

CS 1571 Introduction to AI

Lecture 8

Constraint satisfaction search

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Constraint satisfaction problem (CSP)

Objective:

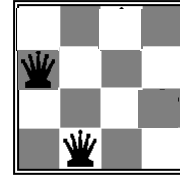
- **Find a configuration** satisfying goal conditions
- Constraint satisfaction problem (CSP) is a configuration search problem where:
 - A **state** is defined by a **set of variables and their values**
 - **Goal condition** is represented by a **set constraints on possible variable values**

CSP example: N-queens

Goal: n queens placed in non-attacking positions on the board

Variables:

- Represent queens, one for each column:
 - Q_1, Q_2, Q_3, Q_4
- Values:
 - Row placement of each queen on the board
 $\{1, 2, 3, 4\}$



$$Q_1 = 2, Q_2 = 4$$

Constraints: $Q_i \neq Q_j$ Two queens not in the same row
 $|Q_i - Q_j| \neq |i - j|$ Two queens not on the same diagonal

CSP example: Satisfiability (SAT) problem

Determine whether a sentence in the conjunctive normal form (CNF) is satisfiable (can evaluate to true)

- Used in the propositional logic (covered later)

$$(P \vee Q \vee \neg R) \wedge (\neg P \vee \neg R \vee S) \wedge (\neg P \vee Q \vee \neg T) \dots$$

Variables:

- Propositional symbols (P, R, T, S)
- Values: *True*, *False*

Constraints:

- Every conjunct must evaluate to true, at least one of the literals must evaluate to true

$$(P \vee Q \vee \neg R) \equiv \text{True}, (\neg P \vee \neg R \vee S) \equiv \text{True}, \dots$$

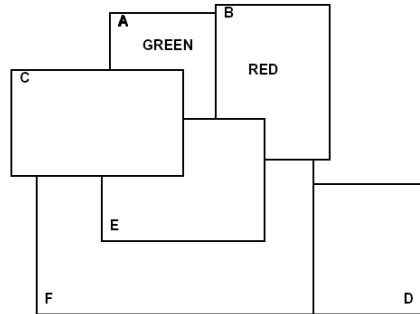
CSP example: Map coloring

Color a map using k different colors such that no adjacent countries have the same color

Variables: ?

- Variable values: ?

Constraints: ?



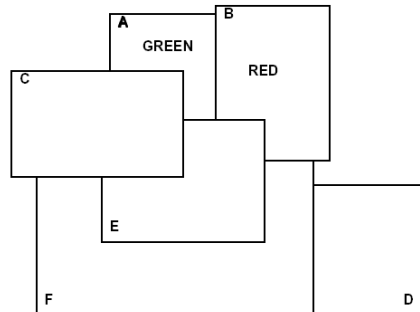
Map coloring

Color a map using k different colors such that no adjacent countries have the same color

Variables:

- Represent countries
 - A, B, C, D, E
- Values:
 - K -different colors
 - $\{\text{Red, Blue, Green, ...}\}$

Constraints: ?

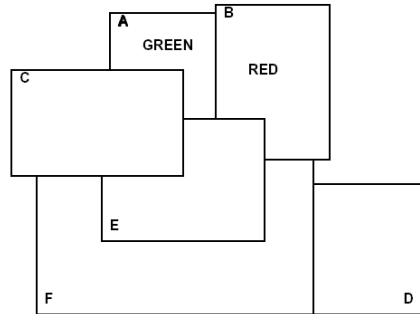


Map coloring

Color a map using k different colors such that no adjacent countries have the same color

Variables:

- Represent countries
 - A, B, C, D, E
- Values:
 - K -different colors
 - $\{\text{Red, Blue, Green, ...}\}$



Constraints: $A \neq B, A \neq C, C \neq E$, etc

An example of a problem with **binary constraints**

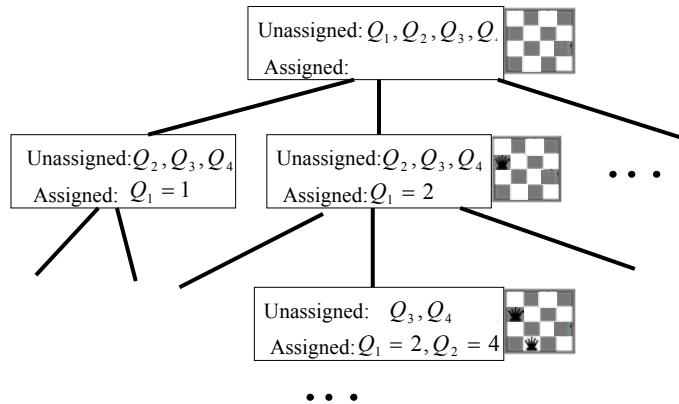
Constraint satisfaction as a search problem

Formulation of a CSP as a search problem:

- **States.** Assignment (partial, complete) of values to variables.
- **Initial state.** No variable is assigned a value.
- **Operators.** Assign a value to one of the unassigned variables.
- **Goal condition.** All variables are assigned, no constraints are violated.
- **Constraints** can be **represented**:
 - **Explicitly** by a set of allowable values
 - **Implicitly** by a function that tests for the satisfaction of constraints

Solving a CSP through standard search

- **Maximum depth of the tree (m): ?**
- **Depth of the solution (d) : ?**
- **Branching factor (b) : ?**

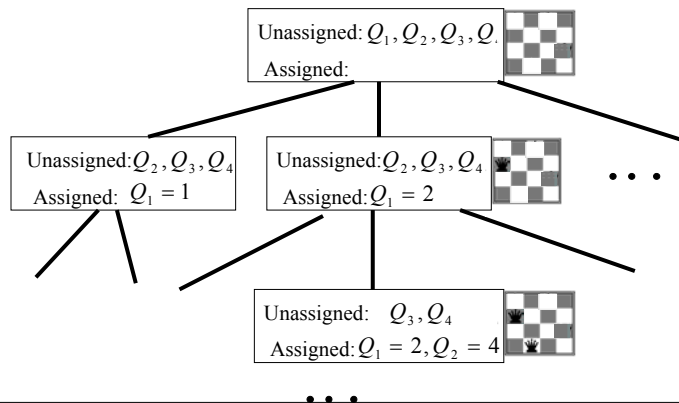


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Solving a CSP through standard search

- **Maximum depth of the tree:** Number of variables in the CSP
- **Depth of the solution:** Number of variables in the CSP
- **Branching factor:** if we fix the order of variable assignments the branch factor depends on the number of their values



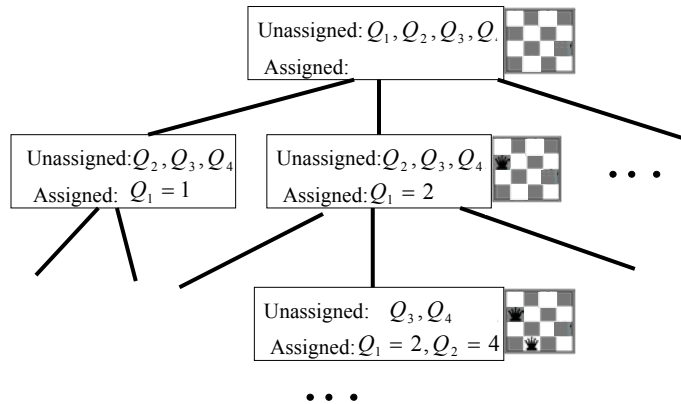
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Solving a CSP through standard search

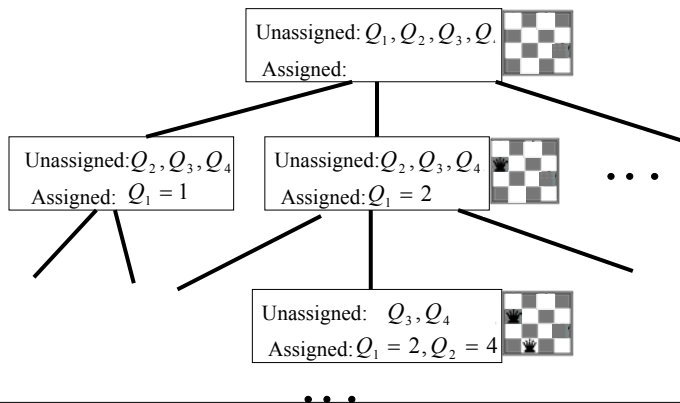
- What search algorithm to use: ?

Depth of the tree = Depth of the solution = number of vars



Solving a CSP through standard search

- What search algorithm to use: **Depth first search !!!**
 - Since we know the depth of the solution
 - We do not have to keep large number of nodes in queues



Backtracking

Depth-first search for CSP is also referred to as **backtracking**

The violation of constraints needs to be checked for each node, either during its generation or before its expansion

Consistency of constraints:

- Current **variable assignments** together with constraints **restrict remaining legal values of unassigned variables**;
- The remaining **legal and illegal values of variables may be inferred** (effect of constraints propagates)
- To prevent “blind” exploration it is necessary to keep track of the remaining legal values, so we know when the constraints are violated and when to terminate the search

Constraint propagation

A **state** (more broadly) is defined by a set of variables, their values and a list of legal and illegal assignments for unassigned variables

Legal and illegal assignments can be represented via: **equations** (value assignments) and **disequations (list of invalid assignments)**

Example: map coloring

Equation $A = \text{Red}$

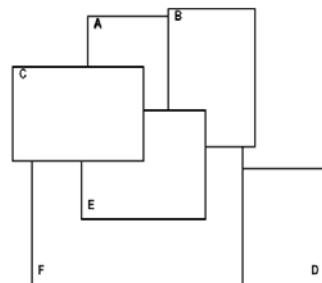
Disequation $C \neq \text{Red}$

Constraints + assignments

can entail new equations and disequations

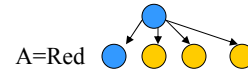
$A = \text{Red} \rightarrow B \neq \text{Red}$

Constraint propagation: the process of inferring of new equations and disequations from existing equations and disequations



Constraint propagation

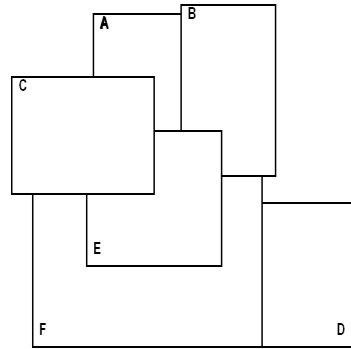
- Assign A=Red



	Red	Blue	Green
A	✓		
B			
C			
D			
E			
F			

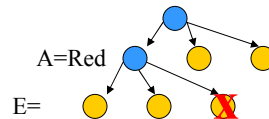


✓ - equations ✗ - disequations



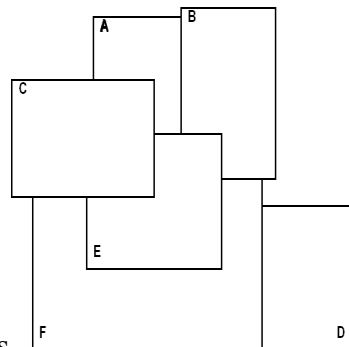
Constraint propagation

- Assign A=Red



	Red	Blue	Green
A	✓		
B	✗		
C	✗		
D			
E	✗		
F			

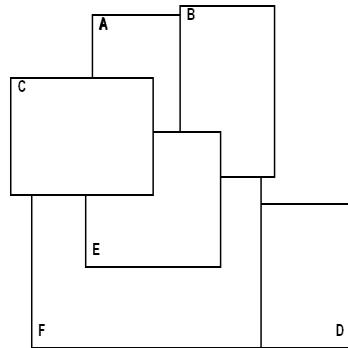
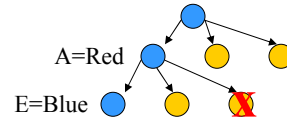
✓ - equations ✗ - disequations



Constraint propagation

- Assign E=Blue

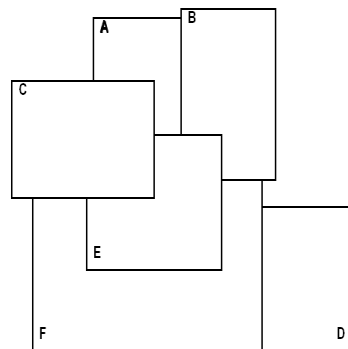
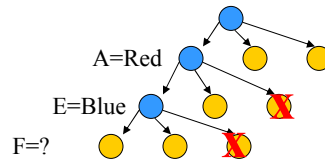
	Red	Blue	Green
A	✓		
B	✗		
C	✗		
D			
E	✗	✓	
F			



Constraint propagation

- Assign E=Blue

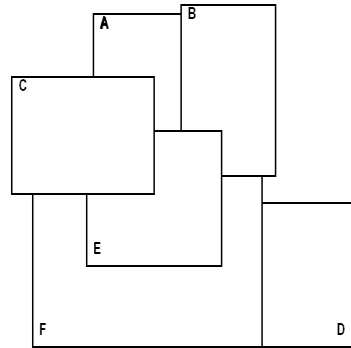
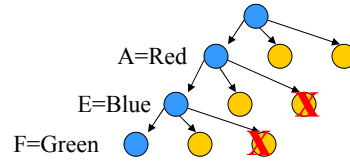
	Red	Blue	Green
A	✓	✗	
B	✗	✗	
C	✗	✗	
D			
E	✗	✓	
F		✗	



Constraint propagation

- Assign F=Green

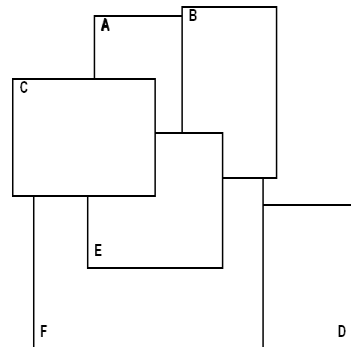
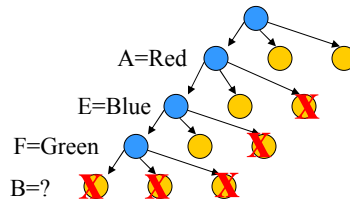
	Red	Blue	Green
A	✓	✗	
B	✗	✗	
C	✗	✗	
D			
E	✗	✓	
F		✗	✓



Constraint propagation

- Assign F=Green

	Red	Blue	Green
A	✓	✗	
B	✗	✗	✗
C	✗	✗	✗
D			✗
E	✗	✓	✗
F		✗	✓

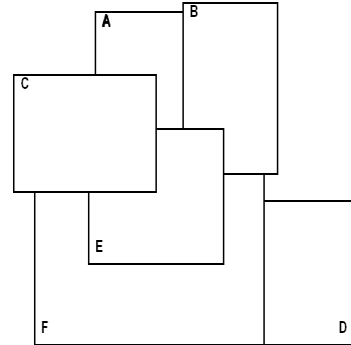
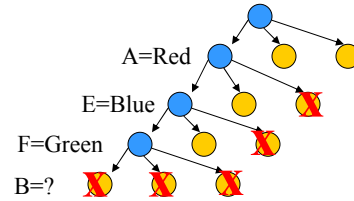


Constraint propagation

- Assign F=Green

	Red	Blue	Green
A	✓	✗	
B	✗	✗	✗
C	✗	✗	✗
D			✗
E	✗	✓	✗
F		✗	✓

Conflict !!! No legal assignments available for B and C



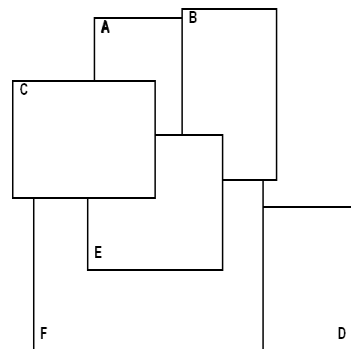
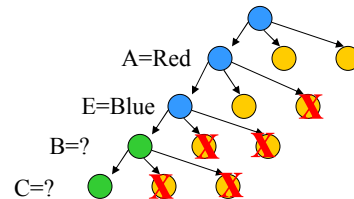
Constraint propagation

- We can derive remaining legal values through propagation

	Red	Blue	Green
A	✓	✗	
B	✗	✗	✓
C	✗	✗	✓
D			
E	✗	✓	
F		✗	

B=Green

C=Green

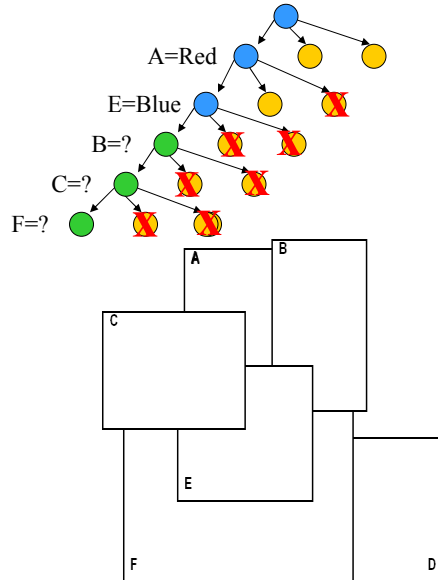


Constraint propagation

- We can derive remaining legal values through propagation

	Red	Blue	Green
A	✓	✗	✗
B	✗	✗	✓
C	✗	✗	✓
D	✗		
E	✗	✓	✗
F	✓	✗	✗

B=Green
 C=Green
 ➔
 F=Red

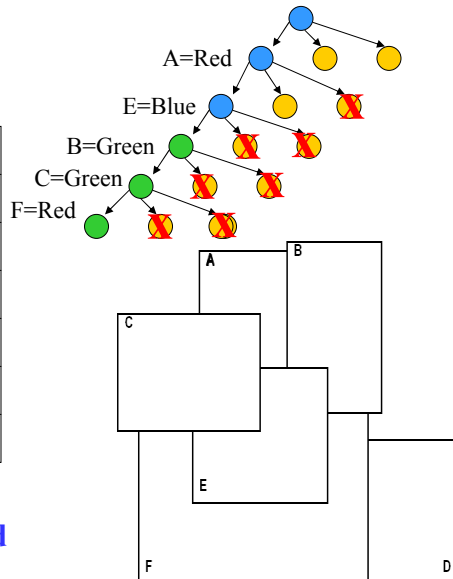


Constraint propagation

- We can derive remaining legal values through propagation

	Red	Blue	Green
A	✓	✗	✗
B	✗	✗	✓
C	✗	✗	✓
D	✗		
E	✗	✓	✗
F	✓	✗	✗

B=Green
 C=Green
 ➔
 F=Red



Constraint propagation

Three known techniques for propagating the effects of past assignments and constraints:

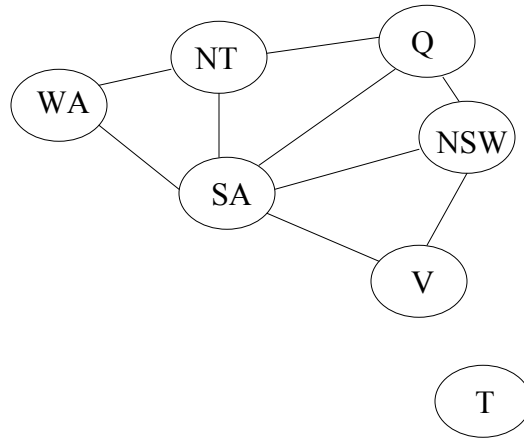
- **Value propagation**
- **Arc consistency**
- **Forward checking**
- **Difference:**
 - Completeness of inferences
 - Time complexity of inferences.

Constraint propagation

1. **Value propagation. Infers:**
 - **equations from** the set of **equations** defining the partial assignment, **and a constraint**
 2. **Arc consistency. Infers:**
 - **disequations from** the set of **equations and disequations** defining the partial assignment, and **a constraint**
 - **equations through** the exhaustion of alternatives
 3. **Forward checking. Infers:**
 - **disequations from** a set of **equations** defining the partial assignment, and a constraint
 - **Equations through** the exhaustion of alternatives
- Restricted forward checking:**
- uses only active constraints (active constraint – only one variable unassigned in the constraint)

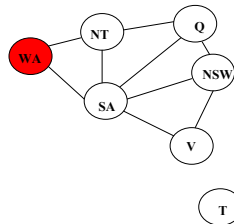
Example

Map coloring of Australia territories



Example: forward checking

Map coloring

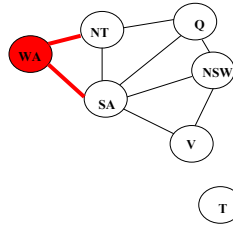


Set: **WA=Red**

<i>vars</i>	WA	NT	Q	NSW	V	SA	T
<i>domain</i>	R G B	R G B	R G B	R G B	R G B	R G B	R G B
<i>WA=Red</i>	R	?	?	?	?	?	?

Example: forward checking

Map coloring

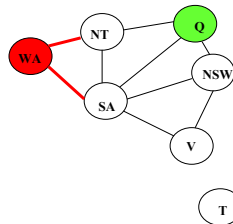


Set: WA=Red

vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B

Example: forward checking

Map coloring

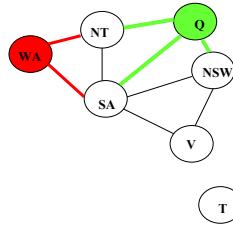


Set: Q=Green

vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B
Q=Green	R	?	G	?	?	?	?

Example: forward checking

Map coloring

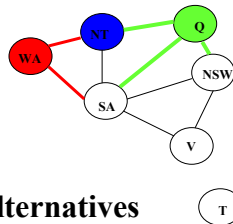


Set: Q=Green

vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B
Q=Green	R	B	G	R B	R G B	B	R G B

Example: forward checking

Map coloring

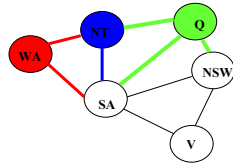


Infer: Exhaustions of alternatives

vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B
Q=Green	R	B	G	R B	R G B	B	R G B
Infer NT	R	B	G	?	?	?	?

Example: forward checking

Map coloring



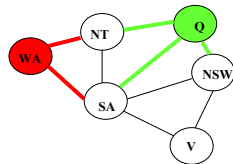
Infer: Exhaustions of alternatives



vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B
Q=Green	R	B	G	R B	R G B	B	R G B
Infer NT	R	B	G	R B	R G B	!	R G B

Example: arc consistency

Map coloring



Set: WA=Red

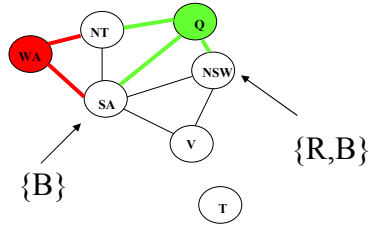
Set: Q=Green



vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B
Q=Green	R	B	G	R B	R G B	B	R G B

Example: arc consistency

Map coloring



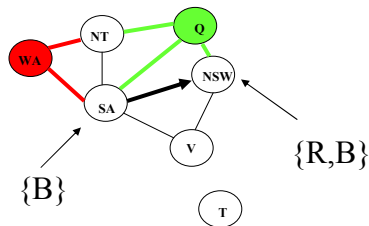
Set: WA=Red

Set: Q=Green

vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B
Q=Green	R	B	G	R B	R G B	B	R G B

Example: arc consistency

Map coloring



Set: WA=Red

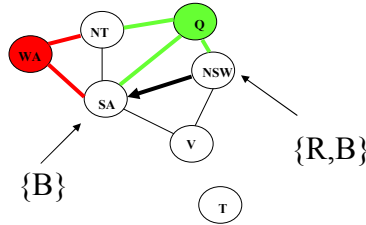
Set: Q=Green

vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B
Q=Green	R	B	G	R B	R G B	B	R G B

Example: arc consistency

Map coloring

Set: WA=Red
Set: Q=Green



NSW=B
SA=!

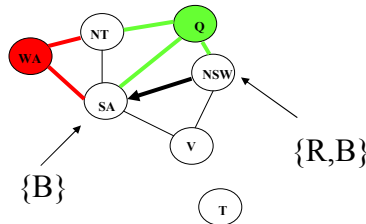
Inconsistent assignment

vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B
Q=Green	R	B	G	R B	R G B	B	R G B

Example: arc consistency

Map coloring

Set: WA=Red
Set: Q=Green



NSW=B
SA=!

Inconsistent assignment

vars	WA	NT	Q	NSW	V	SA	T
domain	R G B	R G B	R G B	R G B	R G B	R G B	R G B
WA=Red	R	G B	R G B	R G B	R G B	G B	R G B
Q=Green	R	B	G	R B R	R G B	B	R G B

Heuristics for CSPs

CSP searches the space in the depth-first manner.

But we still can choose:

- Which variable to assign next?
- Which value to choose first?

Heuristics

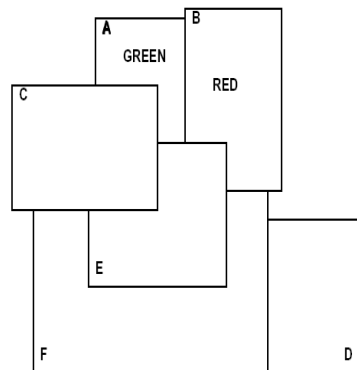
- **Most constrained variable**
 - Which variable is likely to become a bottleneck?
- **Least constraining value**
 - Which value gives us more flexibility later?

Heuristics for CSP

Examples: **map coloring**

Heuristics

- **Most constrained variable**
 - ?
- **Least constraining value**
 - ?

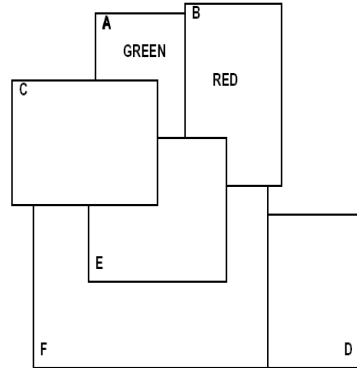


Heuristics for CSP

Examples: **map coloring**

Heuristics

- **Most constrained variable**
 - Country E is the most constrained one (cannot use Red, Green)
- **Least constraining value**
 - ?

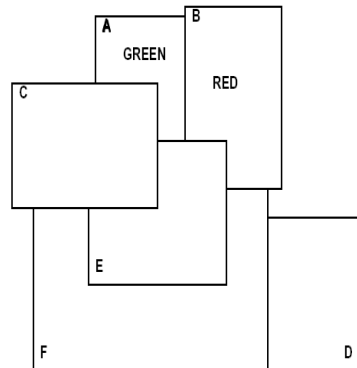


Heuristics for CSP

Examples: **map coloring**

Heuristics

- **Most constrained variable**
 - Country E is the most constrained one (cannot use Red, Green)
- **Least constraining value**
 - Assume we have chosen variable C
 - What color is the least constraining color?

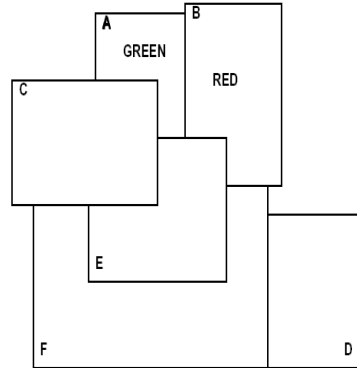


Heuristics for CSP

Examples: **map coloring**

Heuristics

- **Most constrained variable**
 - Country E is the most constrained one (cannot use Red, Green)
- **Least constraining value**
 - Assume we have chosen variable C
 - Red is the least constraining valid color for the future



Finding optimal configurations

Search for the optimal configuration

Constraint satisfaction problem:

Objective: find a configuration that satisfies all constraints



Optimal configuration problem:

Objective: find the best configuration

The quality of a configuration: is defined by some quality measure that reflects our **preference towards each configuration** (or state)

Our goal: optimize the configuration according to the quality measure also referred to as objective function

Search for the optimal configuration

If the space of configurations we search among is

- **Discrete or finite**
 - then it is a **combinatorial optimization problem**
- **Continuous**
 - then it is a **parametric optimization problem**

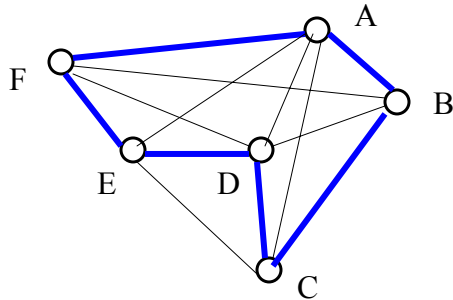
In the following we cover with combinatorial optimization problems

Parametric optimization will be covered next lecture.

Example: Traveling salesman problem

Problem:

- A graph with distances

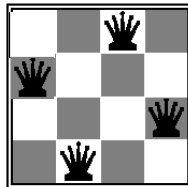


- **Goal:** find the shortest tour which visits every city once and returns to the start

An example of a valid **tour**: ABCDEF

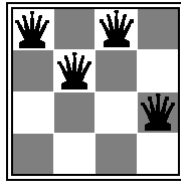
Example: N queens

- A CSP problem
- Is it possible to formulate the problem as an optimal configuration search problem ?

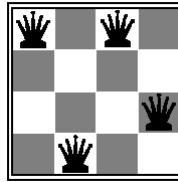


Example: N queens

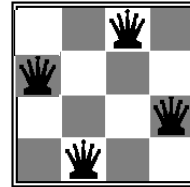
- A CSP problem
- Is it possible to formulate the problem as an optimal configuration search problem ? **Yes.**
- **The quality of a configuration in a CSP** can be measured by the number of violated constraints
- **Solving:** minimize the number of constraint violations



of violations =3



of violations =1



of violations =0

Iterative optimization methods

- Searching systematically for the best configuration with the **DFS** may not be the best solution
- Worst case running time:
 - Exponential in the number of variables
- Solutions to **large ‘optimal’ configuration** problems are often found using iterative optimization methods
- **Methods:**
 - **Hill climbing**
 - **Simulated Annealing**
 - **Genetic algorithms**