

CS 1550

Week 8 – Lab 3
Interrupts
Part 2

Teaching Assistant Henrique Potter

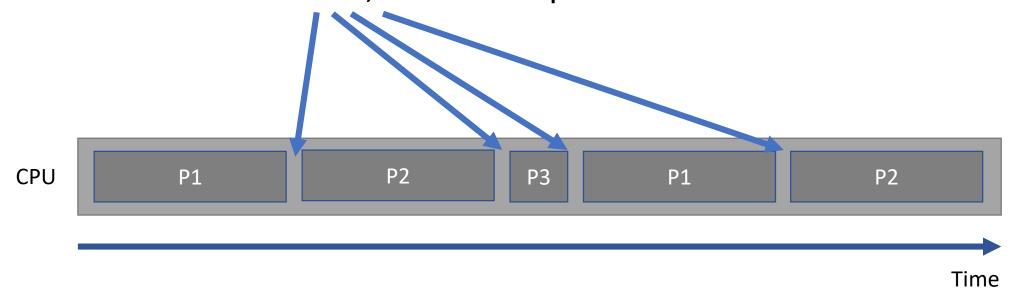
CS 1550 – Dues Dates

- Lab 3: Monday, March 9 @11:59pm
- Project 2: Tuesday, March 3 @11:59pm

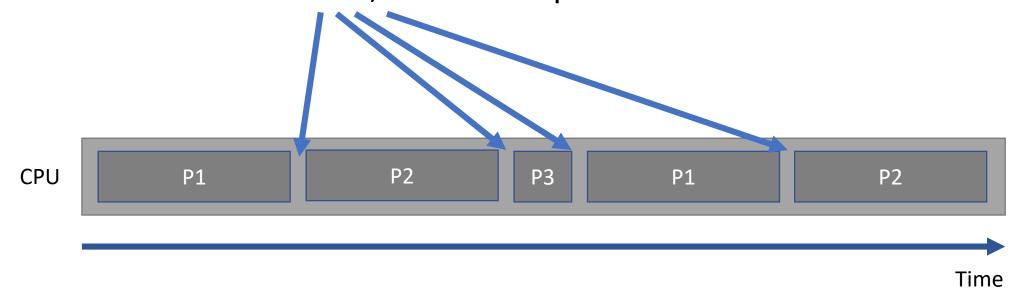
CS 1550 – Dues Dates

- Lab 3: Monday, March 9 @11:59pm
- Project 2: Tuesday, March 3 @11:59pm
- Project 2: Friday, March 6 @11:59pm

- In xv6, an interrupt for the scheduler is generated on every clock tick
- The scheduler is called, and a new process is selected

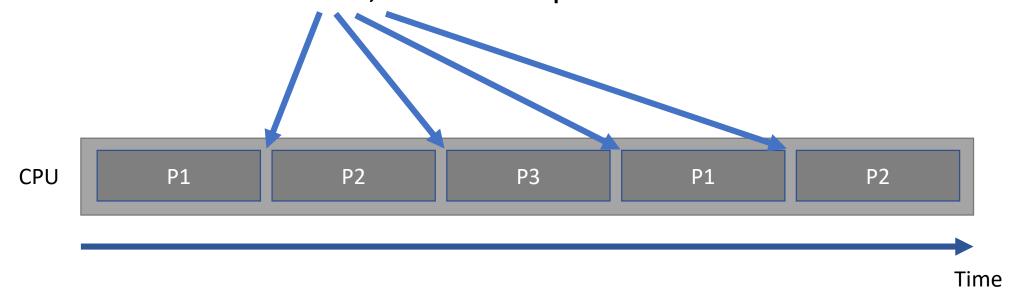


- In xv6, an interrupt for the scheduler is generated on every clock tick
- The scheduler is called, and a new process is selected



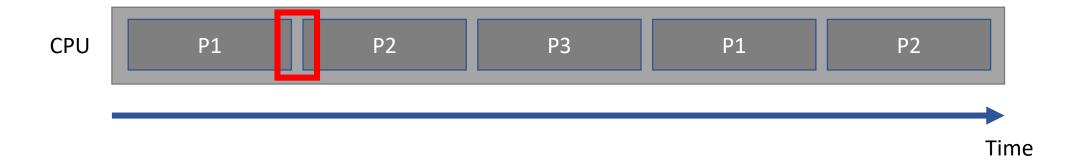
If the scheduler selects new processes in a round robin fashion. What's is wrong with this picture?

- In xv6, an interrupt for the scheduler is generated on every clock tick
- The scheduler is called, and a new process is selected

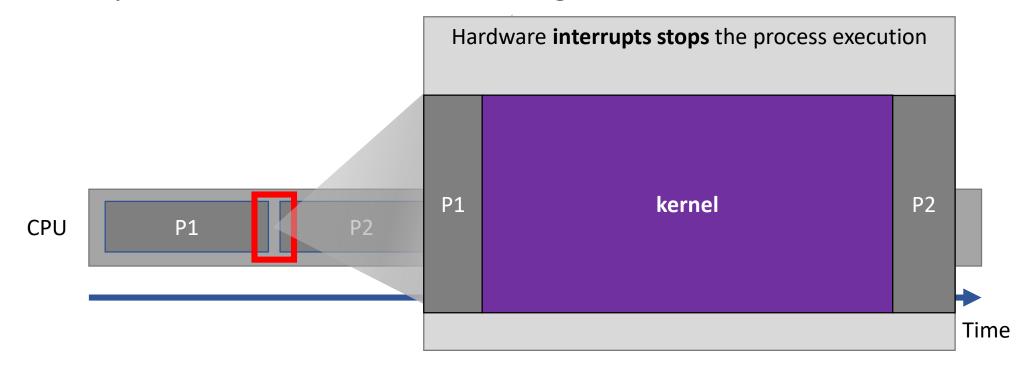


Processes should look more evenly distributed!

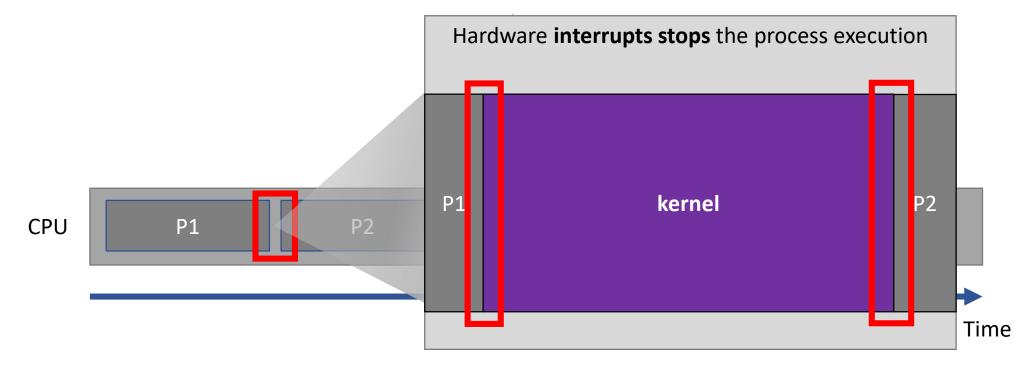
How processes are switched during their execution?



How processes are switched during their execution?



How processes are switched during their execution?



Let's take a deeper look at how the interrupts work!

• trapasm.S file

```
.globl alltraps
alltraps:
 pushl %ds
 pushl %es
 pushl %fs
 pushl %gs
 pushal
 movw $(SEG KDATA<<3), %ax
 movw %ax, %ds
 movw %ax, %es
 # Call trap(tf), where tf=%esp
 pushl %esp
 call trap
 addl $4, %esp
.globl trapret
trapret:
 popl %gs
 popl %fs
 popl %es
 popl %ds
  addl $0x8, %esp # trapno and errcode
  iret
```

```
//PAGEBREAK: 41
void
trap(struct trapframe *tf) ____
                                                Trapframe contains
                                                the process data
```

```
//PAGEBREAK: 41
void
trap(struct trapframe *tf)
₽ {
   if(tf->trapno == T SYSCALL) {
                                                 Call syscall! (Lab 1)
     if (myproc() ->killed)
        exit();
                                                  A user syscall cause a
     myproc() -> tf = tf;
                                                  interrupt!
     syscall();
     if (myproc() ->killed)
        exit();
     return;
```

```
//PAGEBREAK: 41
void
trap(struct trapframe *tf)
  if(tf->trapno == T SYSCALL) {
     if (myproc() ->ki lled)
syscall(void)
       exit();
     myproc()->tf =
                          int num;
                          struct proc *curproc = myproc();
     syscall();
     if (myproc() ->ki
                          num = curproc->tf->eax;
                          if(num > 0 && num < NELEM(syscalls) && syscalls[num]) {</pre>
       exit();
                            curproc->tf->eax = syscalls[num]();
     return;
                           } else {
                            cprintf("%d %s: unknown sys call %d\n",
                                    curproc->pid, curproc->name, num);
                            curproc->tf->eax = -1;
```

```
//PAGEBREAK: 41
trap(struct trapframe *tf)
 if(tf->trapno == T SYSCALL) {
    if (myproc() ->killed)
      exit();
    myproc() -> tf = tf;
    syscall();
    if (myproc() ->killed)
      exit();
    return;
  switch(tf->trapno) {
  case T IRQ0 + IRQ TIMER:
    if(cpuid() == 0){
      acquire (&tickslock);
      ticks++;
      wakeup(&ticks);
      release (&tickslock);
    lapiceoi();
    break;
```

```
//PAGEBREAK: 41
trap(struct trapframe *tf)
 if(tf->trapno == T SYSCALL) {
    if (myproc() ->killed)
      exit();
    myproc() -> tf = tf;
    syscall();
    if (myproc() ->killed)
      exit();
    return;
  switch(tf->trapno) {
 case T IRQ0 + IRQ TIMER:
    if(cpuid() == 0){
                                     Timer interrupt
      acquire(&tickslock);
      ticks++;
                                     Incrementing ticks
      wakeup(&ticks);
      release (&tickslock);
                                    Allow time keeping
    lapiceoi();
    break;
```

```
//PAGEBREAK: 41
trap(struct trapframe *tf)
 if(tf->trapno == T SYSCALL) {
    if (myproc() ->killed)
      exit();
    myproc() -> tf = tf;
    syscall();
    if (myproc() ->killed)
      exit();
    return;
  switch(tf->trapno) {
 case T IRQ0 + IRQ TIMER:
    if(cpuid() == 0){
      acquire (&tickslock);
      ticks++;
      wakeup(&ticks);
      release (&tickslock);
    lapiceoi();
    break;
```

```
switch(tf->trapno) {
case T_IRQ0 + IRQ_TIMER:
   if(cpuid() == 0) {
      acquire(&tickslock);
      ticks++;
      wakeup(&ticks);
      release(&tickslock);
   }
   lapiceoi();
   break;
```

```
switch(tf->trapno) {
case T IRQ0 + IRQ TIMER:
  if(cpuid() == 0){
    acquire (&tickslock);
    ticks++;
    wakeup(&ticks);
                                   Clock interrupts update cpu
    release (&tickslock);
                                   ticks and attempts to
                                   rescheduled a new process!
  lapiceoi();
 break;
if(myproc() && myproc()->state == RUNNING &&
   tf->trapno == T IRQ0+IRQ TIMER)
  yield();
```

```
if(myproc() && myproc()->state == RUNNING &&
    tf->trapno == T_IRQ0+IRQ_TIMER)
    yield();
```

Let's take a deeper look at yield()

```
if(myproc() && myproc()->state == RUNNING &&
    tf->trapno == T_IRQ0+IRQ_TIMER)
    yield();
```

```
if(myproc() && myproc()->state == RUNNING &&
    tf->trapno == T_IRQ0+IRQ_TIMER)
    yield();
```

```
// Give up the CPU for one scheduling round.
void
yield(void)
{
   acquire(&ptable.lock); //DOC: yieldlock
   myproc()->state = RUNNABLE;
   sched();
   release(&ptable.lock);
}
```

```
if(myproc() && myproc()->state == RUNNING &&
    tf->trapno == T_IRQ0+IRQ_TIMER)
    yield();
```

Why do we change the process state to **runnable**?

```
// Give up the CPU for one scheduling round.
void
yield(void)
{
   acquire(&ptable.lock); //DOC: yieldlock
   myproc()->state = RUNNABLE;
   sched();
   release(&ptable.lock);
}
```

```
if(myproc() && myproc()->state == RUNNING &&
    tf->trapno == T_IRQ0+IRQ_TIMER)
    yield();
```

```
// Give up the CPU for one scheduling round.
void
yield(void)
{
   acquire(&ptable.lock); //DOC: yieldlock
   myproc()->state = RUNNABLE;
   sched();
   release(&ptable.lock);
}
```

```
// Give up the CPU for one scheduling round.
void
yield(void)
{
   acquire(&ptable.lock); //DOC: yieldlock
   myproc()->state = RUNNABLE;
   sched();
   release(&ptable.lock);
}
```

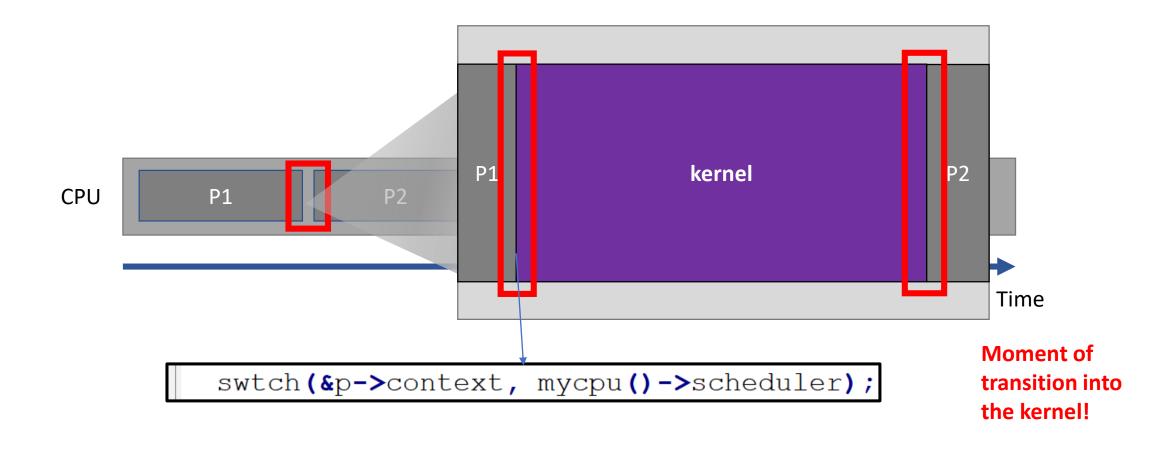
```
void
                 sched (void)
yield(void)
                   int intena;
                    struct proc *p = myproc();
 acquire (&ptable
 myproc()->state
                    if(!holding(&ptable.lock))
 sched();
 release (&ptable
                      panic("sched ptable.lock");
                    if(mycpu()->ncli != 1)
                      panic("sched locks");
                    if(p->state == RUNNING)
                      panic("sched running");
                    if(readeflags()&FL IF)
                      panic("sched interruptible");
                    <u>intena = mycpu()->intena;</u>
                   swtch(&p->context, mycpu()->scheduler);
                   mycpu()->intena = intena;
```

• proc.h file

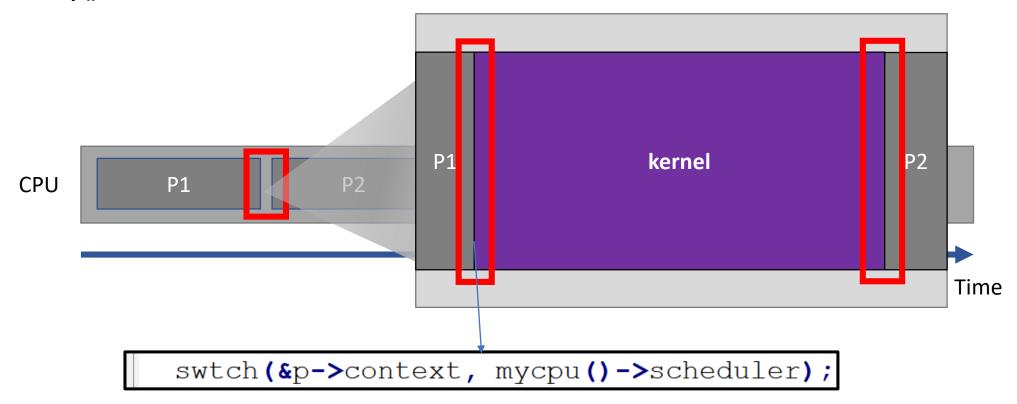
```
// Per-CPU state

struct cpu {
  uchar apicid;
  struct context *scheduler;
  struct taskstate ts;
  struct segdesc gdt[NSEGS];
  volatile uint started;
  int ncli;
  int intena;
  struct proc *proc;
};
```

```
// Per-process state
□struct proc {
                              // Size of process memory (bytes)
  uint sz;
                              // Page table
  pde t* pgdir;
  char *kstack;
                              // Bottom of kernel stack for this process
                              // Process state
  enum procstate state;
  int pid;
                              // Process ID
  struct proc *parent;
                              // Parent process
  struct trapframe *tf;
                              // Trap frame for current syscall
  struct context *context;
                              // swtch() here to run process
  void *chan;
                              // If non-zero, sleeping on chan
                              // If non-zero, have been killed
  int killed;
  struct file *ofile[NOFILE]; // Open files
  struct inode *cwd;  // Current directory
  char name[16];
                          // Process name (debugging)
  int get counts[23];  // Array for get count of syscall
```

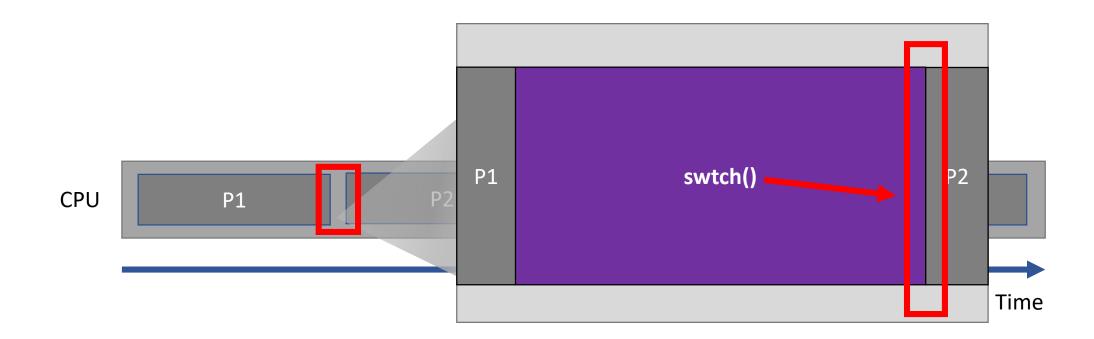


1. What was the context of the trap() execution?



• **proc.c** file

```
void
 scheduler (void)
₽ {
                                            This was executed
   struct proc *p;
                                            by the kernel
   struct cpu *c = mycpu();
  c \rightarrow proc = 0;
  for(;;){
     // Enable interrupts on this processor.
     sti();
     // Loop over process table looking for process to run.
     acquire (&ptable.lock);
     for(p = ptable.proc; p < &ptable.proc[NPROC]; p++){</pre>
       if(p->state != RUNNABLE)
         continue;
       // Switch to chosen process.
       c->proc = p;
       switchuvm(p);
       p->state = RUNNING;
       swtch(&(c->scheduler), p->context);
       switchkvm();
                                   But it was switched
                                   here
```



• proc.c file

The kernel starts

from here since it

stopped at the

previous line!

void

```
scheduler (void)
  struct proc *p;
 struct cpu *c = mycpu();
  c\rightarrow proc = 0;
 for(;;){
    // Enable interrupts on this processor.
    sti();
    // Loop over process table looking for process to run.
    acquire (&ptable.lock);
    for(p = ptable.proc; p < &ptable.proc[NPROC]; p++){</pre>
      if(p->state != RUNNABLE)
        continue;
      // Switch to chosen process.
      c \rightarrow proc = p;
      switchuvm(p);
      p->state = RUNNING;
      swtch(&(c->scheduler), p->context);
      switchkvm();
                                                    This loads the kernel's
                                                    information
      // Process is done running for now.
      c->proc = 0;
    release (&ptable.lock);
```

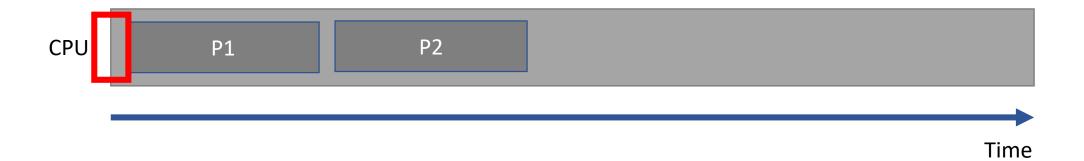
• proc.c file

If this loop is infinite and never breaks when did it start?

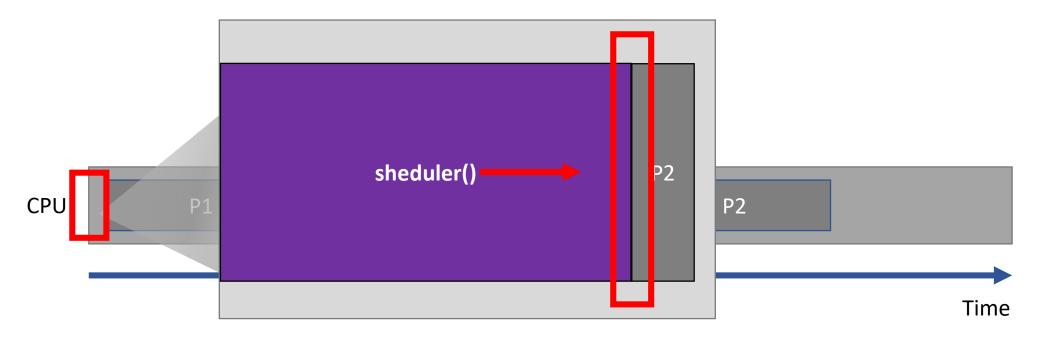
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scheduler (void)
  struct proc *p;
  struct cpu *c = mycpu();
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 for(;;){
    // Enable interrupts on this processor.
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    acquire (&ptable.lock);
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      if(p->state != RUNNABLE)
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      switchuvm(p);
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      switchkvm();
      // Process is done running for now.
      c\rightarrow proc = 0;
    release (&ptable.lock);
```



What is the first program the executes on system boot?

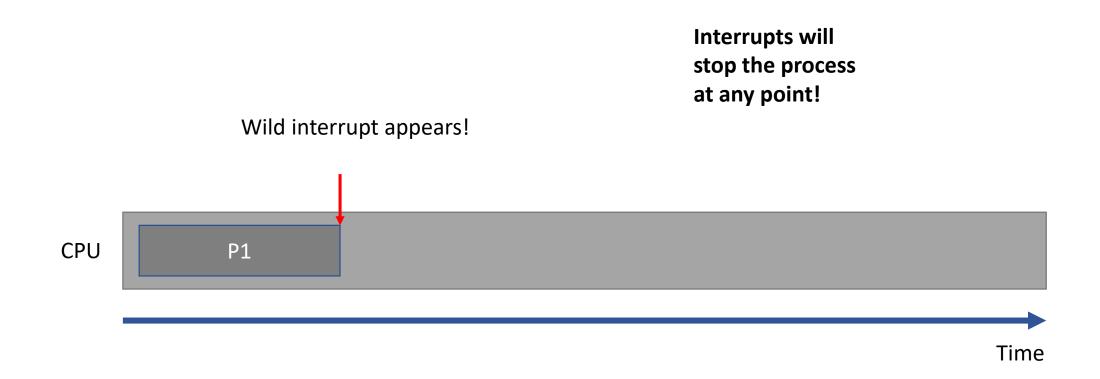


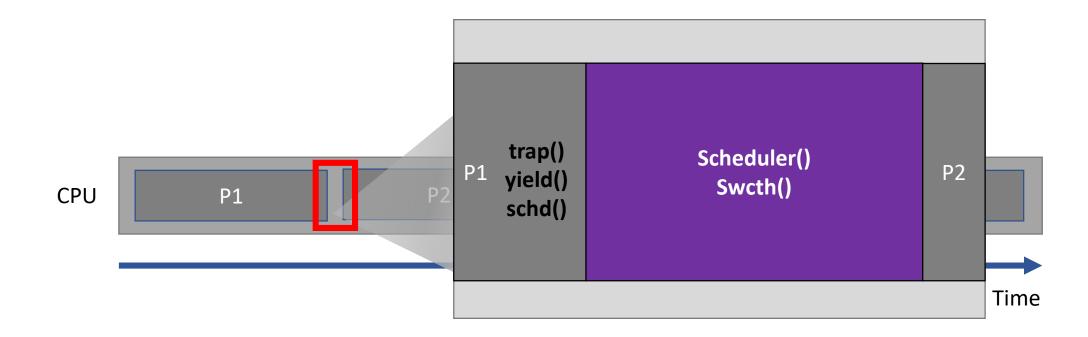
The scheduler started on the system boot process.



The scheduler eventually choses a process to run.







Lab 3 – Priority-based scheduler for XV6

- How are we adding the priority?
- Where we check to avoid an interrupt?

Lab 3 — Priority-based scheduler for XV6

How are we adding the priority?

Lab 3 – Priority-based scheduler for XV6

- The valid priority for a process is in the range of 0 to 200.
- The smaller value represents the higher priority.
- Default priority for a process is 50.
- proc.h:
 - Add an integer field called priority to struct proc.
- proc.c:
 - allocproc function:
 - Set the default priority for a process to 50
 - Scheduler function:
 - Replace the scheduler function with your implementation of a priority-based scheduler.

Lab 3 – part 2: add a syscall to set priority

- Add a new syscall, setpriority, for the process to change its priority.
- Changes the current process's priority and returns the old priority.
- Review lab1 to refresh steps to add a new syscall.



CS 1550

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Interrupts
Part 2

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