#### CS 1571 Introduction to AI Lecture 4

#### Uninformed search methods I.

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#### Problem-solving as search

- Many search problems can be converted to graph search problems
- A graph search problem can be described in terms of:
  - A set of states representing different world situations
  - Initial state
  - Goal condition
  - Operators defining valid moves between states
- Two types of search:
  - Path search: solution is a path to a goal state
  - Configuration search: solution is a state satisfying the goal condition
- **Optimal solution** = a solution with the optimal value
  - shortest path between the two cities, or
  - a desired n-queen configuration

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## Formulating a search problem

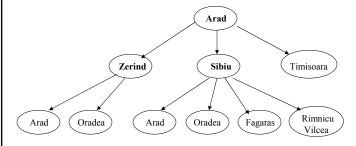
- Search (process)
  - The process of exploration of the search space
- The efficiency of the search depends on:
  - The search space and its size
  - Method used to explore (traverse) the search space
  - Condition to test the satisfaction of the search objective
    (what it takes to determine I found the desired goal object)
- Think twice before solving the problem by search:
  - Choose the search space and the exploration policy

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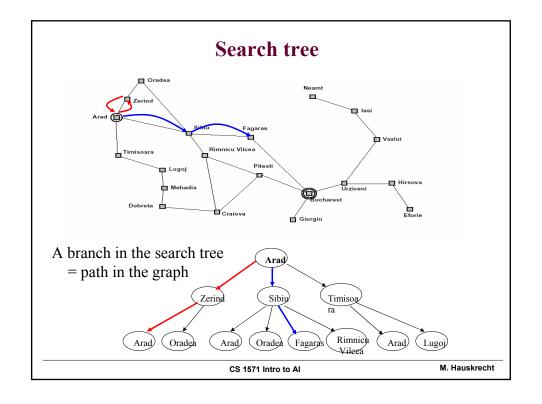
#### Search process

- Exploration of the state space through successive application of operators from the initial state
- A search tree = a kind of (search) exploration trace, branches corresponding to explored paths, and leaf nodes corresponding to exploration fringe



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### Search tree A search tree = a (search) exploration trace - It is different from the graph defining the problem - States can repeat in the search tree Graph **Search tree** Zerind Sibiu Timisoara Rimnicu Arad Oradea) Oradea Fagaras Vilcea CS 1571 Intro to Al M. Hauskrecht



**General-search** (*problem*, *strategy*) **initialize** the search tree with the initial state of *problem* **loop** 

if there are no candidate states to explore return failure choose a leaf node of the tree to expand next according to *strategy* if the node satisfies the goal condition return the solution expand the node and add all of its successors to the tree end loop

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#### General search algorithm

**General-search** (*problem*, *strategy*) **initialize** the search tree with the initial state of *problem* **loop** 

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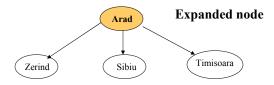
Arad

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**General-search** (problem, strategy)

initialize the search tree with the initial state of problem loop

if there are no candidate states to explore return failure **choose** a leaf node of the tree to expand next according to *strategy* if the node satisfies the goal condition return the solution expand the node and add all of its successors to the tree end loop



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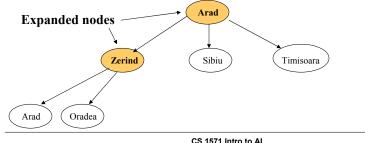
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#### General search algorithm

**General-search** (*problem*, *strategy*)

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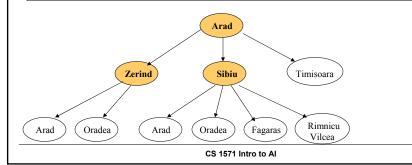


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**General-search** (problem, strategy)

**initialize** the search tree with the initial state of *problem* **loop** 

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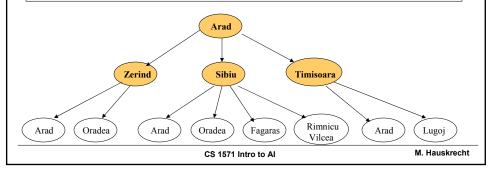
# General search algorithm

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**General-search** (*problem*, *strategy*)

**initialize** the search tree with the initial state of *problem* **loop** 

if there are no candidate states to explore return failure choose a leaf node of the tree to expand next according to *strategy* if the node satisfies the goal condition return the solution expand the node and add all of its successors to the tree end loop



**General-search** (*problem*, *strategy*)

initialize the search tree with the initial state of problem

if there are no candidate states to explore return failure **choose** a leaf node of the tree to expand next according to *strategy* if the node satisfies the goal condition return the solution expand the node and add all of its successors to the tree end loop

Search methods differ in how they explore the space, that is how they choose the node to expand next !!!!!

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#### Implementation of search

Search methods can be implemented using queue structure

General search (problem, Queuing-fn)

 $nodes \leftarrow Make-queue(Make-node(Initial-state(problem)))$ 

if nodes is empty then return failure

 $node \leftarrow Remove-node(nodes)$ 

if Goal-test(problem) applied to State(node) is satisfied then return node

nodes ← Queuing-fn(nodes, Expand(node, Operators(node)))

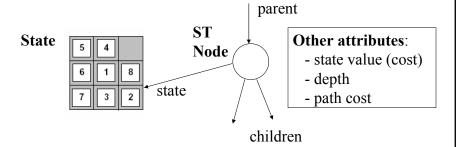
end loop

Candidates are added to *nodes* representing the queue structure

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## Implementation of search

• A **search tree node** is a data-structure constituting part of a search tree



• Expand function – applies Operators to the state represented by the search tree node. Together with Queuing-fn it fills the attributes.

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#### **Uninformed search methods**

- rely only on the information available in the problem definition
  - Breadth first search
  - Depth first search
  - Iterative deepening
  - Bi-directional search

#### For the minimum cost path problem:

- Uniform cost search

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#### Search methods

#### **Properties of search methods:**

- Completeness.
  - Does the method find the solution if it exists?
- Optimality.
  - Is the solution returned by the algorithm optimal? Does it give a minimum length path?
- Space and time complexity.
  - How much time it takes to find the solution?
  - How much memory is needed to do this?

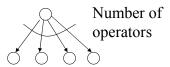
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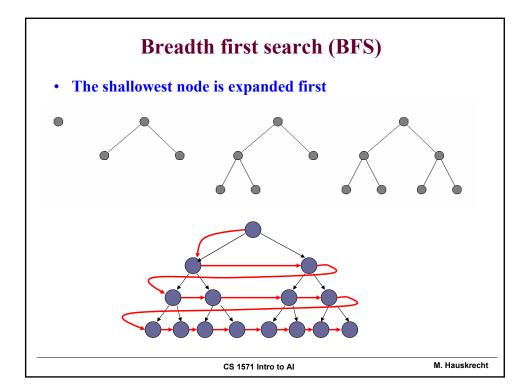
#### Parameters to measure complexities.

- Space and time complexity.
  - Complexities are measured in terms of parameters:
    - b maximum branching factor
    - d depth of the optimal solution
    - m maximum depth of the state space

#### **Branching factor**

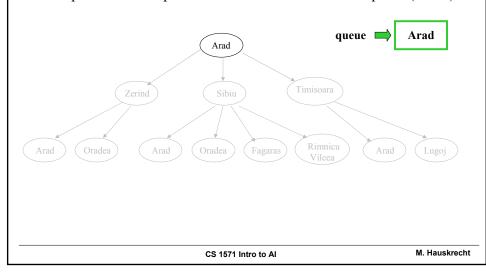


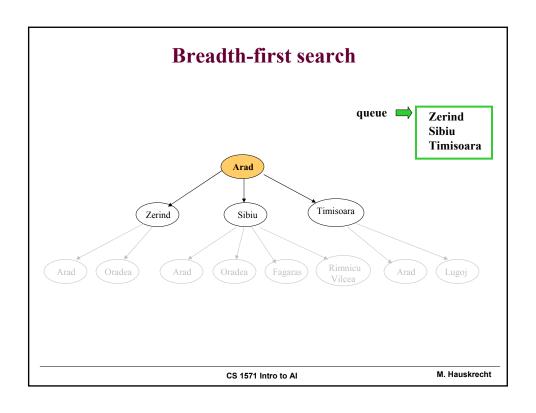
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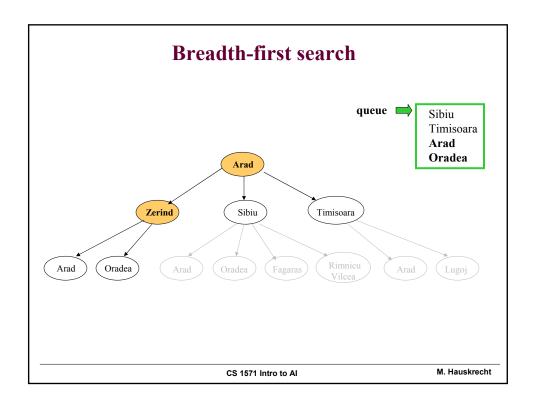


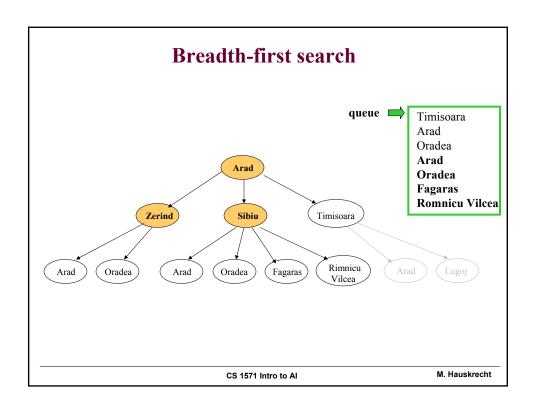
#### **Breadth-first search**

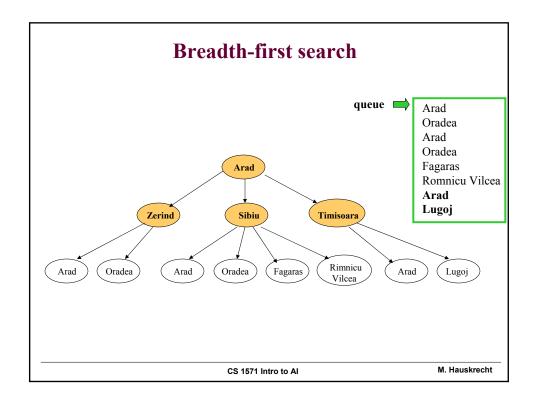
- Expand the shallowest node first
- Implementation: put successors to the end of the queue (FIFO)











- Completeness: ?
- Optimality: ?
- Time complexity: ?
- Memory (space) complexity: ?
  - For complexities use:
    - b maximum branching factor
    - d depth of the optimal solution
    - m maximum depth of the search tree

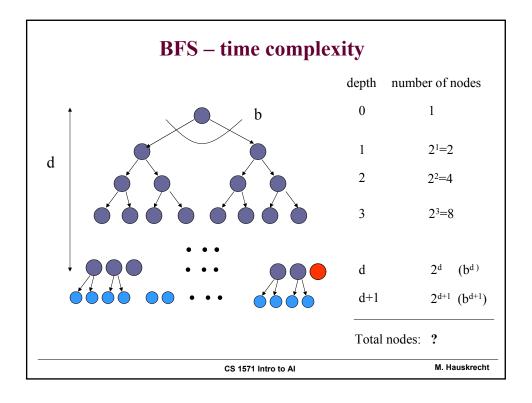
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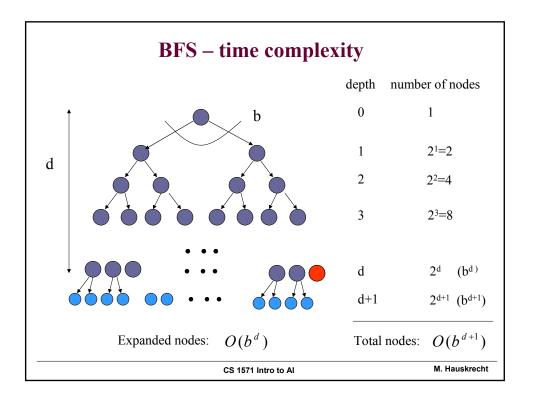
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## Properties of breadth-first search

- Completeness: Yes. The solution is reached if it exists.
- Optimality: Yes, for the shortest path.
- Time complexity: ?
- Memory (space) complexity: ?

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- Completeness: Yes. The solution is reached if it exists.
- Optimality: Yes, for the shortest path.
- Time complexity:

$$1 + b + b^2 + \dots + b^d = O(b^d)$$

exponential in the depth of the solution d

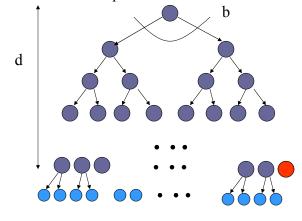
• Memory (space) complexity: ?

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# **BFS** – memory complexity

• Count nodes kept in the tree structure or in the queue



depth number of nodes

0 1

 $1 2^{1}=2$ 

 $2 2^2 = 4$ 

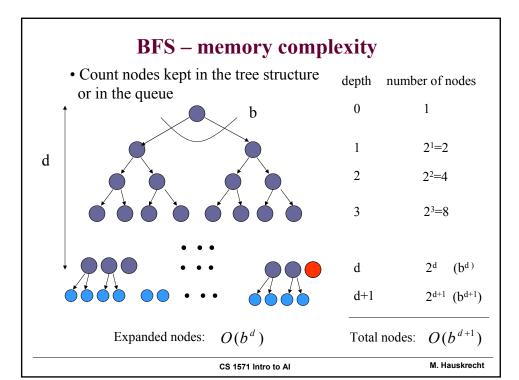
 $3 2^3 = 8$ 

d 2<sup>d</sup> (b<sup>d</sup>)

d+1  $2^{d+1}$   $(b^{d+1})$ 

Total nodes: ?

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- Completeness: Yes. The solution is reached if it exists.
- Optimality: Yes, for the shortest path.
- Time complexity:

$$1 + b + b^2 + \dots + b^d = O(b^d)$$

exponential in the depth of the solution d

• Memory (space) complexity:

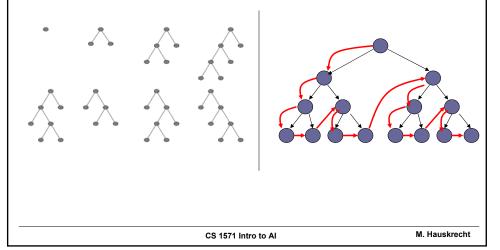
$$O(b^d)$$

nodes are kept in the memory

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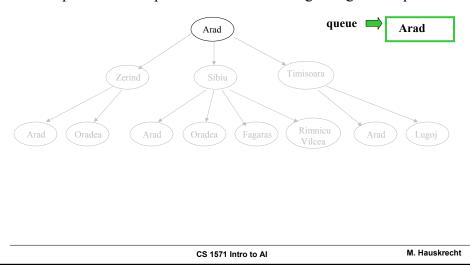
## **Depth-first search (DFS)**

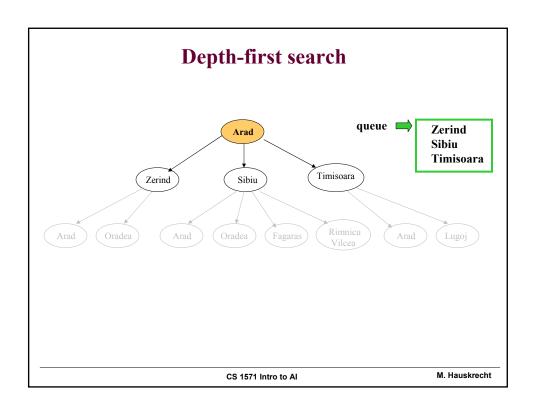
- The deepest node is expanded first
- Backtrack when the path cannot be further expanded

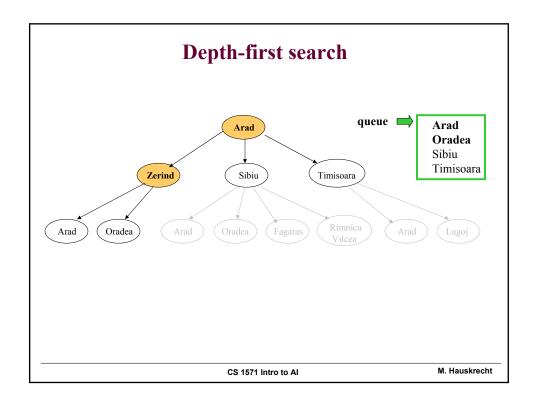


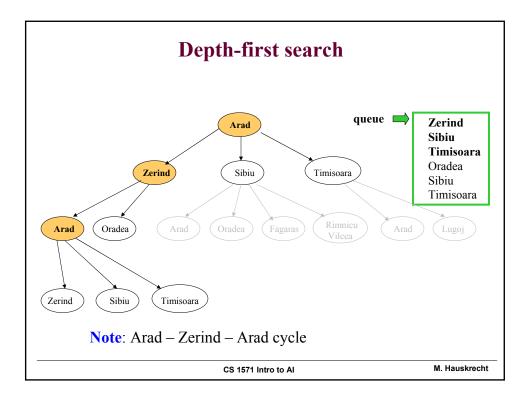
# **Depth-first search**

- The deepest node is expanded first
- Implementation: put successors to the beginning of the queue









- Completeness: Does it always find the solution if it exists?
- Optimality: ?
- Time complexity: ?
- Memory (space) complexity: ?

- Completeness: No. Infinite loops can occur.
  Infinite loops imply -> Infinite depth search tree.
- Optimality: does it find the minimum length path?
- Time complexity: ?
- Memory (space) complexity: ?

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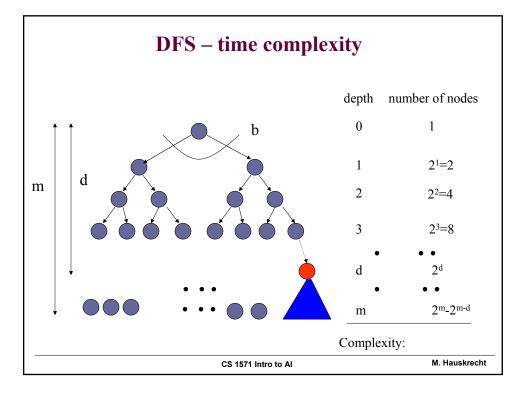
#### Properties of depth-first search

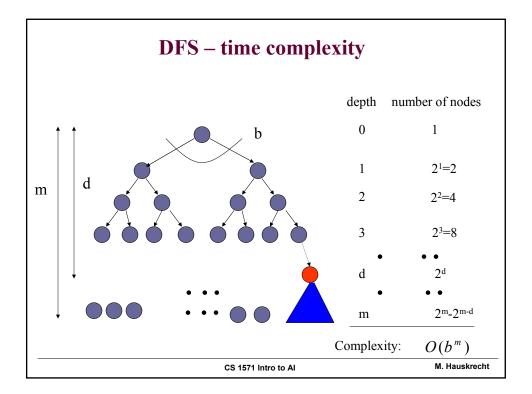
- Completeness: No. Infinite loops can occur.
- **Optimality:** No. Solution found first may not be the shortest possible.
- Time complexity: ?
- Memory (space) complexity: ?

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- Memory (space) complexity: ?

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- Completeness: No. Infinite loops can occur.
- **Optimality:** No. Solution found first may not be the shortest possible.
- Time complexity:

$$O(b^m)$$

exponential in the maximum depth of the search tree m

• Memory (space) complexity: ?

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- Completeness: No. Infinite loops can occur.
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- Time complexity:

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exponential in the maximum depth of the search tree m

• Memory (space) complexity: ?

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# **DFS** – memory complexity

depth number of nodes kept

b

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# DFS – memory complexity

depth number of nodes kept



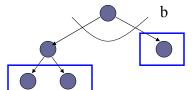
0

2 = b

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# DFS – memory complexity



depth number of nodes kept

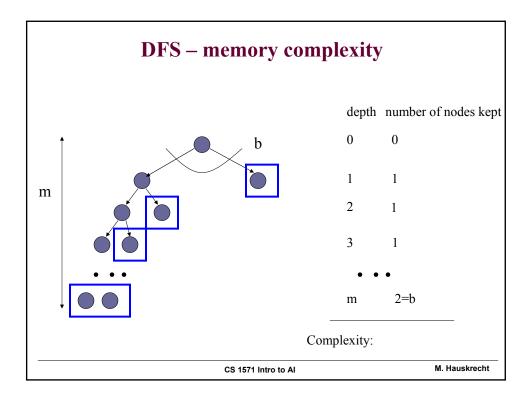
0

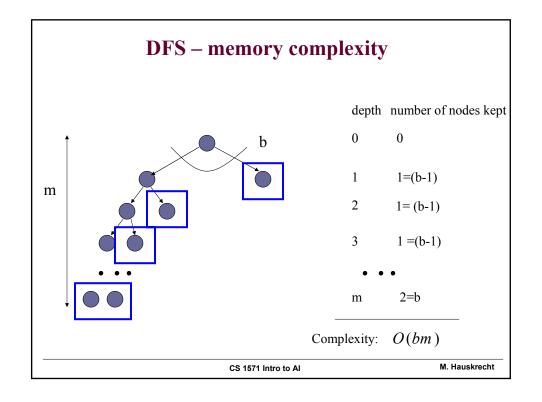
$$1 = (b-1)$$

0

$$2 2 = b$$

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- Completeness: No. Infinite loops can occur.
- **Optimality:** No. Solution found first may not be the shortest possible.
- Time complexity:

$$O(b^m)$$

exponential in the maximum depth of the search tree m

Memory (space) complexity:

O(bm)

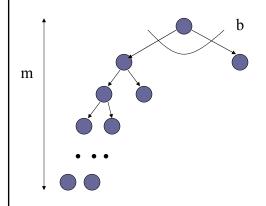
linear in the maximum depth of the search tree m

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# **DFS** – memory complexity

Count nodes kept in the tree structure or the queue



depth number of nodes

0

1 2 = b

1

2 2

3 2

m 2

Total nodes: O(bm)

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- Completeness: No. Infinite loops can occur.
- **Optimality:** No. Solution found first may not be the shortest possible.
- Time complexity:

 $O(b^m)$ 

exponential in the maximum depth of the search tree m

Memory (space) complexity:

O(bm)

the tree size we need to keep is linear in the maximum depth of the search tree *m* 

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