# Discrete Structures for Computer Science

William Garrison

bill@cs.pitt.edu 6311 Sennott Square

Lecture #23: Solving congruences



# Today's Topics

Arithmetic modulo n (reminder)

#### Solving linear congruences

- Modular inverses
- Extended Euclidean algorithm and Bézout numbers

#### Solving systems of congruences

Chinese remainder theorem

Primitive roots and discrete log

# Defining arithmetic restricted to remainders when dividing by m

 $\mathbf{Z}_m$  denotes the set of nonnegative integers less than m

 $\bullet$  i.e., the remainders when dividing by m

Recall that mod "preserves" addition and multiplication

- $(a+b) \mod m = ((a \mod m) + (b \mod m)) \mod m$
- $(ab) \mod m = ((a \mod m)(b \mod m)) \mod m$

Thus, we can define versions of addition and multiplication that are restricted to this set

- $\bullet \ a +_m b = (a+b) \bmod m$
- $a \cdot_m b = (a \cdot b) \mod m$
- These operations form arithmetic modulo m

Modular arithmetic behaves similarly to standard arithmetic: Recap properties from § 4.1

# REPORT OF THE PARTY OF THE PART

# Solving congruences via inverses

Consider the equation  $a + 8 \equiv 2 \pmod{11}$ 

- In standard arithmetic, we'd subtract 8 from both sides
  - ➣i.e., utilize the additive inverse

In modular arithmetic, additive inverses are easy to compute!

- $-8 \equiv 3 \pmod{11}$
- Thus, we can add 3 to both sides:

```
> a + 8 + 3 \equiv 2 + 3 (mod 11)
```

- $\gg$ a + 11  $\equiv$  5 (mod 11)
- $\gg$ a  $\equiv$  5 (mod 11)
- Note that adding any multiple of m preserves the value (mod m)

# Unfortunately, multiplicative inverses are not as simple

We cannot easily "divide by a" mod n

What is the equivalent of 1/a mod n?

Linear congruences are of the form  $ax \equiv b \pmod{m}$ 

- Given values for a and b, how do we solve for x?
- We need a value, say  $\bar{a}$ , where  $a\bar{a} \equiv 1 \pmod{m}$
- If we had this, we could multiply on both sides, then simplify!

Good news: Bézout's theorem says that there exist integers s and t such that gcd(a, m) = sa + tm

Assume a and m are coprime: How does this help us?



# **Extended Euclidean Algorithm**

The extended Euclidean algorithm computes the GCD of a and b, **and** computes the Bézout numbers s and t which satisfy the Bézout identity:

$$gcd(a, b) = sa + tb$$

| Row    | a         | b        | a/b    | a%b | d            | S           | t |
|--------|-----------|----------|--------|-----|--------------|-------------|---|
| 1      | 99        | 78       |        |     |              |             |   |
| Let th | nis repre | esent "a | div b" |     | "a <b>mo</b> | <b>d</b> b" |   |
| 3      |           |          |        |     |              |             |   |
| 4      |           |          |        |     |              |             |   |
| 5      |           |          |        |     |              |             |   |
| 6      |           |          |        |     |              |             |   |
|        |           |          |        |     |              |             |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   |     |   |   |   |
| 2   |    |    |     |     |   |   |   |
| 3   |    |    |     |     |   |   |   |
| 4   |    |    |     |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   |    |    |     |     |   |   |   |
| 3   |    |    |     |     |   |   |   |
| 4   |    |    |     |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   |    |    |     |     |   |   |   |
| 3   |    |    |     |     |   |   |   |
| 4   |    |    |     |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 |     |     |   |   |   |
| 3   |    |    |     |     |   |   |   |
| 4   |    |    |     |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   |     |   |   |   |
| 3   |    |    |     |     |   |   |   |
| 4   |    |    |     |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   |    |    |     |     |   |   |   |
| 4   |    |    |     |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 |     |     |   |   |   |
| 4   |    |    |     |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   |     |   |   |   |
| 4   |    |    |     |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   |    |    |     |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   | 15 | 6  |     |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   | 15 | 6  | 2   |     |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   | 15 | 6  | 2   | 3   |   |   |   |
| 5   |    |    |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   | 15 | 6  | 2   | 3   |   |   |   |
| 5   | 6  | 3  |     |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   | 15 | 6  | 2   | 3   |   |   |   |
| 5   | 6  | 3  | 2   |     |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   | 15 | 6  | 2   | 3   |   |   |   |
| 5   | 6  | 3  | 2   | 0   |   |   |   |
| 6   |    |    |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   | 15 | 6  | 2   | 3   |   |   |   |
| 5   | 6  | 3  | 2   | 0   |   |   |   |
| 6   | 3  | 0  |     |     |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   | 15 | 6  | 2   | 3   |   |   |   |
| 5   | 6  | 3  | 2   | 0   |   |   |   |
| 6   | 3  | 0  |     | _   |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   | 15 | 6  | 2   | 3   |   |   |   |
| 5   | 6  | 3  | 2   | 0   |   |   |   |
| 6   | 3  | 0  |     | _   |   |   |   |
|     |    |    |     |     |   |   |   |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  |   |   |   |
| 2   | 78 | 21 | 3   | 15  |   |   |   |
| 3   | 21 | 15 | 1   | 6   |   |   |   |
| 4   | 15 | 6  | 2   | 3   |   |   |   |
| 5   | 6  | 3  | 2   | 0   |   |   |   |
| 6   | 3  | 0  | _   |     | 3 | 1 | 0 |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  | 3 |   |   |
| 2   | 78 | 21 | 3   | 15  | 3 |   |   |
| 3   | 21 | 15 | 1   | 6   | 3 |   |   |
| 4   | 15 | 6  | 2   | 3   | 3 |   |   |
| 5   | 6  | 3  | 2   | 0   | 3 |   |   |
| 6   | 3  | 0  | _   | _   | 3 | 1 | 0 |

| Row | а  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  | 3 |   |   |
| 2   | 78 | 21 | 3   | 15  | 3 |   |   |
| 3   | 21 | 15 | 1   | 6   | 3 |   |   |
| 4   | 15 | 6  | 2   | 3   | 3 |   |   |
| 5   | 6  | 3  | 2   | 0   | 3 | K |   |
| 6   | 3  | 0  | _   | _   | 3 | 1 | 0 |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  | 3 |   |   |
| 2   | 78 | 21 | 3   | 15  | 3 |   |   |
| 3   | 21 | 15 | 1   | 6   | 3 |   |   |
| 4   | 15 | 6  | 2   | 3   | 3 |   |   |
| 5   | 6  | 3  | 2   | 0   | 3 | 0 |   |
| 6   | 3  | 0  | _   | _   | 3 | 1 | 0 |

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  | 3 |   |   |
| 2   | 78 | 21 | 3   | 15  | 3 |   |   |
| 3   | 21 | 15 | 1   | 6   | 3 |   |   |
| 4   | 15 | 6  | 2   | 3   | 3 |   |   |
| 5   | 6  | 3  | 2   | 0   | 3 | 0 |   |
| 6   | 3  | 0  |     | _   | 3 | 1 | 0 |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  | 3 |   |   |
| 2   | 78 | 21 | 3   | 15  | 3 |   |   |
| 3   | 21 | 15 | 1   | 6   | 3 |   |   |
| 4   | 15 | 6  | 2   | 3   | 3 |   |   |
| 5   | 6  | 3  | 2   | 0   | 3 | 0 | 1 |
| 6   | 3  | 0  |     | _   | 3 | 1 | 0 |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  | 3 |   |   |
| 2   | 78 | 21 | 3   | 15  | 3 |   |   |
| 3   | 21 | 15 | 1   | 6   | 3 |   |   |
| 4   | 15 | 6  | 2   | 3   | 3 |   |   |
| 5   | 6  | 3  | 2   | 0   | 3 | 0 | 1 |
| 6   | 3  | 0  |     | _   | 3 | 1 | 0 |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a  | b  | a/b | a%b | d | S | t |
|-----|----|----|-----|-----|---|---|---|
| 1   | 99 | 78 | 1   | 21  | 3 |   |   |
| 2   | 78 | 21 | 3   | 15  | 3 |   |   |
| 3   | 21 | 15 | 1   | 6   | 3 |   |   |
| 4   | 15 | 6  | 2   | 3   | 3 | 1 |   |
| 5   | 6  | 3  | 2   | 0   | 3 | 0 | 1 |
| 6   | 3  | 0  |     | _   | 3 | 1 | 0 |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a  | b  | a/b | a%b | d | S | t  |
|-----|----|----|-----|-----|---|---|----|
| 1   | 99 | 78 | 1   | 21  | 3 |   |    |
| 2   | 78 | 21 | 3   | 15  | 3 |   |    |
| 3   | 21 | 15 | 1   | 6   | 3 |   |    |
| 4   | 15 | 6  | 2   | 3   | 3 | 1 | -2 |
| 5   | 6  | 3  | 2   | 0   | 3 | 0 | 1  |
| 6   | 3  | 0  |     | _   | 3 | 1 | 0  |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a  | р  | a/b | a%b | d | S  | t  |
|-----|----|----|-----|-----|---|----|----|
| 1   | 99 | 78 | 1   | 21  | 3 |    |    |
| 2   | 78 | 21 | 3   | 15  | 3 |    |    |
| 3   | 21 | 15 | 1   | 6   | 3 | -2 |    |
| 4   | 15 | 6  | 2   | 3   | 3 | 1  | -2 |
| 5   | 6  | 3  | 2   | 0   | 3 | 0  | 1  |
| 6   | 3  | 0  |     | _   | 3 | 1  | 0  |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a  | b  | a/b | a%b | d | S  | t  |
|-----|----|----|-----|-----|---|----|----|
| 1   | 99 | 78 | 1   | 21  | 3 |    |    |
| 2   | 78 | 21 | 3   | 15  | 3 |    |    |
| 3   | 21 | 15 | 1   | 6   | 3 | -2 | 3  |
| 4   | 15 | 6  | 2   | 3   | 3 | 1  | -2 |
| 5   | 6  | 3  | 2   | 0   | 3 | 0  | 1  |
| 6   | 3  | 0  |     | _   | 3 | 1  | 0  |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a  | b  | a/b | a%b | d | S  | t  |
|-----|----|----|-----|-----|---|----|----|
| 1   | 99 | 78 | 1   | 21  | 3 |    |    |
| 2   | 78 | 21 | 3   | 15  | 3 | 3  |    |
| 3   | 21 | 15 | 1   | 6   | 3 | -2 | 3  |
| 4   | 15 | 6  | 2   | 3   | 3 | 1  | -2 |
| 5   | 6  | 3  | 2   | 0   | 3 | 0  | 1  |
| 6   | 3  | 0  | _   | _   | 3 | 1  | 0  |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a  | b  | a/b | a%b | d | S  | t   |
|-----|----|----|-----|-----|---|----|-----|
| 1   | 99 | 78 | 1   | 21  | 3 |    |     |
| 2   | 78 | 21 | 3   | 15  | 3 | 3  | -11 |
| 3   | 21 | 15 | 1   | 6   | 3 | -2 | 3   |
| 4   | 15 | 6  | 2   | 3   | 3 | 1  | -2  |
| 5   | 6  | 3  | 2   | 0   | 3 | 0  | 1   |
| 6   | 3  | 0  |     |     | 3 | 1  | 0   |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a  | b  | a/b | a%b | d | S   | t   |
|-----|----|----|-----|-----|---|-----|-----|
| 1   | 99 | 78 | 1   | 21  | 3 | -11 |     |
| 2   | 78 | 21 | 3   | 15  | 3 | 3   | -11 |
| 3   | 21 | 15 | 1   | 6   | 3 | -2  | 3   |
| 4   | 15 | 6  | 2   | 3   | 3 | 1   | -2  |
| 5   | 6  | 3  | 2   | 0   | 3 | 0   | 1   |
| 6   | 3  | 0  |     | _   | 3 | 1   | 0   |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a  | b  | a/b | a%b | d | S   | t   |
|-----|----|----|-----|-----|---|-----|-----|
| 1   | 99 | 78 | 1   | 21  | 3 | -11 | 14  |
| 2   | 78 | 21 | 3   | 15  | 3 | 3   | -11 |
| 3   | 21 | 15 | 1   | 6   | 3 | -2  | 3   |
| 4   | 15 | 6  | 2   | 3   | 3 | 1   | -2  |
| 5   | 6  | 3  | 2   | 0   | 3 | 0   | 1   |
| 6   | 3  | 0  |     |     | 3 | 1   | 0   |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

| Row | a   | Ь      | a/b   | a%b   | d      | S   | t   |
|-----|---|--------|-------|-------|--------|-----|-----|
| 1   | 99  | 78     | 1     | 21    | 3      | -11 | 14  |
| 2   | 78  | 21     | 3     | 15    | 3      | 3   | -11 |
| 3   | To ch   | neck v | our w | ork v | erify: | _   | 3   |
| 4   | To check your work, verify:<br>99*(-11) + 78*14 = 3 |        |       |       |        |     |     |
| 5   | 6   | 3      | 2     | 0     | 3      | 0   | 1   |
| 6   | 3   | 0      | _     | _     | 3      | 1   | 0   |

$$t = s_{previous} - \left(\frac{a}{b}\right) * t_{previous}$$

# Bézout numbers for modular inverses

If a and m are coprime, then gcd(a, m) = 1

The extended Euclidean algorithm yields:

- 1 = gcd(a, m) = sa + tm
- So sa = 1 tm
- Since  $km \equiv 0 \pmod{m}$  for any k, this means...
- sa  $\equiv$  1 (mod m)

This means that, when a and m are coprime, the Bézout numbers reveal a's (multiplicative) inverse mod m (!)

# PART OF THE PART O

## An example

Solve the following linear congruence:

$$57x \equiv 5 \pmod{98}$$

Using the extended Euclidean algorithm on 98 and 57, we can show that 98 \* (-25) + 57 \* 43 = 1, so 43 is the inverse of 57 (mod 98)

Multiply by 43 on both sides

- $57x * 43 \equiv 5 * 43 \pmod{98}$
- $x \equiv 215 \pmod{98}$
- $x \equiv 19 \pmod{98}$

# THE SECOND

# Solving systems of congruences

The Chinese Remainder Theorem: Let  $m_1$ ,  $m_2$ , ...,  $m_n$  be pairwise coprime positive integers greater than 1 and  $a_1$ ,  $a_2$ , ...,  $a_n$  arbitrary integers. Then the system:

- $x \equiv a_1 \pmod{m_1}$
- $x \equiv a_2 \pmod{m_2}$
- ...
- $x \equiv a_n \pmod{m_n}$

has a unique solution modulo  $m = m_1 m_2 ... m_n$ 

Let m be the product of the moduli, and let  $M_k$  be the product of all but the kth modulus

- Let  $y_k$  be the inverse of  $M_k$  (mod  $m_k$ )
- Now, compute  $x = a_1M_1y_1 + a_2M_2y_2 + ... + a_nM_ny_n$

# THE TOTAL PROPERTY OF THE PARTY OF THE PARTY

# Solving systems of congruences

In *Sunzi Suanjing*, the first known example of such problems was posed:

• 
$$x \equiv 2 \pmod{3}$$
  $m_1 = 3, a_1 = 2$ 

• 
$$x \equiv 3 \pmod{5}$$
  $m_2 = 5, a_2 = 3$ 

• 
$$x \equiv 2 \pmod{7}$$
  $m_3 = 7, a_3 = 2$ 

m = 3\*5\*7 = 105, and  $M_k$  is the product of all but the kth modulus

$$\bullet$$
 M<sub>1</sub> = 5\*7 = 35, M<sub>2</sub> = 3\*7 = 21, M<sub>3</sub> = 3\*5 = 15

•  $y_k$  is the inverse of  $M_k$  (mod  $m_k$ )

$$y_1 = 2 \text{ since } 35^2 = 70 \equiv 1 \pmod{3}$$

$$y_2 = 1 \text{ since } 21*1 = 21 \equiv 1 \pmod{5}$$

$$y_3 = 1 \text{ since } 15*1 = 15 \equiv 1 \pmod{7}$$

• Then,  $x = a_1M_1y_1 + a_2M_2y_2 + a_3M_3y_3$ 

$$> x = 2*35*2 + 3*21*1 + 2*15*1 = 233 \equiv 23 \pmod{105}$$

## **In-class** exercises

**Problem 1:** Find x where  $8x \equiv 3 \pmod{13}$ 

**Problem 2:** Find x where:

- $x \equiv 2 \pmod{3}$
- $x \equiv 3 \pmod{5}$
- $x \equiv 4 \pmod{7}$

# SEVERS TO SEVER SE

## Fermat's Little Theorem

**Theorem:** If p is prime and a is an integer not divisible by p, then  $a^{p-1} \equiv 1 \pmod{p}$ 

• This also means that  $a^p \equiv a \pmod{p}$ 

### Examples:

- Find 7<sup>222</sup> mod 11
  - Since 11 is prime,  $7^{10} \equiv 1 \pmod{11}$
  - Thus,  $7^{10k} \equiv 1 \pmod{11}$  for any integer k
  - So  $7^{220} \equiv 1 \pmod{11}$ , and  $7^{222} \equiv 7^2 \equiv 5 \pmod{11}$
  - $\bullet$  7<sup>222</sup> mod 11 = 5
- Find 5<sup>147</sup> **mod** 13
  - $5^{144} \equiv 1 \pmod{13}$  so  $5^{147} \equiv 5^3 \equiv 8 \pmod{13}$
  - $5^{147}$  mod 13 = 8

# SET THE SET OF THE SET

### Primitive roots

**Definition:** A primitive root modulo a prime p is an integer r in  $\mathbf{Z}_p$  such that every nonzero element of  $\mathbf{Z}_p$  is a power of r

- We sometimes call r a generator, since multiplying r by itself repeatedly can generate every element of  $\mathbf{Z}_p$
- There is a primitive root in Z<sub>D</sub> for every prime p

**Corollary:** If b is an integer in  $\mathbf{Z}_p$  and r is a primitive root modulo p, then there exists a unique exponent e in  $\mathbf{Z}_p$  such that  $\mathbf{r}^e = \mathbf{b}$ 

- i.e., re **mod** p = b
- Here, e is called the discrete log of b modulo p with base r
  > log<sub>r</sub> b = e (where the "mod p" is understood from context)



## The discrete logarithm problem

Given a prime p, a primitive root r modulo p, and a positive integer  $b \in \mathbf{Z}_p$ , find a value e such that  $r^e \mod p = b$ 

### How would you solve this?

No known algorithm in polynomial time

### Takeaways for solving congruences:

- We can invert addition with subtraction
- We can invert multiplication with modular inverses
- Inverting exponentiation is more difficult than it appears



## Final thoughts

- We can solve congruences by inverting operations, similar to standard algebra
  - To do so with multiplication, we use Euclid's algorithm and Bézout numbers to calculate multiplicative modular inverses
- The Chinese Remainder Theorem allows us to solve systems of congruences with coprime moduli
- Fermat's Little Theorem and primitive roots will come up again in cryptography
  - Section 4.5-4.6, next time!