CS 1571 Introduction to AI Lecture 17

Inference in first-order logic. Production systems.

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Sentences in Horn normal form

- Horn normal form (HNF) in the propositional logic
 - a special type of clause with at most one positive literal

$$(A \lor \neg B) \land (\neg A \lor \neg C \lor D)$$

Typically written as: $(B \Rightarrow A) \land ((A \land C) \Rightarrow D)$

- A clause with one literal, e.g. A, is also called a fact
- A clause representing an implication (with a conjunction of positive literals in antecedent and one positive literal in consequent), is also called a rule
- Generalized Modus ponens:
 - is the complete inference rule for KBs in the Horn normal form. Not all KBs are convertible to HNF !!!

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Horn normal form in FOL

First-order logic (FOL)

adds variables and quantifiers, works with terms
 Generalized modus ponens rule:

$$\sigma = \text{a substitution s.t. } \forall i \ SUBST(\sigma, \phi_i') = SUBST(\sigma, \phi_i)$$

$$\underline{\phi_1', \phi_2' \dots, \phi_n', \quad \phi_1 \land \phi_2 \land \dots \phi_n \Rightarrow \tau}$$

$$\underline{SUBST(\sigma, \tau)}$$

Generalized modus ponens:

- is **complete** for the KBs with sentences in Horn form;
- Not all first-order logic sentences can be expressed in this form

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Forward and backward chaining

Two inference procedures based on modus ponens for **Horn KBs**:

Forward chaining

Idea: Whenever the premises of a rule are satisfied, infer the conclusion. Continue with rules that became satisfied.

Typical usage: infer all sentences entailed by the existing KB.

• Backward chaining (goal reduction)

Idea: To prove the fact that appears in the conclusion of a rule prove the premises of the rule. Continue recursively.

Typical usage: If we want to prove that the target (goal) sentence α is entailed by the existing KB.

Both procedures are complete for KBs in Horn form !!!

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Forward chaining example

Forward chaining

Idea: Whenever the premises of a rule are satisfied, infer the conclusion. Continue with rules that became satisfied Assume the KB with the following rules:

KB: R1: Steamboat $(x) \land Sailboat (y) \Rightarrow Faster (x, y)$

R2: Sailboat $(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

F3: RowBoat(PondArrow)

Theorem: Faster (Titanic, PondArrow)

?

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Forward chaining example

KB: R1: Steamboat $(x) \land Sailboat (y) \Rightarrow Faster (x, y)$

R2: $Sailboat(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

F3: RowBoat(PondArrow)

?

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Forward chaining example

KB: R1: Steamboat $(x) \land Sailboat (y) \Rightarrow Faster (x, y)$

R2: Sailboat $(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \wedge Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

F3: RowBoat(PondArrow)

Rule R1 is satisfied:

F4: Faster(Titanic, Mistral)



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Forward chaining example

KB: R1: Steamboat $(x) \land Sailboat (y) \Rightarrow Faster (x, y)$

R2: $Sailboat(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

F3: RowBoat(PondArrow)

Rule R1 is satisfied:

F4: Faster(Titanic, Mistral)



Rule R2 is satisfied:

F5: Faster(Mistral, PondArrow)

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Forward chaining example

KB: R1: Steamboat $(x) \land Sailboat (y) \Rightarrow Faster (x, y)$

R2: Sailboat $(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

F3: *RowBoat(PondArrow)*

Rule R1 is satisfied:

F4: *Faster*(*Titanic*, *Mistral*)

Rule R2 is satisfied:

F5: Faster(Mistral, PondArrow)

Rule R3 is satisfied:

F6: Faster(Titanic, PondArrow)



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Backward chaining example

• Backward chaining (goal reduction)

Idea: To prove the fact that appears in the conclusion of a rule prove the antecedents (if part) of the rule & repeat recursively.

KB: R1: Steamboat $(x) \land Sailboat (y) \Rightarrow Faster (x, y)$

R2: Sailboat $(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

R3: $Faster(x, y) \land Faster(y, z) \Rightarrow Faster(x, z)$

F1: Steamboat (Titanic)

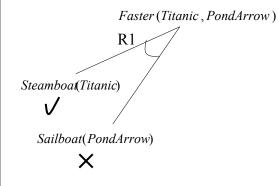
F2: Sailboat (Mistral)

F3: RowBoat(PondArrow)

Theorem: Faster (Titanic, PondArrow)

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Backward chaining example



F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

F3: RowBoat(PondArrow)

Steamboat $(x) \land Sailboat (y) \Rightarrow Faster (x, y)$

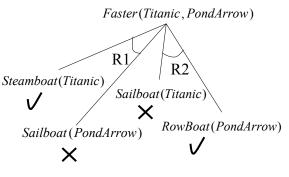
Faster (Titanic, PondArrow)

 $\{x \mid Titanic, y \mid PondArrow\}$

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Backward chaining example



F1: Steamboat (Titanic)

F2: Sailboat (Mistral)

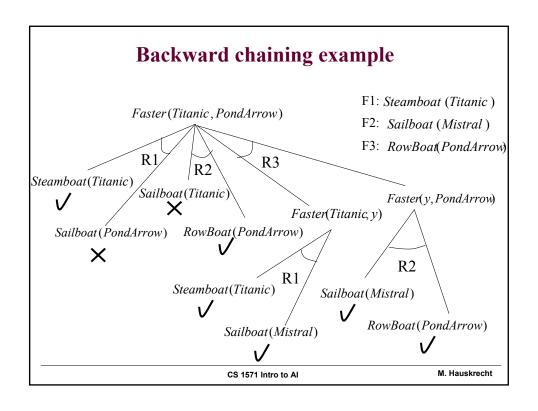
F3: RowBoat(PondArrow)

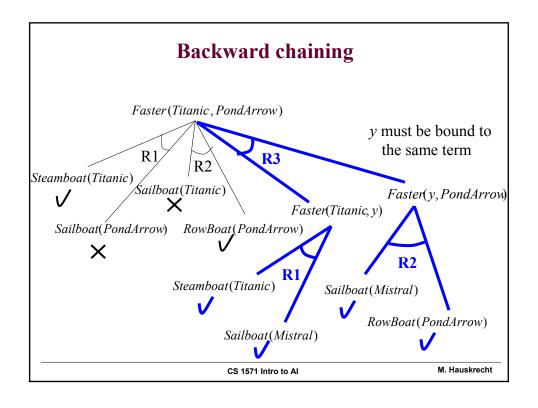
 $Sailboat(y) \land RowBoat(z) \Rightarrow Faster(y, z)$

Faster (Titanic, PondArrow)

 $\{y \mid Titanic, z \mid PondArrow\}$

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Knowledge-based system

Knowledge base

Inference engine

Knowledge base:

- A set of sentences that describe the world in some formal (representational) language (e.g. first-order logic)
- Domain specific knowledge

Inference engine:

- A set of procedures that work upon the representational language and can infer new facts or answer KB queries (e.g. resolution algorithm, forward chaining)
- Domain independent

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Automated reasoning systems

Examples and main differences:

Theorem provers

 Prove sentences in the first-order logic. Use inference rules, resolution rule and resolution refutation.

Deductive retrieval systems

- Systems based on rules (KBs in Horn form)
- Prove theorems or infer new assertions (forward, backward chaining)

Production systems



- Systems based on rules with actions in antecedents
- Forward chaining mode of operation

Semantic networks



 Graphical representation of the world, objects are nodes in the graphs, relations are various links

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Production systems

Based on rules, but different from KBs in the Horn form Knowledge base is divided into:

- A Rule base (includes rules)
- A Working memory (includes facts)

A special type of if – then rule

$$p_1 \wedge p_2 \wedge \dots p_n \Rightarrow a_1, a_2, \dots, a_k$$

- Antecedent: a conjunction of literals
 - facts, statements in predicate logic
- Consequent: a conjunction of actions. An action can:
 - **ADD** the fact to the KB (working memory)
 - **REMOVE** the fact from the KB (consistent with logic?)
 - QUERY the user, etc ...

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Production systems

Based on rules, but different from KBs in the Horn form Knowledge base is divided into:

- A Rule base (includes rules)
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A special type of if – then rule

$$p_1 \wedge p_2 \wedge \dots p_n \Rightarrow a_1, a_2, \dots, a_k$$

- Antecedent: a conjunction of literals
 - facts, statements in predicate logic
- Consequent: a conjunction of actions. An action can:
 - **ADD** the fact to the KB (working memory)
 - REMOVE the fact from the KB ← !!! Different from logic
 - **QUERY** the user, etc ...

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Production systems

- Use forward chaining to do reasoning:
 - If the antecedent of the rule is satisfied (rule is said to be "active") then its consequent can be executed (it is "fired")
- **Problem:** Two or more rules are active at the same time. Which one to execute next?

R27 Conditions R27
$$\checkmark$$
 Actions R27 R105 Conditions R105 \checkmark Actions R105

• Strategy for selecting the rule to be fired from among possible candidates is called **conflict resolution**

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Production systems

- Why is conflict resolution important? Or, why do we care about the order?
- Assume that we have two rules and the preconditions of both are satisfied:

R1:
$$A(x) \wedge B(x) \wedge C(y) \Rightarrow add D(x)$$

R2:
$$A(x) \wedge B(x) \wedge E(z) \Rightarrow delete \ A(x)$$

• What can happen if rules are triggered in different order?

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Production systems

- Why is conflict resolution important? Or, Why do we care about the order?
- Assume that we have two rules and the preconditions of both are satisfied:

R1:
$$A(x) \wedge B(x) \wedge C(y) \Rightarrow add D(x)$$

R2:
$$A(x) \wedge B(x) \wedge E(z) \Rightarrow delete \ A(x)$$

- What can happen if rules are triggered in different order?
 - If R1 goes first, R2 condition is still satisfied and we infer D(x)
 - If R2 goes first we may never infer D(x)

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Production systems

- Problems with production systems:
 - Additions and Deletions can change a set of active rules;
 - If a rule contains variables testing all instances in which the rule is active may require a large number of unifications.
 - Conditions of many rules may overlap, thus requiring to repeat the same unifications multiple times.
- Solution: Rete algorithm
 - gives more efficient solution for managing a set of active rules and performing unifications
 - Implemented in the system OPS-5 (used to implement XCON – an expert system for configuration of DEC computers)

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Rete algorithm

• Assume a set of rules:

$$A(x) \wedge B(x) \wedge C(y) \Rightarrow add \ D(x)$$

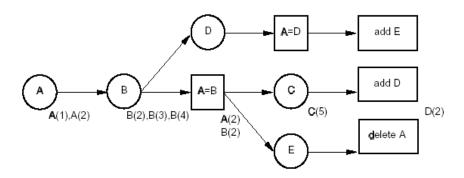
 $A(x) \wedge B(y) \wedge D(x) \Rightarrow add \ E(x)$
 $A(x) \wedge B(x) \wedge E(z) \Rightarrow delete \ A(x)$

- And facts: A(1), A(2), B(2), B(3), B(4), C(5)
- Rete:
 - Compiles the rules to a network that merges conditions of multiple rules together (avoid repeats)
 - Propagates valid unifications
 - Reevaluates only changed conditions

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Rete algorithm. Network.



Rules: $A(x) \wedge B(x) \wedge C(y) \Rightarrow add D(x)$

 $A(x) \wedge B(y) \wedge D(x) \Rightarrow add E(x)$

 $A(x) \wedge B(x) \wedge E(z) \Rightarrow delete A(x)$

Facts: A(1), A(2), B(2), B(3), B(4), C(5)

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Conflict resolution strategies

- **Problem:** Two or more rules are active at the same time. Which one to execute next?
- Solutions:
 - No duplication (do not execute the same rule twice)
 - Recency. Rules referring to facts newly added to the working memory take precedence
 - **Specificity.** Rules that are more specific are preferred.
 - Priority levels. Define priority of rules, actions based on expert opinion. Have multiple priority levels such that the higher priority rules fire first.

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Semantic network systems

- Knowledge about the world described in terms of graphs. Nodes correspond to:
 - Concepts or objects in the domain.

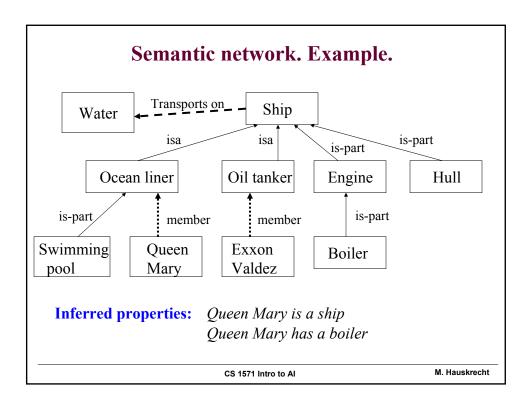
Links to relations. Three kinds:

- Subset links (isa, part-of links)
- Member links (instance links)

Inheritance relation links

- Function links.
- Can be transformed to the first-order logic language
- Graphical representation is often easier to work with
 - better overall view on individual concepts and relations

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Planning: situation calculus

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Representation of actions, situations, events

The world is dynamic:

- What is true now may not be true tomorrow
- Changes in the world may be triggered by our activities

Problems:

- Logic (FOL) as we had it referred to a static world. How to represent the change in the FOL?
- How to represent actions we can use to change the world?

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Planning

Planning problem:

- find a sequence of actions that achieves some goal
- an instance of a search problem
- the state description is typically very complex and relies on a logic-based representation

Methods for modeling and solving a planning problem:

- State space search
- Situation calculus based on FOL
- STRIPS state space search algorithm
- Partial-order planning algorithms

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Situation calculus

Provides a framework for representing change, actions and for reasoning about them

• Situation calculus

- based on the first-order logic,
- a situation variable models new states of the world
- action objects model activities
- uses inference methods developed for FOL to do the reasoning

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Situation calculus

- Logic for reasoning about changes in the state of the world
- The world is described by:
 - Sequences of **situations** of the current state
 - Changes from one situation to another are caused by actions
- The situation calculus allows us to:
 - Describe the initial state and a goal state
 - Build the KB that describes the effect of actions (operators)
 - Prove that the KB and the initial state lead to a goal state
 - extracts a plan as side-effect of the proof

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Situation calculus

The language is based on the First-order logic plus:

- Special variables: s,a objects of type situation and action
- Action functions: return actions.
 - E.g. Move(A, TABLE, B) represents a move action
 - -Move(x,y,z) represents an action schema
- Two special function symbols of type situation
 - $-s_0$ initial situation
 - -DO(a,s) denotes the situation obtained after performing an action a in situation s
- Situation-dependent functions and relations (also called fluents)
 - **Relation:** On(x,y,s) object x is on object y in situation s;
 - **Function:** Above(x,s) object that is above x in situation s.

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Situation calculus. Blocks world example.

	11
	В
A B C	C
Initial state	Goal

 $On(A, Table, s_0)$ $On(B, Table, s_0)$ Find a state (situation) s, such that

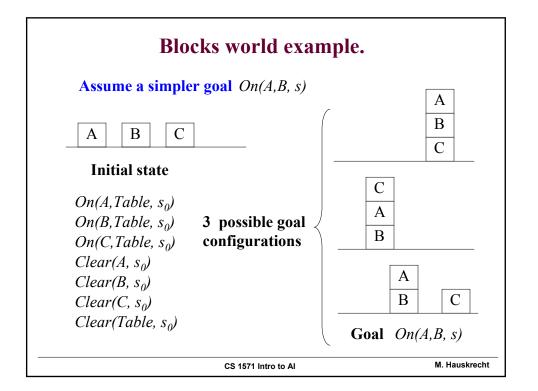
Δ

 $On(C, Table, s_0)$ On(A, B, s) $Clear(A, s_0)$ On(B, C, s) $Clear(B, s_0)$ On(C, Table, s) $Clear(C, s_0)$

Clear(Table, s_0)

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Blocks world example. В В \mathbf{C} C **Initial state** Goal On(A,B,s) $On(A, Table, s_0)$ On(B,C,s) $On(B, Table, s_0)$ On(C, Table, s) $On(C, Table, s_0)$ $Clear(A, s_0)$ **Note:** It is not necessary that Clear(B, s_0) the goal describes all relations $Clear(C, s_0)$ Clear(A, s)Clear(Table, s_0) CS 1571 Intro to Al M. Hauskrecht



Knowledge base: Axioms.

Knowledge base needed to support the reasoning:

• Must represent changes in the world due to actions.

Two types of axioms:

- Effect axioms
 - changes in situations that result from actions
- Frame axioms
 - things preserved from the previous situation

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Blocks world example. Effect axioms.

Effect axioms:

Moving x from y to z. MOVE(x, y, z)

Effect of move changes on On relations

$$On(x, y, s) \land Clear(x, s) \land Clear(z, s) \rightarrow On(x, z, DO(MOVE(x, y, z), s))$$

$$On(x, y, s) \land Clear(x, s) \land Clear(z, s) \rightarrow \neg On(x, y, DO(MOVE(x, y, z), s))$$

Effect of move changes on Clear relations

$$On(x, y, s) \land Clear(x, s) \land Clear(z, s) \rightarrow Clear(y, DO(MOVE(x, y, z), s))$$

$$On(x, y, s) \land Clear(x, s) \land Clear(z, s) \land (z \neq Table)$$

 $\rightarrow \neg Clear(z, DO(MOVE(x, y, z), s))$

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Blocks world example. Frame axioms.

- Frame axioms.
 - Represent things that remain unchanged after an action.

On relations:

$$On(u,v,s) \land (u \neq x) \land (v \neq y) \rightarrow On(u,v,DO(MOVE(x,y,z),s))$$

Clear relations:

$$Clear(u, s) \land (u \neq z) \rightarrow Clear(u, DO(MOVE(x, y, z), s))$$

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Planning in situation calculus

Planning problem:

- find a sequence of actions that lead to a goal
- Planning in situation calculus is converted to the theorem proving problem

Goal state:

$$\exists s \ On(A,B,s) \land On(B,C,s) \land On(C,Table,s)$$

- Possible inference approaches:
 - Inference rule approach
 - Conversion to SAT
- Plan (solution) is a byproduct of theorem proving.
- Example: blocks world

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Planning in a blocks world.

 \mathbf{C} В

Α В \mathbf{C}

Initial state

 $On(A, Table, s_0)$ $On(B, Table, s_0)$ $On(C, Table, s_0)$ $Clear(A, s_0)$ $Clear(B, s_0)$ $Clear(C, s_0)$

Clear(Table, s_0)

On(A,B,s)On(B,C,s)

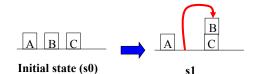
On(C, Table, s)

Goal

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Planning in the blocks world.



 $s_0 =$

 $On(A, Table, s_0)$ Clear (A, s_0) Clear (Table, s_0)

s1

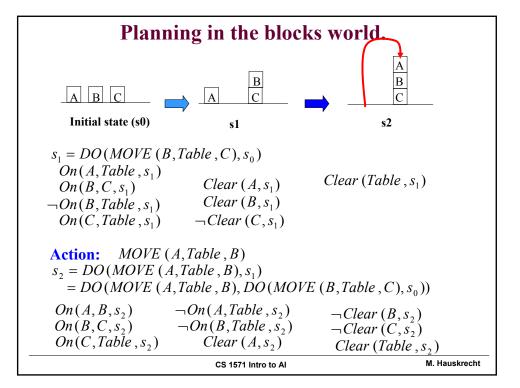
 $On(B, Table, s_0)$ $Clear(B, s_0)$ $On(C, Table, s_0)$ Clear (C, s_0)

Action: MOVE(B, Table, C) $s_1 = DO(MOVE(B, Table, C), s_0)$

 $On(A, Table, s_1)$ Clear (A, s_1) Clear (Table, s_1) $On(B,C,s_1)$

Clear (B, s_1) $\neg On(B, Table, s_1)$ $\neg Clear(C, s_1)$ $On(C, Table, s_1)$

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Planning in situation calculus.

Planning problem:

- Find a sequence of actions that lead to a goal
- Is a special type of a search problem
- Planning in situation calculus is converted to theorem proving.

Problems:

- Large search space
- Large number of axioms to be defined for one action
- Proof may not lead to the best (shortest) plan.

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