CS 1571 Introduction to AI Lecture 15

Inferences in first-order logic. Knowledge based systems.

Milos Hauskrecht

milos@cs.pitt.edu5329 Sennott Square

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Administration announcements

Midterm:

- Thursday, October 25, 2012
- In-class
- Closed book

What does it cover?

· All material covered by the end of lecture today

Homework 7:

• out today to practice inferences in FOL

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Tic-tac-toe competion

The programs were played against each other 10 times (5 times as first and 5 times as second player)

Winners:

- 1. Rishi Sadhir and Yuriy Koziy
- 3. Nathan Essel

Reward: 20% credit for the homework

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Sentences in Horn normal form

- Horn normal form (HNF) in the propositional logic
 - a special type of clause with at most one positive literal

$$(A \vee \neg B) \wedge (\neg A \vee \neg C \vee D)$$

Implicative form:

$$(B \Rightarrow A) \land ((A \land C) \Rightarrow D)$$

- A clause with one literal, e.g. A, is also called a fact
- A clause representing an implication (with a conjunction of positive literals in antecedent and one positive literal in consequent), is also called **a rule**
- Resolution rule (for clausal form) and modus ponens (for implicative form):
 - Both are complete inference rule for unit inferences for KBs in the Horn normal form.
 - Recall: Not all KBs are convertible to HNF !!!

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Horn normal form in FOL

- First-order logic (FOL)
 - adds variables and quantifiers, works with terms, predicates
- **HNF in FOL:** primitive sentences (propositions) are formed by predicates
- Inference rules: generalized versions with substitutions Example: modus ponens

$$\sigma = \text{a substitution s.t. } \forall i \ SUBST(\sigma, \phi_i') = SUBST(\sigma, \phi_i)$$

$$\frac{\phi_1', \phi_2' \dots, \phi_n', \quad \phi_1 \wedge \phi_2 \wedge \dots \phi_n \Rightarrow \tau}{SUBST(\sigma, \tau)}$$

- Generalized resolution and generalized modus ponens:
 - is **complete** for unit inferences for the KBs in HN;

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Forward and backward chaining

Two inference procedures based on modus ponens for **Horn KBs**:

- Forward chaining
 - **Idea:** Whenever the premises of a rule are satisfied, infer the conclusion. Continue with rules that became satisfied.
 - **Typical usage:** If we want to infer all sentences entailed by the existing KB.
- Backward chaining (goal reduction)
 - **Idea:** To prove the fact that appears in the conclusion of a rule prove the premises of the rule. Continue recursively.
 - **Typical usage:** If we want to prove that the target (goal) sentence α is entailed by the existing KB.

Both procedures are complete for unit inferences in KBs in Horn form !!!

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Forward chaining example

Forward chaining

Idea: Whenever the premises of a rule are satisfied, infer the conclusion. Continue with rules that became satisfied

Assume the KB with the following rules:

KB: R1: Steamboat $(x) \land Sailboat(y) \Rightarrow Faster(x, y)$

R2: Sailboat $(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

F3: RowBoat(PondArrow)

Theorem: Faster (Titanic, PondArrow)

?

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Forward chaining example

KB: R1: Steamboat $(x) \land Sailboat(y) \Rightarrow Faster(x, y)$

R2: $Sailboat(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

F3: *RowBoat(PondArrow)*

?

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Forward chaining example

KB: R1: Steamboat $(x) \land Sailboat(y) \Rightarrow Faster(x, y)$

> R2: $Sailboat(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

Sailboat (Mistral) F2:

RowBoat(PondArrow) F3:

Rule R1 is satisfied:

F4: Faster(Titanic, Mistral)



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Forward chaining example

KB: R1: Steamboat $(x) \land Sailboat (y) \Rightarrow Faster (x, y)$

> R2: $Sailboat(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

Sailboat (Mistral) F2:

RowBoat(PondArrow) F3:

Rule R1 is satisfied:

Faster(Titanic, Mistral)

Rule R2 is satisfied:

Faster(*Mistral*, *PondArrow*) ◀ F5:



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Forward chaining example

KB: R1: Steamboat $(x) \land Sailboat(y) \Rightarrow Faster(x, y)$

R2: $Sailboat(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

F3: RowBoat(PondArrow)

Rule R1 is satisfied:

F4: *Faster*(*Titanic*, *Mistral*)

-

Rule R2 is satisfied:

F5: Faster(Mistral, PondArrow)



Rule R3 is satisfied:

F6: *Faster*(*Titanic*, *PondArrow*)



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Backward chaining example

• Backward chaining (goal reduction)

Idea: To prove the fact that appears in the conclusion of a rule prove the antecedents (if part) of the rule & repeat recursively.

KB: R1: Steamboat $(x) \land Sailboat (y) \Rightarrow Faster (x, y)$

R2: Sailboat $(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

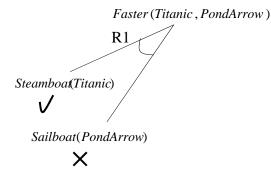
F2: Sailboat (Mistral)

F3: RowBoat(PondArrow)

Theorem: Faster (Titanic, PondArrow)

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Backward chaining example



F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

F3: *RowBoat(PondArrow)*

Steamboat $(x) \land Sailboat (y) \Rightarrow Faster (x, y)$

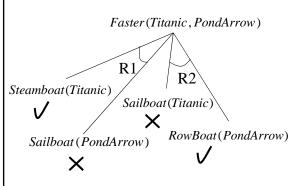
Faster (Titanic, PondArrow)

 $\{x / Titanic, y / PondArrow\}$

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Backward chaining example



F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

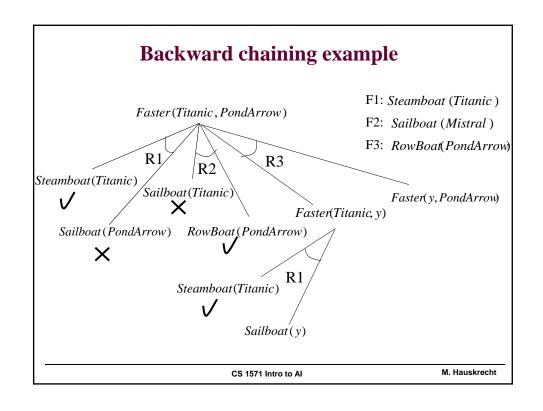
F3: *RowBoat(PondArrow)*

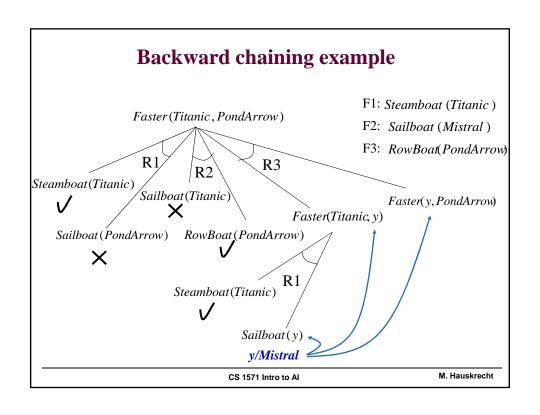
 $Sailboat(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

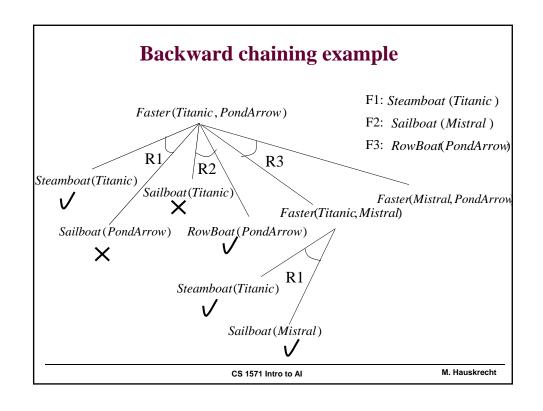
Faster (Titanic, PondArrow)

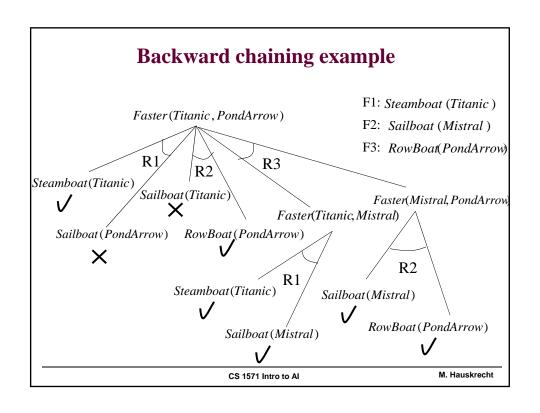
{ y / Titanic , z / PondArrow }

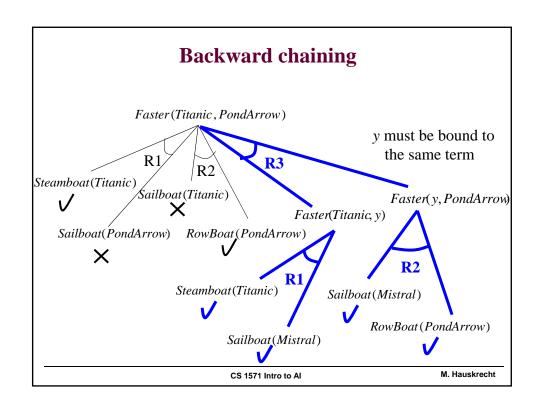
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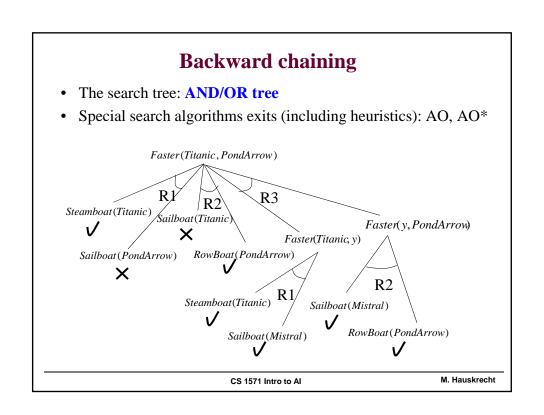












Knowledge-based system

Knowledge base

Inference engine

Knowledge base:

- A set of sentences that describe the world in some formal (representational) language (e.g. first-order logic)
- Domain specific knowledge

• Inference engine:

- A set of procedures that work upon the representational language and can infer new facts or answer KB queries (e.g. resolution algorithm, forward chaining)
- Domain independent

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Automated reasoning systems

Examples and main differences:

Theorem provers

 Prove sentences in the first-order logic. Use inference rules, resolution rule and resolution refutation.

Deductive retrieval systems

- Systems based on rules (KBs in Horn form)
- Prove theorems or infer new assertions (forward, backward chaining)

Production systems



- Systems based on rules with actions in antecedents
- Forward chaining mode of operation

Semantic networks



 Graphical representation of the world, objects are nodes in the graphs, relations are various links

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Production systems

Based on rules, but different from KBs in the Horn form Knowledge base is divided into:

- A Rule base (includes rules)
- A Working memory (includes facts)

A special type of if – then rule

$$p_1 \wedge p_2 \wedge \dots p_n \Rightarrow a_1, a_2, \dots, a_k$$

- Antecedent: a conjunction of literals
 - facts, statements in predicate logic
- Consequent: a conjunction of actions. An action can:
 - **ADD** the fact to the KB (working memory)
 - **REMOVE** the fact from the KB (consistent with logic?)
 - **QUERY** the user, etc ...

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Production systems

Based on rules, but different from KBs in the Horn form Knowledge base is divided into:

- A Rule base (includes rules)
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A special type of if – then rule

$$p_1 \wedge p_2 \wedge \dots p_n \Rightarrow a_1, a_2, \dots, a_k$$

- Antecedent: a conjunction of literals
 - facts, statements in predicate logic
- Consequent: a conjunction of actions. An action can:
 - **ADD** the fact to the KB (working memory)
 - **REMOVE** the fact from the KB ← !!! Different from logic
 - **QUERY** the user, etc ...

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Production systems

- Use forward chaining to do reasoning:
 - If the antecedent of the rule is satisfied (rule is said to be "active") then its consequent can be executed (it is "fired")
- **Problem:** Two or more rules are active at the same time. Which one to execute next?

R27 Conditions R27
$$\checkmark$$
 Actions R27 \checkmark R105 Conditions R105 \checkmark Actions R105

 Strategy for selecting the rule to be fired from among possible candidates is called conflict resolution

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Production systems

- Why is conflict resolution important? Or, why do we care about the order?
- Assume that we have two rules and the preconditions of both are satisfied:

R1:
$$A(x) \wedge B(x) \wedge C(y) \Rightarrow add D(x)$$

R2:
$$A(x) \wedge B(x) \wedge E(z) \Rightarrow delete \ A(x)$$

What can happen if rules are triggered in different order?

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Production systems

- Why is conflict resolution important? Or, Why do we care about the order?
- Assume that we have two rules and the preconditions of both are satisfied:

R1: $A(x) \wedge B(x) \wedge C(y) \Rightarrow add D(x)$

R2: $A(x) \wedge B(x) \wedge E(z) \Rightarrow delete \ A(x)$

- What can happen if rules are triggered in different order?
 - If R1 goes first, R2 condition is still satisfied and we infer D(x)
 - If R2 goes first we may never infer D(x)

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Production systems

- Problems with production systems:
 - Additions and Deletions can change a set of active rules;
 - If a rule contains variables, testing all instances in which the rule is active may require a large number of unifications.
 - Conditions of many rules may overlap, thus requiring to repeat the same unifications multiple times.
- Solution: Rete algorithm
 - gives more efficient solution for managing a set of active rules and performing unifications
 - Implemented in the system OPS-5 (used to implement XCON – an expert system for configuration of DEC computers)

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Rete algorithm

• Assume a set of rules:

$$A(x) \wedge B(x) \wedge C(y) \Rightarrow add D(x)$$

$$A(x) \wedge B(y) \wedge D(x) \Rightarrow add \ E(x)$$

$$A(x) \wedge B(x) \wedge E(z) \Rightarrow delete \ A(x)$$

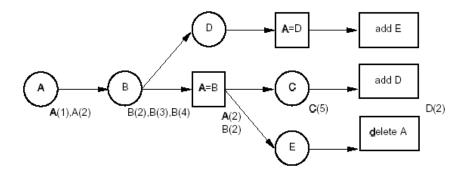
• And facts:

- Rete:
 - Compiles the rules to a network that merges conditions of multiple rules together (avoid repeats)
 - Propagates valid unifications
 - Reevaluates only changed conditions

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Rete algorithm. Network.



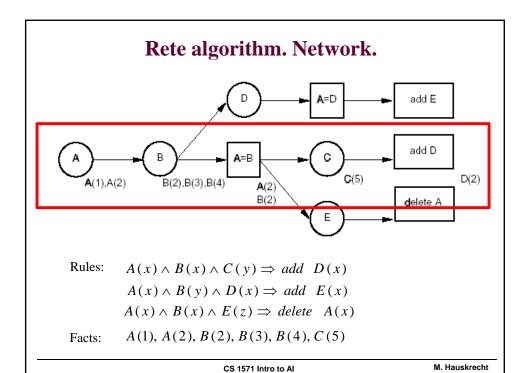
Rules: $A(x) \wedge B(x) \wedge C(y) \Rightarrow add D(x)$

 $A(x) \wedge B(y) \wedge D(x) \Rightarrow add E(x)$

 $A(x) \wedge B(x) \wedge E(z) \Rightarrow delete \ A(x)$

Facts: A(1), A(2), B(2), B(3), B(4), C(5)

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Conflict resolution strategies

- **Problem:** Two or more rules are active at the same time. Which one to execute next?
- Solutions:
 - **No duplication** (do not execute the same rule twice)
 - Recency. Rules referring to facts newly added to the working memory take precedence
 - **Specificity.** Rules that are more specific are preferred.
 - Priority levels. Define priority of rules, actions based on expert opinion. Have multiple priority levels such that the higher priority rules fire first.

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Semantic network systems

- Knowledge about the world described in terms of graphs. Nodes correspond to:
 - Concepts or objects in the domain.

Links to relations. Three kinds:

- Subset links (isa, part-of links)
- Inheritance relation links
- Member links (instance links)
- Function links.
- Can be transformed to the first-order logic language
- Graphical representation is often easier to work with
 - better overall view on individual concepts and relations

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