

# CS 1571 Introduction to AI

## Lecture 6

### Constraint-satisfaction search.

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### Search problem

**A search problem:**

- **Search space (or state space):** a set of objects among which we conduct the search;
- **Initial state:** an object we start to search from;
- **Operators (actions):** transform one state in the search space to the other;
- **Goal condition:** describes the object we search for
- **Possible metric on a search space:**
  - measures the quality of the object with regard to the goal

Search problems occur in planning, optimizations, learning

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## Constraint satisfaction problem (CSP)

Two types of search:

- **path search** (a path from the initial state to a state satisfying the goal condition)
- **configuration search** (a configuration satisfying the goal)

**Constraint satisfaction problem (CSP)** is a **configuration search problem** where:

- A state is defined in terms of *variables* and their *value assignments*
- Goal condition is represented *by a set constraints* on possible variable values
- **Search:** find *variable values satisfying the constraints*

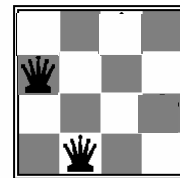
Special properties of the CSP allow more specific procedures to be designed and applied for solving them

## Example of a CSP: N-queens

**Goal:** n queens placed in non-attacking positions on the board

**Variables:**

- Represent queens, one for each column:
  - $Q_1, Q_2, Q_3, Q_4$
- Values:
  - Row placement of each queen on the board  
 $\{1, 2, 3, 4\}$



$$Q_1 = 2, Q_2 = 4$$

**Constraints:**  $Q_i \neq Q_j$  Two queens not in the same row

$|Q_i - Q_j| \neq |i - j|$  Two queens not on the same diagonal

## Satisfiability (SAT) problem

Determine whether a sentence in the conjunctive normal form (CNF) is satisfiable (can evaluate to true)

- Used in the propositional logic (covered later)

$$(P \vee Q \vee \neg R) \wedge (\neg P \vee \neg R \vee S) \wedge (\neg P \vee Q \vee \neg T) \dots$$

### Variables:

- Propositional symbols (P, R, T, S)
- Values: *True*, *False*

### Constraints:

- Every conjunct must evaluate to true, at least one of the literals must evaluate to true

$$(P \vee Q \vee \neg R) \equiv \text{True}, (\neg P \vee \neg R \vee S) \equiv \text{True}, \dots$$

## Other real world CSP problems

### Scheduling problems:

- E.g. telescope scheduling
- High-school class schedule

### Design problems:

- Hardware configurations
- VLSI design

### More complex problems may involve:

- **real-valued variables**
- **additional preferences on variable assignments** – the optimal configuration is sought

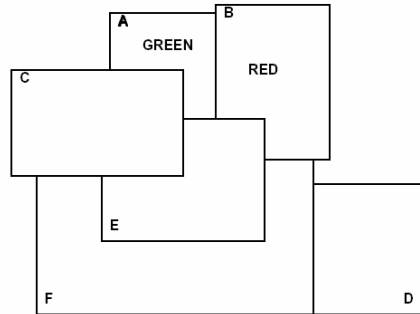
## Map coloring

Color a map using  $k$  different colors such that no adjacent countries have the same color

**Variables: ?**

- Variable values: ?

**Constraints: ?**



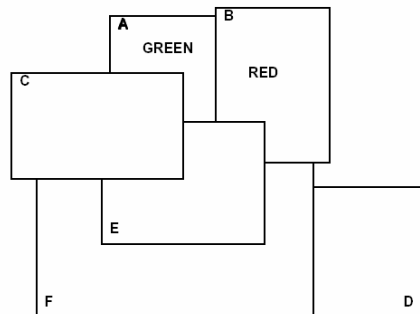
## Map coloring

Color a map using  $k$  different colors such that no adjacent countries have the same color

**Variables:**

- Represent countries
  - $A, B, C, D, E$
- Values:
  - $K$  -different colors
  - $\{\text{Red, Blue, Green, ...}\}$

**Constraints: ?**

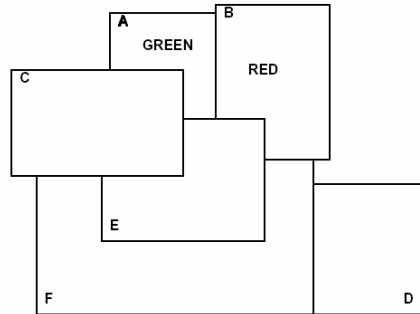


## Map coloring

Color a map using  $k$  different colors such that no adjacent countries have the same color

### Variables:

- Represent countries
  - $A, B, C, D, E$
- Values:
  - $K$ -different colors
  - $\{\text{Red, Blue, Green, ...}\}$



**Constraints:**  $A \neq B, A \neq C, C \neq E$ , etc

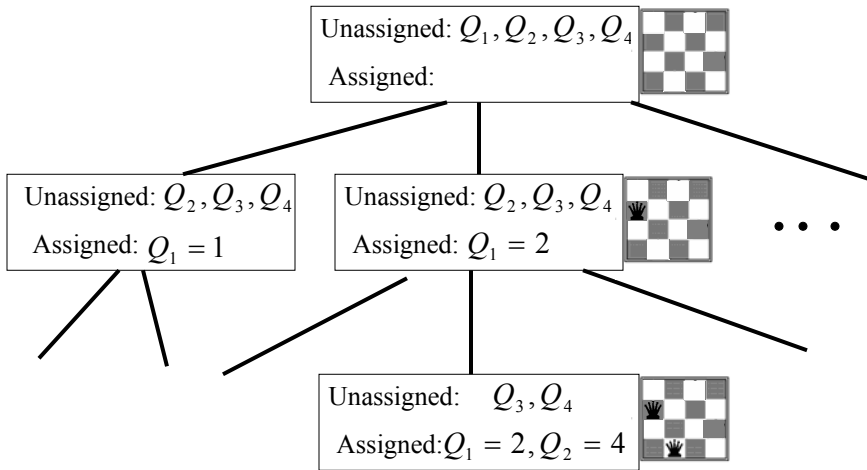
An example of a problem with **binary constraints**

## Constraint satisfaction as a search problem

### Formulation of a CSP as a search problem:

- **States.** Assignment (partial, complete) of values to variables.
- **Initial state.** No variable is assigned a value.
- **Operators.** Assign a value to one of the unassigned variables.
- **Goal condition.** All variables are assigned, no constraints are violated.
- **Constraints** can be **represented**:
  - **Explicitly** by a set of allowable values
  - **Implicitly** by a function that tests for the satisfaction of constraints

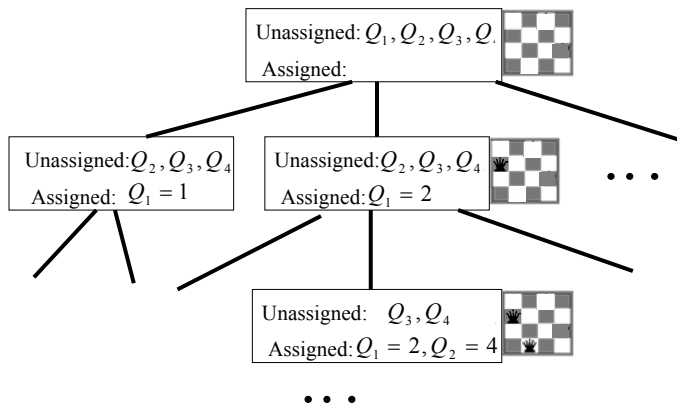
## Solving CSP as a standard search



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## Solving a CSP through standard search

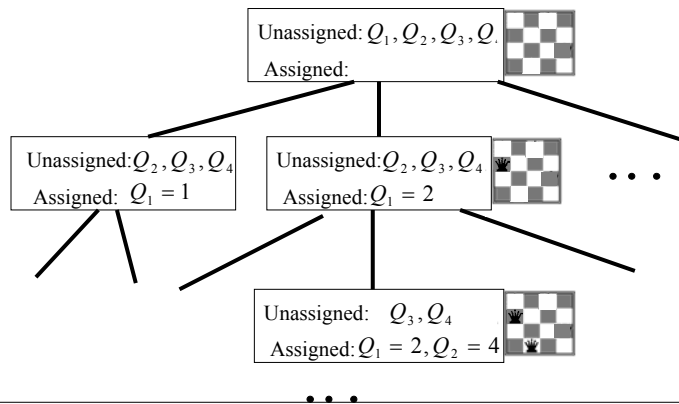
- **Maximum depth of the tree (m): ?**
- **Depth of the solution (d) : ?**
- **Branching factor (b) : ?**



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## Solving a CSP through standard search

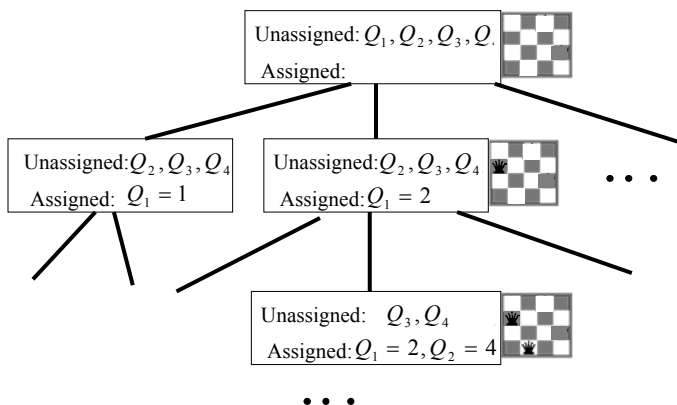
- **Maximum depth of the tree:** Number of variables of the CSP
- **Depth of the solution:** Number of variables of the CSP
- **Branching factor:** if we fix the order of variable assignments the branch factor depends on the number of their values



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## Solving a CSP through standard search

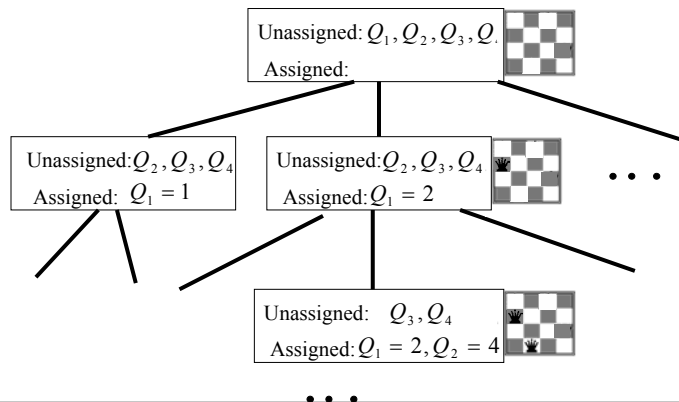
- **What search algorithm to use: ?**  
Depth of the tree = Depth of the solution = number of vars



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## Solving a CSP through standard search

- What search algorithm to use: **Depth first search !!!**
  - Since we know the depth of the solution
  - We do not have to keep large number of nodes in queues



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## Backtracking

**Depth-first search for CSP** is also referred to as **backtracking**

The violation of constraints needs to be checked for each node, either during its generation or before its expansion

### Consistency of constraints:

- Current **variable assignments** together **with constraints** **restrict remaining legal values of unassigned variables**;
- The remaining **legal and illegal values of variables** may be **inferred** (effect of constraints propagates)
- To prevent “blind” exploration it is necessary to keep track of the remaining legal values, so we know when the constraints are violated and when to terminate the search

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## Constraint propagation

A **state** (more broadly) is defined by a set variables and their legal and illegal assignments

Legal and illegal assignments can be represented through variable **equations** and variable **disequations**

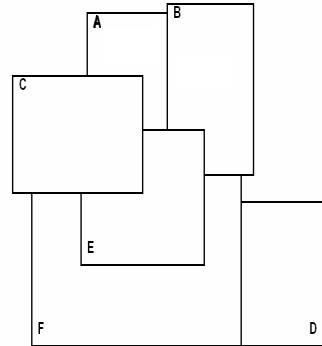
Example: **map coloring**

Equation  $A = \text{Red}$

Disequation  $C \neq \text{Red}$

**Constraints + assignments**  
can entail new equations and disequations

$A = \text{Red} \rightarrow B \neq \text{Red}$



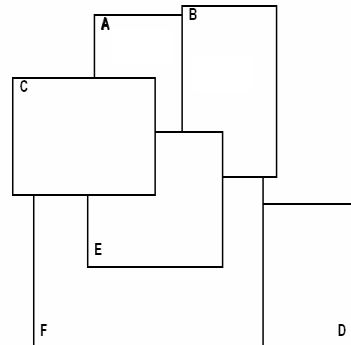
## Constraint propagation

- Assign  $A = \text{Red}$

	Red	Blue	Green
A	✓		
B			
C			
D			
E			
F			



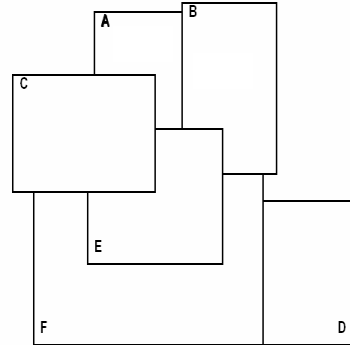
✓ - equations    ✗ - disequations



## Constraint propagation

- Assign A=Red

	Red	Blue	Green
A	✓		
B	✗		
C	✗		
D			
E	✗		
F			

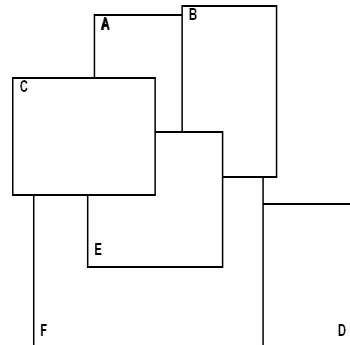


✓ - equations ✗ - disequations

## Constraint propagation

- Assign E=Blue

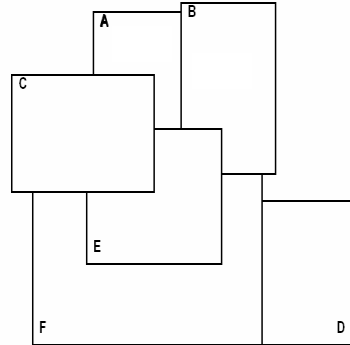
	Red	Blue	Green
A	✓		
B	✗		
C	✗		
D			
E	✗	✓	
F			



## Constraint propagation

- Assign E=Blue

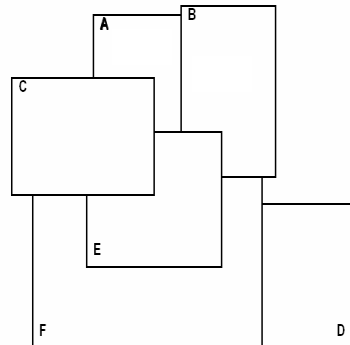
	Red	Blue	Green
A	✓	✗	
B	✗	✗	
C	✗	✗	
D			
E	✗	✓	
F		✗	



## Constraint propagation

- Assign F=Green

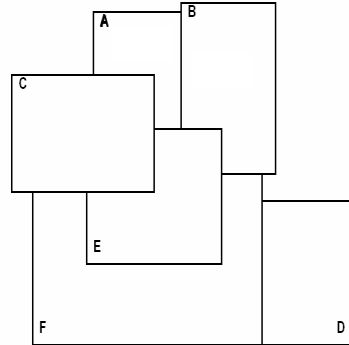
	Red	Blue	Green
A	✓	✗	
B	✗	✗	
C	✗	✗	
D			
E	✗	✓	
F		✗	✓



## Constraint propagation

- Assign F=Green

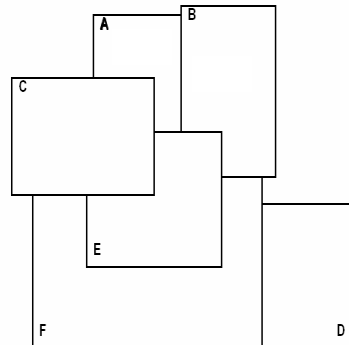
	Red	Blue	Green
A	✓	✗	
B	✗	✗	✗
C	✗	✗	✗
D			✗
E	✗	✓	✗
F		✗	✓



## Constraint propagation

- Assign F=Green

	Red	Blue	Green
A	✓	✗	
B	✗	✗	✗
C	✗	✗	✗
D			✗
E	✗	✓	✗
F		✗	✓

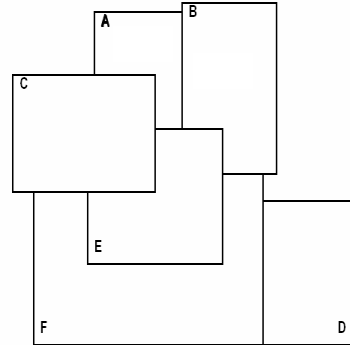


**Conflict !!! No legal assignments available for B and C**

## Constraint propagation

- We can derive remaining legal values through propagation

	Red	Blue	Green
A	✓	✗	
B	✗	✗	✓
C	✗	✗	✓
D			
E	✗	✓	
F		✗	



B=Green

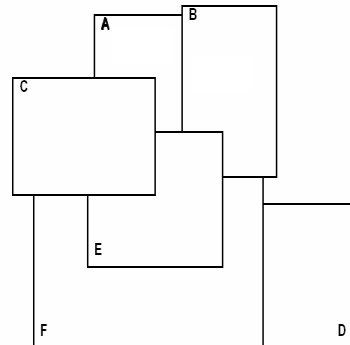
C=Green

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## Constraint propagation

- We can derive remaining legal values through propagation

	Red	Blue	Green
A	✓	✗	✗
B	✗	✗	✓
C	✗	✗	✓
D	✗		
E	✗	✓	✗
F	✓	✗	✗



B=Green

C=Green



F=Red

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## Constraint propagation

Three known techniques for propagating the effects of past assignments and constraints:

- **Value propagation**
- **Arc consistency**
- **Forward checking**
- **Difference:**
  - Completeness of inferences
  - Time complexity of inferences.

## Constraint propagation

1. **Value propagation. Infers:**
    - **equations from** the set of **equations** defining the partial assignment, **and a constraint**
  2. **Arc consistency. Infers:**
    - **disequations from** the set of **equations and disequations** defining the partial assignment, and **a constraint**
    - **equations through** the **exhaustion of alternatives**
  3. **Forward checking. Infers:**
    - **disequations from** a set of **equations** defining the partial assignment, and a constraint
    - **Equations through** the **exhaustion of alternatives**
- Restricted forward checking:**
- uses only active constraints (active constraint – only one variable unassigned in the constraint)

## Heuristics for CSPs

**Backtracking** searches the space in the depth-first manner.

But we can choose:

- Which variable to assign next?
- Which value to choose first?

### Heuristics

- **Most constrained variable**
  - Which variable is likely to become a bottleneck?
- **Least constraining value**
  - Which value gives us more flexibility later?

## Heuristics for CSP

Examples: **map coloring**

### Heuristics

- **Most constrained variable**
  - Country E is the most constrained one (cannot use Red, Green)
- **Least constraining value**
  - Assume we have chosen variable C
  - Red is the least constraining valid color for the future

