CS441 – Discrete Structures for Computer Science Instructor: Dr.Litman

Problems from Section 1.5

- 4. a) We have taken the conjunction of two propositions and asserted one of them. This is, according to Table 1, simplification.
 - b) We have taken the disjunction of two propositions and the negation of one of them, and asserted the other. This is, according to Table 1, disjunctive syllogism. See Table 1 for the other parts of this exercise as well.
 - c) modus ponens
- d) addition
- e) hypothetical syllogism
- 6. Let r be the proposition "It rains," let f be the proposition "It is foggy," let s be the proposition "The sailing race will be held," let l be the proposition "The life saving demonstration will go on," and let t be the proposition "The trophy will be awarded." We are given premises $(\neg r \lor \neg f) \to (s \land l)$, $s \to t$, and $\neg t$. We want to conclude r. We set up the proof in two columns, with reasons, as in Example 6. Note that it is valid to replace subexpressions by other expressions logically equivalent to them.

| Step | Reason |
|---|-------------------------------------|
| 1. ¬ <i>t</i> | Hypothesis |
| $2. s \rightarrow t$ | Hypothesis |
| 3. ¬s | Modus tollens using (1) and (2) |
| 4. $(\neg r \lor \neg f) \to (s \land l)$ | Hypothesis |
| 5. $(\neg(s \land l)) \rightarrow \neg(\neg r \lor \neg f)$ | Contrapositive of (4) |
| 6. $(\neg s \lor \neg l) \to (r \land f)$ | De Morgan's law and double negative |
| $7. \neg s \lor \neg l$ | Addition, using (3) |
| 8. $r \wedge f$ | Modus ponens using (6) and (7) |
| 9. r | Simplification using (8) |

- 14. In each case we set up the proof in two columns, with reasons, as in Example 6.
 - a) Let c(x) be "x is in this class," let r(x) be "x owns a red convertible," and let t(x) be "x has gotten a speeding ticket." We are given premises c(Linda), r(Linda), $\forall x(r(x) \to t(x))$, and we want to conclude $\exists x(c(x) \land t(x))$.

| Step | Reason |
|--|--------------------------------------|
| 1. $\forall x (r(x) \rightarrow t(x))$ | Hypothesis |
| 2. $r(Linda) \rightarrow t(Linda)$ | Universal instantiation using (1) |
| 3. $r(Linda)$ | Hypothesis |
| 4. $t(Linda)$ | Modus ponens using (2) and (3) |
| 5. $c(Linda)$ | Hypothesis |
| 6. $c(Linda) \wedge t(Linda)$ | Conjunction using (4) and (5) |
| 7. $\exists x (c(x) \land t(x))$ | Existential generalization using (6) |