid A):=P\left(X_{1} \in B \mid X_{0} \in A\right)$. Then:

$$
\operatorname{Ncut}(A, \bar{A})=P(\bar{A} \mid A)+P(A \mid \bar{A}) .
$$

## RANDOM WALK

- Random walks and Ncut

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$$
\operatorname{Ncut}(A, \bar{A})=P(\bar{A} \mid A)+P(A \mid \bar{A}) .
$$

Proof. First of all observe that

$$
\begin{aligned}
& P\left(X_{0} \in A, X_{1} \in B\right)=\sum_{i \in A, j \in B} P\left(X_{0}=i, X_{1}=j\right)=\sum_{i \in A, j \in B} \pi_{i} p_{i j} \\
& =\sum_{i \in A, j \in B} \frac{d_{i}}{\operatorname{vol}(V)} \frac{w_{i j}}{d_{i}}=\frac{1}{\operatorname{vol}(V)} \sum_{i \in A, j \in B} w_{i j}
\end{aligned}
$$

Using this we obtain

$$
\begin{gathered}
P\left(X_{1} \in B \mid X_{0} \in A\right)=\frac{P\left(X_{0} \in A, X_{1} \in B\right)}{P\left(X_{0} \in A\right)} \\
=\left(\frac{1}{\operatorname{vol}(V)} \sum_{i \in A, j \in B} w_{i j}\right)\left(\frac{\operatorname{vol}(A)}{\operatorname{vol}(V)}\right)^{-1}=\frac{\sum_{i \in A, j \in B} w_{i j}}{\operatorname{vol}(A)}
\end{gathered}
$$

Now the proposition follows directly with the definition of Ncut.

## RANDOM WALK

- Random walks and Ncut

Proposition 5 (Ncut via transition probabilities) Let $G$ be connected and non bipartite. Assume that we run the random walk $\left(X_{t}\right)_{t \in N}$ starting with $X_{0}$ in the stationary distribution $\pi$. For disjoint subsets $A, B \subset V$, denote by $P(B \mid A):=P\left(X_{1} \in B \mid X_{0} \in A\right)$. Then:

$$
\operatorname{Ncut}(A, \bar{A})=P(\bar{A} \mid A)+P(A \mid \bar{A})
$$

It tells us that when minimizing Ncut, we actually look for a cut through the graph such that $\mathbf{A}$ random walk seldom transitions from $A$ to $\bar{A}$ and vice versa.

## RANDOM WALK

- What is commute distance

The commute distance (resistance distance) $c_{i j}$ between two vertices $v_{i}$ and $v_{j}$ is the expected time it takes the random walk to travel from vertex $v_{i}$ to vertex $v_{j}$ and back.

The commute distance between two vertices decrease if there are many different short ways to get from vertex $v_{i}$ to vertex $v_{j}$.

Points which are connected by a short path in the graph and lie in the same high-density region of the graph are considered closer to each other than points which are connected by a short path but lie in different high-density regions of the graph.

## RANDOM WALK

- How to calculate commute distance

Generalized inverse (also called pseudo-inverse or Moore-Penrose inverse)
$L$ can be decomposed as $L=U \wedge U^{\prime}$, and $L$ is not invertible.
Define generalized inverse as $L^{\dagger}=U \wedge^{\dagger} U^{\prime}$, and $\Lambda^{\dagger}$ is the diagonal matrix with the eigenvalues $\lambda_{1}, \ldots, \lambda_{n}$ on the diagonal entries $1 / \lambda_{i}$ if $\lambda_{i} \neq 0$ and 0 if $\lambda_{i}=0$.

The entries of $L^{\dagger}$ can be computed as $l_{i j}^{\dagger}=\sum_{k=2}^{n} \frac{1}{\lambda_{k}} u_{i k} u_{j k}$

## RANDOM WALK

- How to calculate commute distance

Proposition 6 (Commute distance) Let $G=(V, E)$ a connected, undirected graph. Denote by $c_{i j}$ the commute distance between vertex $v_{i}$ and vertex $v_{j}$, and by $L^{\dagger}=\left(l_{i j}^{\dagger}\right)_{i, j=1, \ldots, n}$ the generalized inverse of $L$. Then we have:

$$
\begin{aligned}
c_{i j}=\operatorname{vol}(V)\left(l_{i i}^{\dagger}-2 l_{i j}^{\dagger}+l_{j j}^{\dagger}\right) & =\operatorname{vol}(V)\left(e_{i}-e_{j}\right)^{\prime} L^{\dagger}\left(e_{i}-e_{j}\right) \\
e_{i} & =(0, \ldots, 0,1,0, \ldots, 0)^{\prime} \text { as the } i \text {-th unit vector. }
\end{aligned}
$$

This is result has been published by Klein and Randic(1993), where it has been proved by methods of electrical network theory.

## RANDOM WALK

- Proposition 6's consequence

$$
c_{i j}=\operatorname{vol}(V)\left(l_{i i}^{\dagger}-2 l_{i j}^{\dagger}+l_{j j}^{\dagger}\right)=\operatorname{vol}(V)\left(e_{i}-e_{j}\right)^{\prime} L^{\dagger}\left(e_{i}-e_{j}\right)
$$

Construct
an embedding
Commute Distance - Euclidean Distance
between $v_{i}$ and $v_{j}$ between $z_{i}$ and $z_{j}$

Choose $z_{i}$ as the point in $\mathbb{R}^{n}$ corresponding to the $i$-th row of the matrix $U\left(\wedge^{\dagger}\right)^{1 / 2}$

$$
\left\langle z_{i}, z_{j}\right\rangle=e_{i}^{\prime} L^{\dagger} e_{j} \text { and } c_{i j}=\operatorname{vol}(V)\left\|z_{i}-z_{j}\right\|^{2}
$$

## RANDOM WALK

- A loose relation between spectral clustering and commute distance.


## Spectral Clustering

1. Map the vertices of the graph on the rows $y_{i}$ of the matrix $U$
2. Only take the first $k$ columns of the matrix

## Commute Distance

1. Map the vertices on the rows $z_{i}$ of the matrix $\left(\wedge^{\dagger}\right)^{1 / 2} U$
2. Commute time embedding takes all columns

Several authors justify that spectral clustering constructs clusters based on the Euclidean distances between the $y_{i}$ can be interpreted as building clusters of the vertices in the graph based on the commute distance.

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## PERTURBATION THEORY

- Perturbation theory studies the question of how eigenvalues and eigenvectors of a matrix $A$ change if we add a small perturbation $H$.

$$
\text { perturbed matrix } \tilde{A}:=A+H
$$

Perturbation theorems state that a certain distance between eigenvalues or eigenvectors of $A$ and $\bar{A}$ is bounded by a constant times a norm of $H$.

## PERTURBATION THEORY

Theorem 7 (Davis-Kahan) Let $A, H \in \mathbb{R}^{n \times n}$ be symmetric matrices, and let $\|\cdot\|$ be the Frobenius norm or the two-norm for matrices, respectively. Consider $\tilde{A}:=A+H$ as a perturbed version of $A$. Let $S_{1} \subset \mathbb{R}$ be an interval. Denote by $\sigma_{S_{1}}(A)$ the set of eigenvalues of $A$ which are contained in $S_{1}$, and by $V_{1}$ the eigenspace corresponding to all those eigenvalues (more formally, $V_{1}$ is the image of the spectral projection induced by $\sigma_{S_{1}}(A)$ ). Denote by $\sigma_{S_{1}}(\tilde{A})$ and $\tilde{V}_{1}$ the analogous quantities for $\tilde{A}$. Define the distance between $S_{1}$ and the spectrum of $A$ outside of $S_{1}$ as

$$
\delta=\min \left\{|\lambda-s| ; \lambda \text { eigenvalue of } A, \lambda \notin S_{1}, s \in S_{1}\right\} .
$$

Then the distance $d\left(V_{1}, \tilde{V}_{1}\right):=\left\|\sin \boldsymbol{\Theta}\left(V_{1}, \tilde{V}_{1}\right)\right\|$ between the two subspaces $V_{1}$ and $\tilde{V}_{1}$ is bounded by

$$
d\left(V_{1}, \tilde{V}_{1}\right) \leq \frac{\|H\|}{\delta} .
$$

The smaller the perturbation $H=L-\tilde{L}$ and the larger the eigengap $\left|\lambda_{k}-\lambda_{k+1}\right|$ is.

Below we will see that the size of the eigengap can also be used in a different context as a quality criterion for spectral clustering, namely when choosing the number $k$ of clusters to construct.

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## PRACTICAL DETAILS

- Constructing the similarity graph

1. Similarity Function Itself Make sure that points which are considered to be "very similar" by the similarity function are also closely related in the application the data comes from.
2. Type of Similarity Graph Which one to choose from those three types. General recommendation: $k$-nearest neighbor graph.
3. Parameters of Similarity $\operatorname{Graph}(k$ or $\varepsilon$ )
4. KNN: $k$ in order of $\log (n)$;
5. mutual KNN: k significantly larger than standard KNN ;
6. $\varepsilon$-neighborhood graph: longest edge of MST;
7. fully connected graph: $\sigma$ in order of the mean distance of a point to its $k$-th nearest neighbor. Or choose $k=\varepsilon$.

## PRACTICAL DETAILS

- Computing Eigenvectors

1. How to compute the first eigenvectors efficiently for large $L$
2. Numerical eigensolvers converge to some orthonormal basis of the eigenspace.

- Number of Clusters

1. General Strategies
2. Eigengap heuristic(Choose the number $k$ such that all eigenvalues
$\lambda_{1}, \ldots \lambda_{k}$ are very small, but $\lambda_{k+1}$ is relatively large)

## PRACTICAL DETAILS



Eigenvalues


Well Separated


Eigenvalues


More Blurry


Eigenvalues


Overlap So Much

Eigengap Heuristic usually works well if the data contains very well pronounced clusters, but in ambiguous cases it also returns ambiguous results.

## PRACTICAL DETAILS

- The k-means step

It is not necessary. People also use other techniques

- Which graph Laplacian should be used?

Look at the degree distribution. There are several arguments which advocate for using normalized rather than unnormalized spectral clustering, and in the normalized case to use the eigenvectors of $L_{r w}$ rather than those of $L_{\text {sym }}$

## PRACTICAL DETAILS

- Which graph Laplacian should be used?

Why normalized is better than unnormalized spectral clustering?

## Objective1:

1. We want to find a partition such that points in different clusters are dissimilar to each other, that is we want to minimize the between-cluster similarity. In the graph setting, this means to minimize $\operatorname{cut}(A, \bar{A})$.

## Both RatioCut and Ncut directly implement

Objective2:
2. We want to find a partition such that points in the same cluster are similar to each other, that is we want to maximize the within-cluster similarities $W(A, A)$, and $W(\bar{A}, \bar{A})$.

## Only Ncut implements

Normalized spectral clustering implements both clustering objectives mentioned above, while unnormalized spectral clustering only implements the first obejctive.

## PRACTICAL DETAILS

- Which graph Laplacian should be used?

Why the eigenvectors of $L_{r w}$ are better than those of $L_{s y m}$ ?

1. Eigenvectors of $L_{r w}$ are cluster indicator vectors $\mathbb{I}_{A_{i}}$, while the eigenvectors of $L_{\text {sym }}$ are additionally multiplied with $D^{1 / 2}$, which might lead to undesired artifacts.
2. Using $L_{\text {sym }}$ also does not have any computational advantages.

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## SPECTRAL CLUSTERING

## Thank you

