

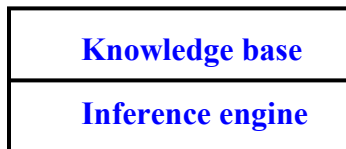
# CS 1571 Introduction to AI

## Lecture 11

### Knowledge Representation. Propositional logic.

**Milos Hauskrecht**  
[milos@cs.pitt.edu](mailto:milos@cs.pitt.edu)  
5329 Sennott Square

### Knowledge-based agent



- **Knowledge base (KB):**
  - A set of sentences that describe facts about the world in some formal (representational) language
  - **Domain specific**
- **Inference engine:**
  - A set of procedures that use the representational language to infer new facts from known ones or answer a variety of KB queries. Inferences typically require search.
  - **Domain independent**

## Example: MYCIN

- MYCIN: an expert system for diagnosis of bacterial infections
- **Knowledge base** represents
  - Facts about a specific patient case
  - Rules describing relations between entities in the bacterial infection domain

<b>If</b>	1. The stain of the organism is gram-positive, and 2. The morphology of the organism is coccus, and 3. The growth conformation of the organism is chains
<b>Then</b>	the identity of the organism is streptococcus

- **Inference engine:**
  - manipulates the facts and known relations to answer diagnostic queries (consistent with findings and rules)

## Knowledge representation

- The objective of knowledge representation is to express the knowledge about the world in a computer-tractable form
- Key aspects of knowledge representation languages:
  - **Syntax:** describes how sentences are formed in the language
  - **Semantics:** describes the meaning of sentences, what is it the sentence refers to in the real world
  - **Computational aspect:** describes how sentences and objects are manipulated in concordance with semantical conventions

**Many KB systems rely on some variant of logic**

## Logic

A formal language for expressing knowledge and for making logical inferences

**Logic** is defined by:

- **A set of sentences**
  - A sentence is constructed from a set of primitives according to syntax rules.
- **A set of interpretations**
  - An interpretation gives a semantic to primitives. It associates primitives with values.
- **The valuation (meaning) function  $V$** 
  - Assigns a value (typically the truth value) to a given sentence under some interpretation
$$V : \text{sentence} \times \text{interpretation} \rightarrow \{True, False\}$$

## Propositional logic

- **The simplest logic**
- **Definition:**
  - A **proposition** is a statement that is either true or false.
- **Examples:**
  - Pitt is located in the Oakland section of Pittsburgh.
    - (T)

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    - (T)
  - $5 + 2 = 8$ .
    - (F)
  - It is raining today.
    - ?

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    - (F)
  - It is raining today.
    - (either T or F)

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- Examples (cont.):
  - How are you?
    - ?

## Propositional logic

- **Examples (cont.):**
  - How are you?
    - a question is not a proposition
  - $x + 5 = 3$ 
    - ?

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  - $x + 5 = 3$ 
    - since  $x$  is not specified, neither true nor false
  - 2 is a prime number.
    - ?

## Propositional logic

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  - How are you?
    - a question is not a proposition
  - $x + 5 = 3$ 
    - since  $x$  is not specified, neither true nor false
  - 2 is a prime number.
    - (T)
  - She is very talented.
    - ?

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  - How are you?
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  - $x + 5 = 3$ 
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    - (T)
  - She is very talented.
    - since she is not specified, neither true nor false
  - There are other life forms on other planets in the universe.
    - ?

## Propositional logic

- **Examples (cont.):**
  - How are you?
    - a question is not a proposition
  - $x + 5 = 3$ 
    - since  $x$  is not specified, neither true nor false
  - 2 is a prime number.
    - (T)
  - She is very talented.
    - since she is not specified, neither true nor false
  - There are other life forms on other planets in the universe.
    - either T or F

## Propositional logic. Syntax

- **Formally propositional logic P:**
  - Is defined by Syntax+interpretation+semantics of P

### Syntax:

- **Symbols (alphabet)** in P:
  - Constants: *True, False*
  - Propositional symbols

Examples:

- $P$
- *Pitt is located in the Oakland section of Pittsburgh.,*
- *It rains outside,* etc.

- **A set of connectives:**

$\neg, \wedge, \vee, \Rightarrow, \Leftrightarrow$



## Propositional logic. Syntax

### Sentences in the propositional logic:

- **Atomic sentences:**
  - **Constructed from constants and propositional symbols**
  - True, False are (atomic) sentences
  - $P, Q$  or *Light in the room is on*, *It rains outside* are (atomic) sentences
- **Composite sentences:**
  - **Constructed from valid sentences via connectives**
  - If  $A, B$  are sentences then
$$\neg A \quad (A \wedge B) \quad (A \vee B) \quad (A \Rightarrow B) \quad (A \Leftrightarrow B)$$
or  $(A \vee B) \wedge (A \vee \neg B)$ are sentences

## Propositional logic. Semantics.

**The semantic gives the meaning to sentences.**

the semantics in the propositional logic is defined by:

- 1. Interpretation of propositional symbols and constants**
  - Semantics of atomic sentences
- 2. Through the meaning of connectives**
  - Meaning (semantics) of composite sentences

## Semantic: propositional symbols

### A **propositional symbol**

- a statement about the world that is either true or false

Examples:

- *Pitt is located in the Oakland section of Pittsburgh*
- *It rains outside*
- *Light in the room is on*

- An **interpretation** maps symbols to one of the two values: **True (T)**, or **False (F)**, depending on whether the symbol is satisfied in the world

**I**: *Light in the room is on* -> **True**, *It rains outside* -> **False**

**I'**: *Light in the room is on* -> **False**, *It rains outside* -> **False**

## Semantic: propositional symbols

The **meaning (value)** of the propositional symbol for a specific interpretation is given by its interpretation

**I**: *Light in the room is on* -> **True**, *It rains outside* -> **False**

$V(\text{Light in the room is on}, \mathbf{I}) = \text{True}$

$V(\text{It rains outside}, \mathbf{I}) = \text{False}$

**I'**: *Light in the room is on* -> **False**, *It rains outside* -> **False**

$V(\text{Light in the room is on}, \mathbf{I}') = \text{False}$

## Semantics: constants

- **The meaning (truth) of constants:**
  - True and False constants are always (under any interpretation) assigned the corresponding *True, False* value

$$\left. \begin{array}{l} V(\text{True}, \mathbf{I}) = \text{True} \\ V(\text{False}, \mathbf{I}) = \text{False} \end{array} \right\} \text{For any interpretation } \mathbf{I}$$

## Semantics: composite sentences.

- **The meaning (truth value) of complex propositional sentences.**
  - Determined using the standard rules of logic:

$P$	$Q$	$\neg P$	$P \wedge Q$	$P \vee Q$	$P \Rightarrow Q$	$P \Leftrightarrow Q$
<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>True</i>	<i>False</i>	<i>False</i>	<i>False</i>	<i>True</i>	<i>False</i>	<i>False</i>
<i>False</i>	<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>	<i>True</i>	<i>False</i>
<i>False</i>	<i>False</i>	<i>True</i>	<i>False</i>	<i>False</i>	<i>True</i>	<i>True</i>

## Translation

### Translation of English sentences to propositional logic:

- (1) identify atomic sentences that are propositions
- (2) Use logical connectives to translate more complex composite sentences that consist of many atomic sentences

### Assume the following sentence:

- It is not sunny this afternoon and it is colder than yesterday.

### Atomic sentences:

- $p$  = It is sunny this afternoon
- $q$  = it is colder than yesterday

**Translation:**  $\neg p \wedge q$

## Translation

### Assume the following sentences:

- It is not sunny this afternoon and it is colder than yesterday.
- We will go swimming only if it is sunny.
- If we do not go swimming then we will take a canoe trip.
- If we take a canoe trip, then we will be home by sunset.

### Denote:

- $p$  = It is sunny this afternoon
- $q$  = it is colder than yesterday
- $r$  = We will go swimming
- $s$  = we will take a canoe trip
- $t$  = We will be home by sunset

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### Assume the following sentences:

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- We will go swimming only if it is sunny.  $r \rightarrow p$
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- If we do not go swimming then we will take a canoe trip.  $\neg r \rightarrow s$
- If we take a canoe trip, then we will be home by sunset.  $s \rightarrow t$

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## Contradiction and Tautology

Some composite sentences may always (under any interpretation) evaluate to a single truth value:

- **Contradiction** (always *False*)

$$P \wedge \neg P$$

- **Tautology** (always *True*)

$$P \vee \neg P$$

$$\left. \begin{array}{l} \neg(P \vee Q) \Leftrightarrow (\neg P \wedge \neg Q) \\ \neg(P \wedge Q) \Leftrightarrow (\neg P \vee \neg Q) \end{array} \right\} \text{DeMorgan's Laws}$$

## Model, validity and satisfiability

- A **model (in logic)**: An interpretation is a model for a set of sentences if it assigns true to each sentence in the set.
- A sentence is **satisfiable** if it has a model;
  - There is at least one interpretation under which the sentence can evaluate to True.
- A sentence is **valid** if it is *True* in all interpretations
  - i.e., if its negation is **not satisfiable** (leads to contradiction)

<i>P</i>	<i>Q</i>	$P \vee Q$	$(P \vee Q) \wedge \neg Q$	$((P \vee Q) \wedge \neg Q) \Rightarrow P$
<i>True</i>	<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>
<i>True</i>	<i>False</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>False</i>	<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>
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<i>True</i>	<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>
<i>True</i>	<i>False</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>False</i>	<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>
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### Satisfiable sentence

$P$	$Q$	$P \vee Q$	$(P \vee Q) \wedge \neg Q$	$((P \vee Q) \wedge \neg Q) \Rightarrow P$
<i>True</i>	<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>
<i>True</i>	<i>False</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>False</i>	<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>
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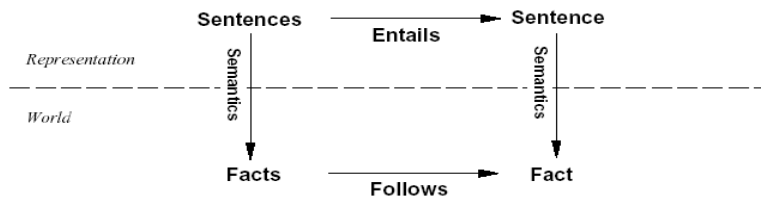
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		Satisfiable sentence		Valid sentence
$P$	$Q$	$P \vee Q$	$(P \vee Q) \wedge \neg Q$	$((P \vee Q) \wedge \neg Q) \Rightarrow P$
True	True	True	False	True
True	False	True	True	True
False	True	True	False	True
False	False	False	False	True

## Entailment

- **Entailment** reflects the relation of one fact in the world following from the others



- Entailment  $KB \models \alpha$
- Knowledge base KB entails sentence  $\alpha$  if and only if  $\alpha$  is true in all worlds where KB is true

## Sound and complete inference.

**Inference** is a process by which conclusions are reached.

- We want to implement the inference process on a computer !!

Assume an **inference procedure**  $i$  that

- derives a sentence  $\alpha$  from the KB :  $KB \vdash_i \alpha$

### Properties of the inference procedure in terms of entailment

- **Soundness:** An inference procedure is **sound**

If  $KB \vdash_i \alpha$  then it is true that  $KB \models \alpha$

- **Completeness:** An inference procedure is **complete**

If  $KB \models \alpha$  then it is true that  $KB \vdash_i \alpha$

## Logical inference problem

### Logical inference problem:

- **Given:**
  - a knowledge base KB (a set of sentences) and
  - a sentence  $\alpha$  (called **a theorem**),
- **Does a KB semantically entail  $\alpha$  ?**  $KB \models \alpha$  ?

In other words: In all interpretations in which sentences in the KB are true, is also  $\alpha$  true?

**Question:** Is there a procedure (program) that can decide this problem in a finite number of steps?

**Answer:** Yes. Logical inference problem for the propositional logic is **decidable**.

## Solving logical inference problem

In the following:

**How to design the procedure that answers:**

$$KB \models \alpha ?$$

**Three approaches:**

- **Truth-table approach**
- **Inference rules**
- **Conversion to the inverse SAT problem**
  - **Resolution-refutation**

## Truth-table approach

**Problem:**  $KB \models \alpha ?$

- We need to check all possible interpretations for which the KB is true (models of KB) whether  $\alpha$  is true for each of them

**Truth table:**

- enumerates truth values of sentences for all possible interpretations (assignments of True/False values to propositional symbols)

**Example:**

		KB		$\alpha$
$P$	$Q$	$P \vee Q$	$P \Leftrightarrow Q$	$(P \vee \neg Q) \wedge Q$
True	True	True	True	True
True	False	True	False	False
False	True	True	False	False
False	False	False	True	False

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## Truth-table approach

A two steps procedure:

1. Generate table for all possible interpretations
2. Check whether the sentence  $\alpha$  evaluates to true whenever  $KB$  evaluates to true

**Example:**  $KB = (A \vee C) \wedge (B \vee \neg C)$        $\alpha = (A \vee B)$

$A$	$B$	$C$	$A \vee C$	$(B \vee \neg C)$	$KB$	$\alpha$
<i>True</i>	<i>True</i>	<i>True</i>				
<i>True</i>	<i>True</i>	<i>False</i>				
<i>True</i>	<i>False</i>	<i>True</i>				
<i>True</i>	<i>False</i>	<i>False</i>				
<i>False</i>	<i>True</i>	<i>True</i>				
<i>False</i>	<i>True</i>	<i>False</i>				
<i>False</i>	<i>False</i>	<i>True</i>				
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$A$	$B$	$C$	$A \vee C$	$(B \vee \neg C)$	$KB$	$\alpha$
<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>True</i>	<i>True</i>	<i>False</i>	<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>True</i>	<i>False</i>	<i>True</i>	<i>True</i>	<i>False</i>	<i>False</i>	<i>True</i>
<i>True</i>	<i>False</i>	<i>False</i>	<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>False</i>	<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>	<i>True</i>
<i>False</i>	<i>True</i>	<i>False</i>	<i>False</i>	<i>True</i>	<i>False</i>	<i>True</i>
<i>False</i>	<i>False</i>	<i>True</i>	<i>True</i>	<i>False</i>	<i>False</i>	<i>False</i>
<i>False</i>	<i>False</i>	<i>False</i>	<i>False</i>	<i>True</i>	<i>False</i>	<i>False</i>

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True	True	True	True	True	True	True
True	True	False	True	True	True	True
True	False	True	True	False	False	True
True	False	False	True	True	True	True
False	True	True	True	True	True	True
False	True	False	False	True	False	True
False	False	True	True	False	False	False
False	False	False	False	True	False	False



## Truth-table approach

$$KB = (A \vee C) \wedge (B \vee \neg C) \quad \alpha = (A \vee B)$$

$A$	$B$	$C$	$A \vee C$	$(B \vee \neg C)$	$KB$	$\alpha$
True	True	True	True	True	True	True
True	True	False	True	True	True	True
True	False	True	True	False	False	True
True	False	False	True	True	True	True
False	True	True	True	True	True	True
False	True	False	False	True	False	True
False	False	True	True	False	False	False
False	False	False	False	True	False	False

KB entails  $\alpha$

- The **truth-table approach** is **sound and complete** for the propositional logic!!