



CS 1501: Algorithm Implementation

LZW Data Compression

Data Compression

- Reduce the size of data
 - **Reduces storage space** and hence the **cost**
 - Reduces time to retrieve and transmit data

Compression ratio = original data size / compressed data size

Lossless and Lossy Compression

- $\text{compressedData} = \text{compress}(\text{originalData})$
- $\text{decompressedData} = \text{decompress}(\text{compressedData})$
- **Lossless** compression $\rightarrow \text{originalData} = \text{decompressedData}$
- **Lossy** compression $\rightarrow \text{originalData} \neq \text{decompressedData}$

Lossless and Lossy Compression

- **Lossy compressors** generally obtain much higher compression ratios as compared to **lossless compressors**
e.g. 100 vs. 2
- **Lossless compression** is essential in applications such as **text** file compression.
- **Lossy compression** is acceptable in many **audio** and **imaging** applications
 - In video transmission, a slight loss in the transmitted video is not noticable by the human eye.

Text Compression

- Lossless compression is essential
- Popular text compressors such as **zip** and Unix's **compress** are based on the **LZW** (Lempel-Ziv-Welch) method.

LZW Compression

- Character sequences in the original text are replaced by codes that are **dynamically** determined.
- The **code table** is not encoded into the compressed text, because it may be reconstructed from the compressed text during decompression.

LZW Compression

- Assume the letters in the text are limited to {a, b}
 - In practice, the alphabet may be the 256 character ASCII set.
- The characters in the alphabet are assigned code numbers beginning at 0
- The initial code table is:

code	0	1
key	a	b

Compression

- `public static void compress() {`
- `String input = BinaryStdIn.readString();`
- `TST<Integer> st = new TST<Integer>();`
- `for (int i = 0; i < R; i++)`
- `st.put("" + (char) i, i);`
- `int code = R+1; // R is codeword for EOF`
- `while (input.length() > 0) {`
- `String s = st.longestPrefixOf(input); // Find max prefix match s.`
- `BinaryStdOut.write(st.get(s), W); // Print s's encoding.`
- `int t = s.length();`
- `if (t < input.length() && code < L) // Add s to symbol table.`
- `st.put(input.substring(0, t + 1), code++);`
- `input = input.substring(t); // Scan past s in input.`
- `}`
- `BinaryStdOut.write(R, W);`
- `BinaryStdOut.close();`
- `}`

LZW Compression

code	0	1
key	a	b

- Original text = **abababbabaabbabbaabba**
- Compression is done by scanning the original text from left to right.
- Find longest prefix **p** for which there is a code in the code table.
- Represent **p** by its code **pCode** and assign the next available code number to **pc**, where **c** is the next character in the text to be compressed.

LZW Compression

code	0	1	2
key	a	b	ab

- Original text = **abababbabaabbabbaabba**
- $p = a$ $pCode = 0$ Compressed text = 0
- $c = b$
- Enter $p \mid c = ab$ into the code table.

code	0	1	2	3
key	a	b	ab	ba

- Original text = abababbabaabbabbaabba
- Compressed text = 0

- $p = b$ $pCode = 1$ Compressed text = 01
- $c = a$
- Enter $p | c = ba$ into the code table.

code	0	1	2	3	4
key	a	b	ab	ba	aba

- Original text = abababbabaabbabbaabba
- Compressed text = 01

- $p = ab$ $pCode = 2$ Compressed text = 012
- $c = a$
- Enter aba into the code table

code	0	1	2	3	4	5
key	a	b	ab	ba	aba	abb

- Original text = abababbabaabbabbaabba
- Compressed text = 012

- $p = ab$ $pCode = 2$ Compressed text = 0122
- $c = b$
- Enter **abb** into the code table.

code	0	1	2	3	4	5	6
key	a	b	ab	ba	aba	abb	bab

- Original text = ababab**b**abaabbabbaabba
- Compressed text = 0122

- $p = ba$ $pCode = 3$ Compressed text = 01223
- $c = b$
- Enter **bab** into the code table.

code	0	1	2	3	4	5	6	7
key	a	b	ab	ba	aba	abb	bab	baa

- Original text = abababbabaabbabbaabba
- Compressed text = 01223

- $p = ba$ $pCode = 3$ Compressed text = 012233
- $c = a$
- Enter baa into the code table.

code	0	1	2	3	4	5	6	7	8
key	a	b	ab	ba	aba	abb	bab	baa	abba

- Original text = abababbaba**abbabbaabba**
- Compressed text = 012233

- **p = abb** **pCode = 5** Compressed text = 0122335
- **c = a**
- Enter **abba** into the code table.

code	0	1	2	3	4	5	6	7	8	9
key	a	b	ab	ba	aba	abb	bab	baa	abba	abbaa

- Original text = abababbabaabb**abba**abba
- Compressed text = 0122335

- $p = \text{abba}$ $p\text{Code} = 8$ Compressed text = 01223358
- $c = a$
- Enter **abbaa** into the code table.

code	0	1	2	3	4	5	6	7	8	9
key	a	b	ab	ba	aba	abb	bab	baa	abba	abbaa

- Original text = abababbabaabbabba**abba**
- Compressed text = 01223358

- **p = abba** **pCode = 8** Compressed text = 012233588
- **c = null**
- No need to enter anything to the table

LZW Decompression

code	0	1
key	a	b

- Original text = abababbabaabbabbaabba
- Compressed text = 012233588

- Convert codes to text from left to right.
- pCode=0 represents p=a DecompressedText = a

Expand

- `public static void expand() {`
- `String[] st = new String[L];`
- `int i; // next available codeword value`
- `// initialize symbol table with all 1-character strings`
- `for (i = 0; i < R; i++)`
- `st[i] = "" + (char) i;`
- `st[i++] = ""; // (unused) lookahead for EOF`
- `int codeword = BinaryStdIn.readInt(W);`
- `if (codeword == R) return; // expanded message is empty string`
- `String val = st[codeword];`
- `while (true) {`
- `BinaryStdOut.write(val);`
- `codeword = BinaryStdIn.readInt(W);`
- `if (codeword == R) break;`
- `String s = st[codeword];`
- `if (i == codeword) s = val + val.charAt(0); // special case hack`
- `if (i < L) st[i++] = val + s.charAt(0);`
- `val = s;`
- `}`
- `BinaryStdOut.close();`
- `}`

LZW Decompression

code	0	1	2
key	a	b	ab

- Original text = abababbabaabbabbaabba
- Compressed text = 012233588

- pCode=1 represents p = b DecompressedText = ab
- lastP = a followed by first character of p is entered into the code table (a | b)

code	0	1	2	3
key	a	b	ab	ba

- Original text = abababbabaabbabbaabba
- Compressed text = 012233588

- pCode = 2 represents p=ab DecompressedText = abab
- lastP = b followed by first character of p is entered into the code table (b | a)

code	0	1	2	3	4
key	a	b	ab	ba	aba

- Original text = abababbabaabbabbaabba
 - Compressed text = 012233588
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- pCode = 2 represents p = ab DecompressedText = ababab.
 - lastP = ab followed by first character of p is entered into the code table (ab | a)

code	0	1	2	3	4	5
key	a	b	ab	ba	aba	abb

- Original text = abababbabaabbabbaabba
- Compressed text = 012233588

- pCode = 3 represents p = ba DecompressedText = abababba.
- lastP = ab followed by first character of p is entered into the code table

code	0	1	2	3	4	5	6
key	a	b	ab	ba	aba	abb	bab

- Original text = abababbabaabbabbaabba
- Compressed text = 012233588

- pCode = 3 represents p = ba DecompressedText = abababbaba.
- lastP = ba followed by first character of p is entered into the code table.

code	0	1	2	3	4	5	6	7
key	a	b	ab	ba	aba	abb	bab	baa

- Original text = abababbabaabbabbaabba
- Compressed text = 012233588

- pCode = 5 represents p = abb DecompressedText = abababbabaabb
- lastP = ba followed by first character of p is entered into the code table.

code	0	1	2	3	4	5	6	7	8
key	a	b	ab	ba	aba	abb	bab	baa	abba

- Original text = abababbabaabbabbaabba
- Compressed text = 012233588

- 8 represents ???
- When a code is not in the table (then, it is the last one entered), and its key is lastP followed by first character of lastP
- lastP = abb
- So 8 represents p = abba Decompressed text = abababbabaabbabba

code	0	1	2	3	4	5	6	7	8	9
key	a	b	ab	ba	aba	abb	bab	baa	abba	abbaa

- Original text = abababbabaabbabbaabba
- Compressed text = 012233588

- pCode= 8 represents p=abba DecompressedText=abababbabaabbabbaabba.
- lastP = abba followed by first character of p is entered into the code table (abba | a)

Time Complexity

- Compression

Expected time = $O(n)$

where n is the length of the text that is being compressed.

- Decompression

$O(n)$, where n is the length of the decompressed text.