

# CS 1571 Introduction to AI

## Lecture 9

### Finding optimal configurations

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### Search for the optimal configuration

#### Objective:

- find the optimal configuration

#### Optimality:

- Defined by some **quality measure**

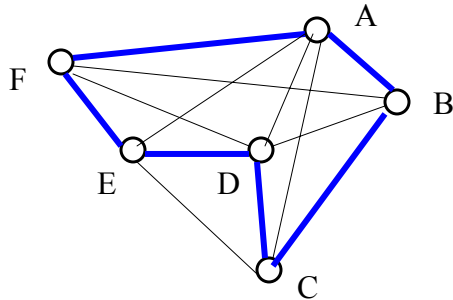
#### Quality measure

- reflects our **preference towards each configuration** (or state)

## Example: Traveling salesman problem

### Problem:

- A graph with distances

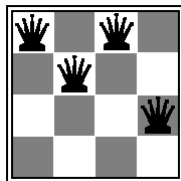


- **Goal:** find the shortest tour which visits every city once and returns to the start

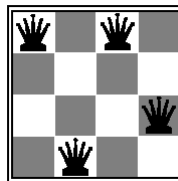
An example of a valid **tour**: ABCDEF

## Example: N queens

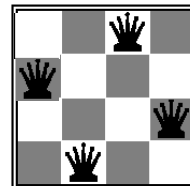
- A CSP problem can be converted to the ‘optimal’ configuration problem
- **The quality of a configuration in a CSP**  
= the number of constraints violated
- **Solving:** minimize the number of constraint violations



# of violations =3



# of violations =1



# of violations =0

## Iterative optimization methods

- Searching systematically for the best configuration with the **DFS** may not be the best solution
- Worst case running time:
  - Exponential in the number of variables
- Solutions to **large ‘optimal’ configuration** problems are often found using iterative optimization methods
- **Methods:**
  - **Hill climbing**
  - **Simulated Annealing**
  - **Genetic algorithms**

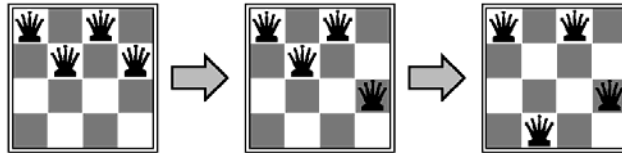
## Iterative optimization methods

### Properties:

- Search **the space of “complete” configurations**
- **Take advantage of local moves**
  - Operators make “local” changes to “complete” configurations
- **Keep track of just one state (the current state)**
  - no memory of past states
  - **!!! No search tree is necessary !!!**

## Example: N-queens

- “Local” operators for generating the next state:
  - Select a variable (a queen)
  - Reallocate its position

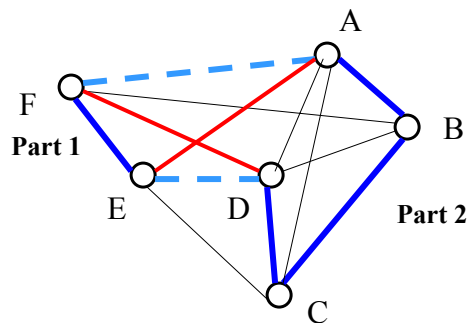


## Example: Traveling salesman problem

- “Local” operator for generating the next state:
  - divide the existing tour into two parts,
  - reconnect the two parts in the opposite order

**Example:**

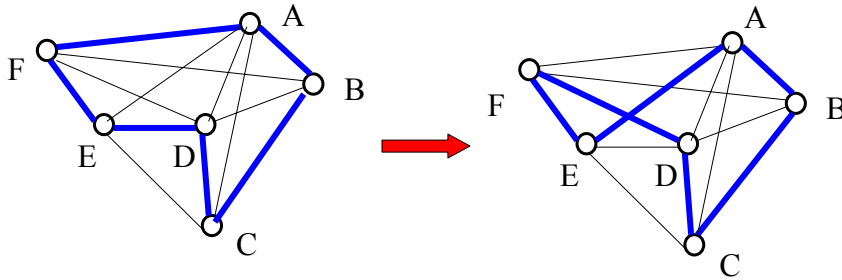
ABCDEF  
↓  
ABCD | EF |  
↓  
ABCDFE



## Example: Traveling salesman problem

### “Local” operator:

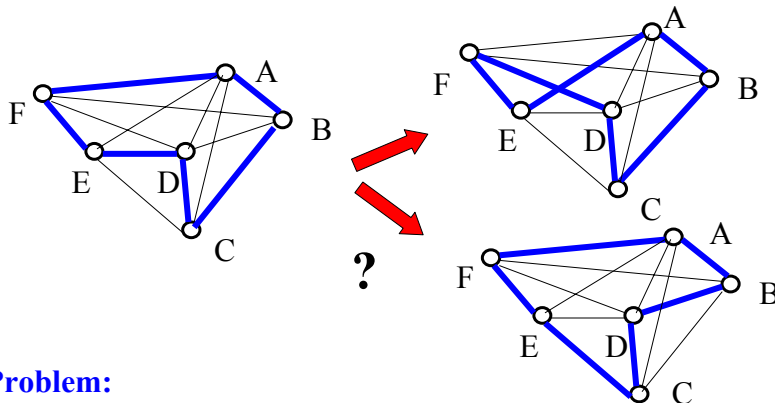
- generates the next configuration (state)



## Searching the configuration space

### Search algorithms

- keep only one configuration (the current configuration)



### Problem:

- How to decide about which operator to apply?

# Search algorithms

Two strategies to choose the configuration (state) to be visited next:

- Hill climbing
- Simulated annealing

- Later: Extensions to multiple current states:

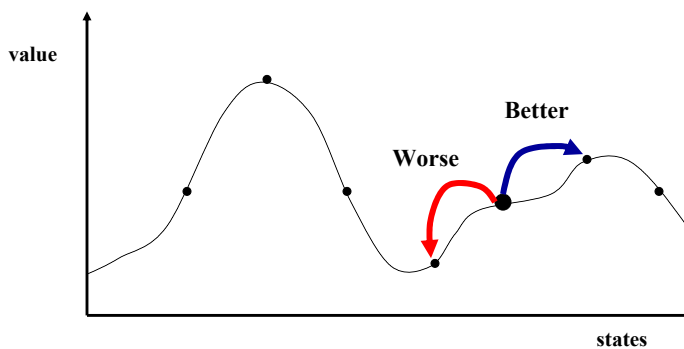
- Genetic algorithms

- **Note:** Maximization is inverse of the minimization

$$\min f(X) \Leftrightarrow \max [-f(X)]$$

## Hill climbing

- Look around at states in the local neighborhood and choose the one with the best value
- Assume: we want to maximize the



## Hill climbing

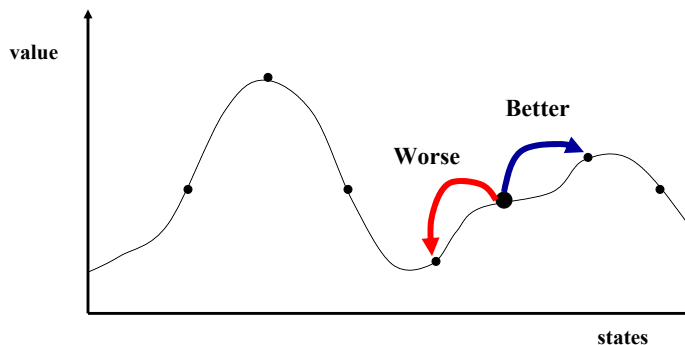
- Always choose the next best successor state
- Stop when no improvement possible

```
function HILL-CLIMBING(problem) returns a solution state
  inputs: problem, a problem
  static: current, a node
         next, a node

  current ← MAKE-NODE(INITIAL-STATE[problem])
  loop do
    next ← a highest-valued successor of current
    if VALUE[next] < VALUE[current] then return current
    current ← next
  end
```

## Hill climbing

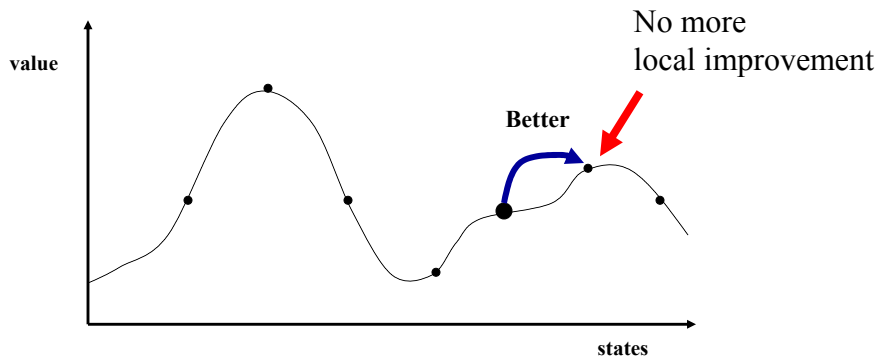
- Look around at states in the local neighborhood and choose the one with the best value



- What can go wrong?

## Hill climbing

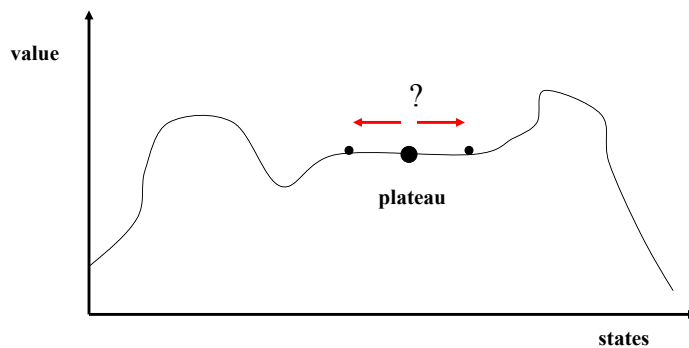
- Hill climbing can get trapped in the local optimum



- What can go wrong?

## Hill climbing

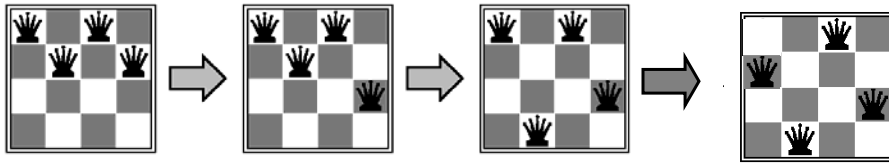
- Hill climbing can get clueless on plateaus





## Hill climbing and n-queens

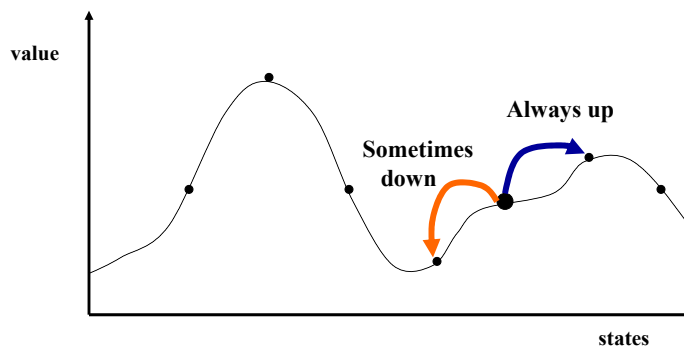
- The quality of a configuration is given by the number of constraints violated
- **Then: Hill climbing** reduces the number of constraints
- **Min-conflict strategy (heuristic):**
  - Choose randomly a variable with conflicts
  - Choose its value such that it violates the fewest constraints



Success !! But not always!!! The local optima problem!!!

## Simulated annealing

- Permits “bad” moves to states with lower value, thus escape the local optima
- **Gradually decreases** the frequency of such moves and their size (parameter controlling it – **temperature**)



## Simulated annealing algorithm

The probability of making a move:

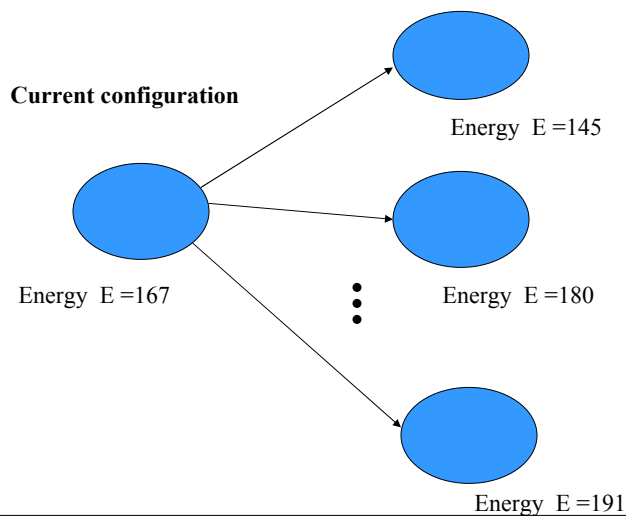
- The probability of moving into a state with a higher value is 1
- The probability of moving into a state with a lower value is

$$p(\text{Accept } NEXT) = e^{\Delta E / T} \quad \text{where} \quad \Delta E = E_{NEXT} - E_{CURRENT}$$

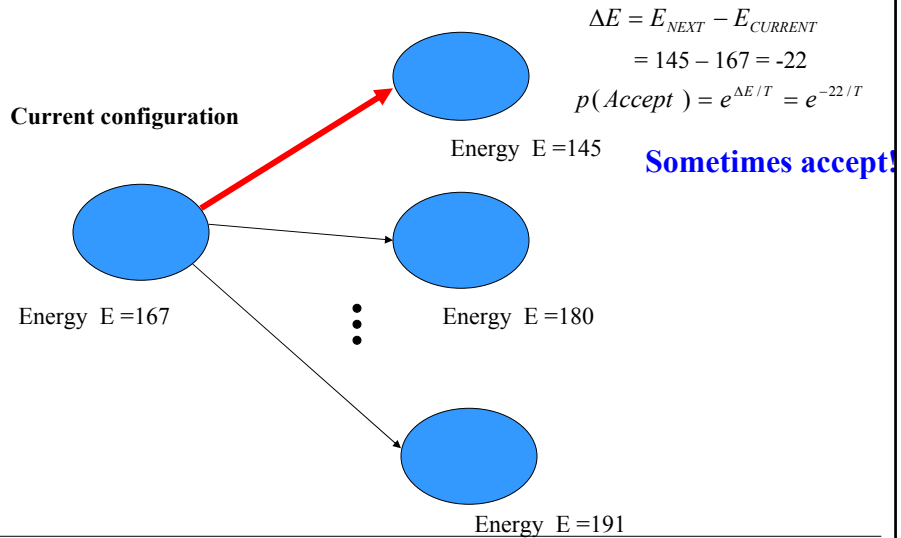
– The probability is:

- **Proportional to the energy difference**

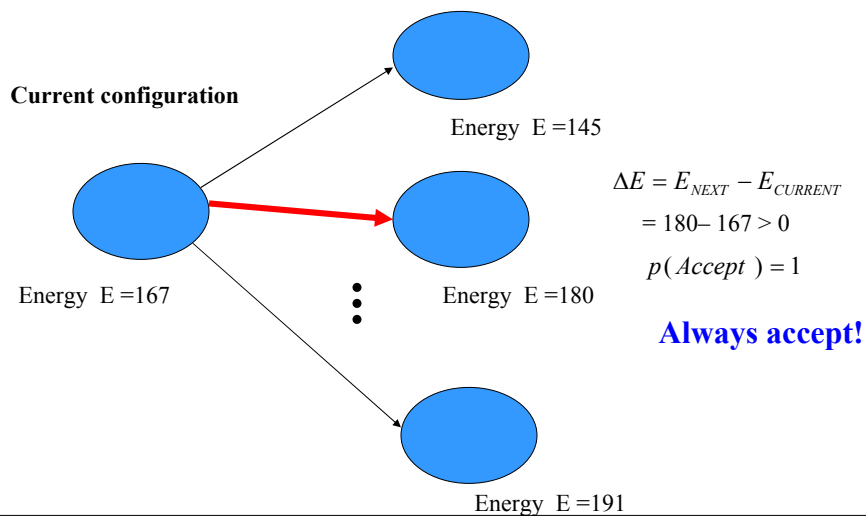
## Simulated annealing algorithm



## Simulated annealing algorithm



## Simulated annealing algorithm



## Simulated annealing algorithm

The probability of moving into a state with a lower value is

$$p(\text{Accept}) = e^{\Delta E / T} \quad \text{where} \quad \Delta E = E_{\text{NEXT}} - E_{\text{CURRENT}}$$

The probability is:

- **Modulated through a temperature parameter T:**
  - for  $T \rightarrow \infty$  the probability of any move approaches 1
  - for  $T \rightarrow 0$  the probability that a state with smaller value is selected goes down and approaches 0
- **Cooling schedule:**
  - Schedule of changes of a parameter T over iteration steps

## Simulated annealing

```
function SIMULATED-ANNEALING(problem, schedule) returns a solution state
  inputs: problem, a problem
           schedule, a mapping from time to “temperature”
  static: current, a node
           next, a node
           T, a “temperature” controlling the probability of downward steps

  current ← MAKE-NODE(INITIAL-STATE[problem])
  for t ← 1 to ∞ do
    T ← schedule[t]
    if T=0 then return current
    next ← a randomly selected successor of current
     $\Delta E \leftarrow \text{VALUE}[\textit{next}] - \text{VALUE}[\textit{current}]$ 
    if  $\Delta E > 0$  then current ← next
    else current ← next only with probability  $e^{\Delta E / T}$ 
```

## Simulated annealing algorithm

- **Simulated annealing algorithm**
  - developed originally for modeling physical processes (Metropolis et al, 53)
- **Properties:**
  - **If T is decreased slowly enough the best configuration (state) is always reached**
- **Applications:**
  - VLSI design
  - airline scheduling

## Simulated evolution and genetic algorithms

- Limitations of **simulated annealing**:
  - Pursues one state configuration at the time;
  - Changes to configurations are typically local

### Can we do better?

- Assume we have two configurations with good values that are quite different
- We expect that the combination of the two individual configurations may lead to a configuration with higher value (**Not guaranteed !!!**)

This is the idea behind **genetic algorithms** in which we grow a population of individual combinations

# Genetic algorithms

## Algorithm idea:

- **Create a population of random configurations**
  - **Create a new population through:**
    - Biased selection of pairs of configurations from the previous population
    - Crossover (combination) of pairs
    - Mutation of resulting individuals
  - **Evolve the population over multiple generation cycles**
- 
- **Selection of configurations to be combined:**
    - **Fitness function = value function**  
measures the quality of an individual (a state) in the population

## Reproduction process in GA

- Assume that a state configuration is defined by a set variables with two values, represented as 0 or 1

